

ON THE  $n$ -PARTITE TOURNAMENTS WITH EXACTLY  
 $n - m + 1$  CYCLES OF LENGTH  $m$

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**Abstract**

Gutin and Rafiey [*Multipartite tournaments with small number of cycles*, Australas J. Combin. 34 (2006) 17–21] raised the following two problems: (1) Let  $m \in \{3, 4, \dots, n\}$ . Find a characterization of strong  $n$ -partite tournaments having exactly  $n - m + 1$  cycles of length  $m$ ; (2) Let  $3 \leq m \leq n$  and  $n \geq 4$ . Are there strong  $n$ -partite tournaments, which are not themselves tournaments, with exactly  $n - m + 1$  cycles of length  $m$  for two values of  $m$ ? In this paper, we discuss the strong  $n$ -partite tournaments  $D$  containing exactly  $n - m + 1$  cycles of length  $m$  for  $4 \leq m \leq n - 1$ . We describe the substructure of such  $D$  satisfying a given condition and we also show that, under this condition, the second problem has a negative answer.

**Keywords:** multipartite tournaments, tournaments, cycles.

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1. INTRODUCTION

An  $n$ -partite or *multipartite tournament* is an orientation of a complete  $n$ -partite graph. A *tournament* is an  $n$ -partite tournament with exactly  $n$  vertices. A digraph  $D$  is *transitive* if, for every pair of arcs  $xy$  and  $yz$  in  $D$  such that  $x \neq z$ , the arc  $xz$  is also in  $D$ . It is easy to show that a tournament is transitive if and only if it is acyclic.

A digraph  $D$  is said to be *strong*, if for every pair of vertices  $x$  and  $y$ ,  $D$  contains a path from  $x$  to  $y$  and a path from  $y$  to  $x$ . A directed path from  $x$  to  $y$  in  $D$  is denoted by an  $(x, y)$ -path. An  $l$ -cycle is a cycle of length  $l$ . A cycle or path in a digraph  $D$  is Hamiltonian if it includes all the vertices of  $D$ .

In 1966, Moon discussed the number of  $m$ -cycle in a strong tournament.

**Theorem 1** (Moon [9]). *Let  $T$  be a strong tournament of order  $n$ . Then  $T$  contains at least  $n - m + 1$  cycles of length  $m$  for  $3 \leq m \leq n$ .*

The tournaments which gain the lower bound in Theorem 1 were characterized by Burzio and Demaria [2] for  $m = 3$ , Douglas [3] for  $m = n$  and Las Vergnas [8] for  $4 \leq m \leq n - 1$ . We list the result of Las Vergnas especially because we will use it to prove our main results.

**Theorem 2** (Las Vergnas [8]). *Every strong tournament of order  $n$  having exactly  $n - m + 1$  cycles of given length  $m$  with  $4 \leq m \leq n - 1$  is isomorphic to  $Q_n$ , where  $Q_n$  is a tournament of order  $n \geq 3$  obtained by reversing the arcs in the unique Hamiltonian path of a transitive tournament.*

In 2002, Volkmann extended Theorem 1 from tournaments to multipartite tournaments.

**Theorem 3** (Volkmann [10]). *Let  $D$  be a strong  $n$ -partite tournament. Then  $D$  contains at least  $n - m + 1$  cycles of length  $m$  for  $3 \leq m \leq n$ .*

It is notable that the bound in Theorem 3 is sharp which can be seen in [4] for  $n = 3$  and in [6] for  $4 \leq m \leq n$ . In addition, Gutin and Rafiey [6] raised the following two interesting and natural problems.

**Problem 4** (Gutin and Rafiey [6]). *Given  $m \in \{3, 4, \dots, n\}$ , find a characterization of strong  $n$ -partite tournaments having exactly  $n - m + 1$  cycles of length  $m$ .*

**Problem 5** (Gutin and Rafiey [6]). *Let  $3 \leq m \leq n$  and  $n \geq 4$ . Are there strong  $n$ -partite tournaments, which are not themselves tournaments, with exactly  $n - m + 1$  cycles of length  $m$  for two values of  $m$ ?*

Problem 4 seems to be especially interesting for the case  $m = n$  which was already solved by Gutin *et al.* in [7]. In this paper, we investigate strong  $n$ -partite tournaments  $D$ , which are not themselves tournaments and contain exactly  $n - m + 1$  cycles of length  $m$  for any given  $4 \leq m \leq n - 1$ . We prove that if  $D$  has an  $(n - 1)$ -cycle with no pair of vertices from the same partite set, then  $D$  must contain some given multipartite tournament as its subdigraph.

As far as Problem 5, Gutin and Rafiey [6] gave a negative answer for two values  $n - 1$  and  $n$  of  $m$ . This also implies that Problem 5 has a negative answer for  $n = 4$ . In this paper, we give a necessary condition to Problem 5 and show that if a strong  $n$ -partite tournament  $D$ , which is not itself a tournament, contains exactly  $n - m + 1$  cycles of length  $m$  for two values of  $m \in \{4, 5, \dots, n - 1\}$ , then there is no an  $(n - 1)$ -cycle with no pair of vertices from the same partite set in  $D$ .

2. TERMINOLOGY AND PRELIMINARIES

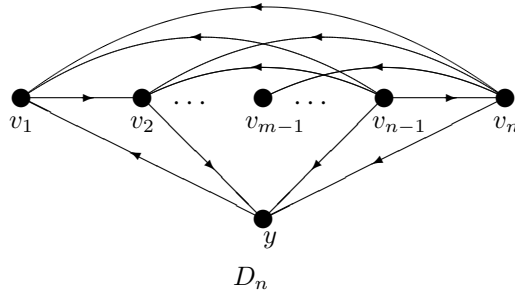
We assume that the reader is familiar with the standard terminology on digraphs and refer the reader to [1].

Let  $D$  be a digraph with the vertex set  $V(D)$  and the arc set  $A(D)$ . We call the number of vertices of  $D$  the *order* of  $D$ . A subdigraph induced by a subset  $A \subseteq V(D)$  is denoted by  $D\langle A \rangle$ . We use  $V(D) \setminus V(A)$  to stand for the set of vertices which are in  $V(D)$  but not in  $V(A)$ .

If  $xy$  is an arc in  $D$ , then we say that  $x$  *dominates*  $y$  and write  $x \rightarrow y$ . For two disjoint subsets  $X$  and  $Y$  of  $V(D)$ , if every vertex of  $X$  dominates every vertex of  $Y$ , we say  $X$  *dominates*  $Y$  and write  $X \rightarrow Y$ . Furthermore,  $X \Rightarrow Y$  denotes the property that there is no arc from  $Y$  to  $X$ .

The *out-neighborhood*  $N^+(x)$  of a vertex  $x$  is the set of vertices dominated by  $x$  and the *in-neighborhood*  $N^-(x)$  of a vertex  $x$  is the set of vertices dominating  $x$ . The numbers  $d^+(x) = |N^+(x)|$  and  $d^-(x) = |N^-(x)|$  are the *outdegree* and *indegree* of  $x$ , respectively. The *global irregularity* of  $D$  is defined as  $i_g(D) = \max\{\max\{d^+(x), d^-(x)\} - \min\{d^+(y), d^-(y)\} : x, y \in V(D)\}$ . We denote by  $D^{-1}$  the inverse digraph of  $D$ .

In order to present our main results, we define a class of  $n$ -partite tournaments  $D_n$  of order  $n + 1$  as described in the following figure, where  $3 \leq m \leq n$ ,  $\{v_2, \dots, v_{m-2}, v_m, \dots, v_n\} \rightarrow y \rightarrow v_1$ ,  $y$  and  $v_{m-1}$  belong to the same partite set and  $v_i \rightarrow v_j$  for all  $1 < j + 1 < i \leq n$ .



The following two theorems on cycles in strong  $n$ -partite tournaments are very useful to prove our main results.

**Theorem 6** (Guo and Volkmann [5]). *Every partite set of a strong  $n$ -partite tournament,  $n \geq 3$ , contains a vertex which lies on an  $m$ -cycle for each  $m \in \{3, 4, \dots, n\}$ .*

**Theorem 7** (Gutin and Rafiey [6]). *Let  $D$  be a strong  $n$ -partite tournament containing exactly  $n - m + 1$  cycles of length  $m$  for some  $m \in \{3, 4, \dots, n\}$ . Then every  $m$ -cycle of  $D$  has no pair of vertices from the same partite set.*

## 3. MAIN RESULTS

Before presenting the main results, we first prove the following lemma.

**Lemma 8.** *Let  $D$  be a strong  $n$ -partite tournament,  $n \geq 5$ , containing exactly  $n - m + 1$  cycles of length  $m$  for some  $3 \leq m \leq n - 1$ . If  $D$  has an  $(n - 1)$ -cycle  $C$  with no pair of vertices from the same partite set, then the following statements hold.*

- (a) *There are no two vertices  $u, w$  in  $V(D) \setminus V(C)$  such that  $C \Rightarrow u \rightarrow w \Rightarrow C$ .*
- (b) *There exists a vertex  $v \notin V(C)$  such that  $D \langle V(C) \cup \{v\} \rangle$  is strong.*

**Proof.** Let  $V_1, V_2, \dots, V_n$  be the partite sets of  $D$ . Suppose, without loss of generality, that  $C = v_1 v_2 \cdots v_{n-1} v_1$  with  $v_i \in V_i$ ,  $i = 1, 2, \dots, n - 1$ . By Theorem 1,  $D \langle V(C) \rangle$  contains at least  $(n - 1) - m + 1 = n - m$  cycles  $C_1, C_2, \dots, C_{n-m}$  of length  $m$ .

(a) Suppose to the contrary that there exist two vertices  $u, w \in V(D) \setminus V(C)$  such that  $C \Rightarrow u \rightarrow w \Rightarrow C$ . Obviously,  $u \notin V_n$  or  $w \notin V_n$ . Assume, without loss of generality, that  $w \in V_j$  for some  $1 \leq j \leq n - 1$ . Since  $n \geq 5$ , there exist at least two vertices  $v_k$  and  $v_t$ , such that  $u, w, v_k$  and  $v_t$  are in different partite sets. If  $m = 3$ , then  $u w v_k u$  and  $u w v_t u$  are two  $m$ -cycles different from  $C_1, C_2, \dots, C_{n-m}$ . This contradicts the fact that  $D$  contains exactly  $n - m + 1$  cycles of length  $m$ . If  $m \geq 4$ , then  $u w v_{j-1} v_j \cdots v_{j+m-4} u$  (if  $u \notin V_{j+m-4}$ ) or  $u w v_{j+1} v_{j+2} \cdots v_{j+m-2} u$  (if  $u \in V_{j+m-4}$ ) is an  $m$ -cycle with  $w, v_j \in V_j$  or  $u, v_{j+m-4} \in V_{j+m-4}$  (where all indices are modulo  $n - 1$ ). This is impossible by Theorem 7.

(b) Assume that there is no vertex  $v \in V(D) \setminus V(C)$  such that  $D \langle V(C) \cup \{v\} \rangle$  is strong. Let  $S = \{x \in V(D) \setminus V(C) : C \Rightarrow x\}$  and  $T = \{z \in V(D) \setminus V(C) : z \Rightarrow C\}$ . Since  $D$  is strong, we have that  $S$  and  $T$  are non-empty and there are vertices  $u \in S$  and  $w \in T$  such that  $u \rightarrow w$ . Thus, we have  $C \Rightarrow u \rightarrow w \Rightarrow C$ , which contradicts (a).  $\blacksquare$

**Theorem 9.** *Let  $D$  be a strong  $n$ -partite tournament which is not itself a tournament and contains exactly  $n - m + 1$  cycles of length  $m$  for some  $4 \leq m \leq n - 1$ . If  $D$  has an  $(n - 1)$ -cycle  $C$  with no pair of vertices from the same partite set, then  $D$  contains some  $D_i$  or  $D_i^{-1}$  as its subdigraph for  $i \in \{n - 1, n\}$ , where  $D_i$  is defined in Section 2.*

**Proof.** Let  $V_1, V_2, \dots, V_n$  be the partite sets of  $D$  and let  $C = v_1 v_2 \cdots v_{n-1} v_1$ ,  $v_i \in V_i$ ,  $i = 1, 2, \dots, n - 1$ . By Theorem 1,  $D \langle V(C) \rangle$  contains at least  $n - m$  cycles  $C_1, C_2, \dots, C_{n-m}$  of length  $m$ . By Theorem 6, there exists a vertex in  $V_n$ , say  $x$ , which lies on an  $m$ -cycle  $C_{n-m+1}$  different from  $C_1, C_2, \dots, C_{n-m}$ . We consider the following two cases.

*Case 1.*  $D \langle V(C) \cup \{x\} \rangle$  is not strong. Since  $D$  contains exactly  $n - m + 1$  cycles of length  $m$ , we have that  $D \langle V(C) \rangle$  contains exactly  $n - m$  cycles of length  $m$ .  $\blacksquare$

$m$ . By Theorem 2,  $D\langle V(C) \rangle$  is isomorphic to  $Q_{n-1}$ . So we may assume that  $v_i \rightarrow v_j$  for all  $1 < j + 1 < i \leq n - 1$ . Since  $D\langle V(C) \cup \{x\} \rangle$  is not strong, we have that  $C \rightarrow x$  or  $x \rightarrow C$ .

First we consider the case  $C \rightarrow x$ . Let  $S = \{u \in V(D) \setminus (V(C) \cup \{x\}) : D\langle V(C) \cup \{u\} \rangle \text{ is strong} \}$ . By Lemma 8(b),  $S$  is not empty. Since  $D$  is strong, there is a path from  $x$  to  $S$ . Let  $P = x_1 x_2 \cdots x_t$  ( $x_1 = x$ ) be such a path and assume that the  $P$  is of minimum length. That is,  $x_t \in S$  and  $D\langle V(C) \cup \{x_i\} \rangle$  is not strong for each  $i \in \{1, 2, \dots, t - 1\}$ . Since  $C \rightarrow x_1$  and  $x_1 \rightarrow x_2$ , we have  $x_2 \notin V(C)$ . If  $t > 2$ , then by Lemma 8(a), we have  $C \Rightarrow x_2$ . Successively, we can get that  $x_i \notin V(C)$  and  $C \Rightarrow x_i$  for all  $i \in \{2, 3, \dots, t - 1\}$  when  $t > 2$ .

If there exist two vertices  $v_i, v_j$  on  $C$  such that  $x_t \rightarrow \{v_i, v_j\}$ , then, when  $t > m - 1$ , we have that  $x_t v_i x_{t-(m-2)} x_{t-(m-3)} \cdots x_t$  (if  $x_{t-(m-2)} \notin V_i$ ) or  $x_t v_j x_{t-(m-2)} x_{t-(m-3)} \cdots x_t$  (if  $x_{t-(m-2)} \in V_i$ ) is an  $m$ -cycle different from  $C_1, C_2, \dots, C_{n-m+1}$ , a contradiction; when  $t = m - 1$ , it is clear that  $x_t v_i x_1 \cdots x_t$  and  $x_t v_j x_1 \cdots x_t$  are two  $m$ -cycles different from  $C_1, C_2, \dots, C_{n-m}$ , a contradiction; when  $t \leq m - 2$ , it is easy to see that  $x_t v_i v_{i+1} \cdots v_{i+(m-1-t)} x_1 \cdots x_t$  and  $x_t v_j v_{j+1} \cdots v_{j+(m-1-t)} x_1 \cdots x_t$  are two  $m$ -cycles different from  $C_1, C_2, \dots, C_{n-m}$ , a contradiction.

Therefore,  $x_t$  has only one out-neighbor on  $C$ . We will show that  $x_t \rightarrow v_1$ . In fact, if  $x_t \rightarrow v_i$  and  $i \geq 2$ , then we have that  $x_t v_i \cdots v_{i+m-2} x_t$  (when  $i + m - 2 \leq n - 1$  and  $x_t \notin V_{i+m-2}$ ) or  $x_t v_i \cdots v_{i+m-3} v_{i-1} x_t$  (when  $i + m - 2 \leq n - 1$  and  $x_t \in V_{i+m-2}$ ) or  $x_t v_i \cdots v_{n-1} v_1 \cdots v_{m-n+i-1} x_t$  (when  $i + m - 2 \geq n$  and  $x_t \notin V_{m-n+i-1}$ ) or  $x_t v_i \cdots v_{n-1} v_2 \cdots v_{m-n+i} x_t$  (when  $i + m - 2 \geq n$  and  $x_t \in V_{m-n+i-1}$ ) is an  $m$ -cycle different from  $C_1, C_2, \dots, C_{n-m+1}$ , a contradiction. So we have  $\{v_2, v_3, \dots, v_{n-1}\} \Rightarrow x_t$ . Furthermore, if  $v_{m-1} \rightarrow x_t$ , then  $x_t v_1 v_2 \cdots v_{m-1} x_t$  is an  $m$ -cycle different from  $C_1, C_2, \dots, C_{n-m+1}$ , a contradiction. So  $x_t \in V_{m-1}$ . Let  $x_t = y$ . Then  $D$  contains  $D_{n-1}$  as its subdigraph.

For the case  $x \rightarrow C$ , by considering the inverse of  $D$ , it is easy to see that  $D$  contains  $D_{n-1}^{-1}$  as its subdigraph.

*Case 2.*  $D\langle V(C) \cup \{x\} \rangle$  is strong. In this case,  $D\langle V(C) \cup \{x\} \rangle$  is a strong tournament of order  $n$ . By Theorem 1,  $D\langle V(C) \cup \{x\} \rangle$  contains at least  $n - m + 1$  cycles of length  $m$ . Note that  $D$  contains exactly  $n - m + 1$  cycles of length  $m$ . We have that  $D\langle V(C) \cup \{x\} \rangle$  contains exactly  $n - m + 1$  cycles of length  $m$ . By Theorem 2,  $D\langle V(C) \cup \{x\} \rangle$  is isomorphic to  $Q_n$ . So we may assume that  $C' = v_1 v_2 \cdots v_n v_1$  is an  $n$ -cycle of  $D\langle V(C) \cup \{x\} \rangle$  satisfying  $v_i \in V_i$  and  $v_i \rightarrow v_j$  for all  $1 < j + 1 < i \leq n$ . Obviously,  $C_1 = v_1 v_2 \cdots v_m v_1, C_2 = v_2 v_3 \cdots v_{m+1} v_2, \dots, C_{n-m+1} = v_{n-m+1} v_{n-m+2} \cdots v_n v_{n-m+1}$  are  $n - m + 1$  cycles of length  $m$  of  $D$ .

**Claim 10.** *There exists a vertex  $y \in V(D) \setminus V(C')$  such that  $D\langle V(C') \cup \{y\} \rangle$  is strong.*

**Proof.** Assume that there is no vertex  $y \in V(D) \setminus V(C')$  such that  $D\langle V(C') \cup \{y\} \rangle$  is strong. Let  $S = \{x \in V(D) \setminus V(C') : C' \Rightarrow x\}$  and  $T = \{z \in V(D) \setminus V(C') : z \Rightarrow V(C')\}$ . Since  $D$  is strong, we have that  $S$  and  $T$  are non-empty and there are vertices  $u \in S$  and  $w \in T$  such that  $u \rightarrow w$ . Suppose that  $u \in V_i$  and  $w \in V_j$  for  $1 \leq i \neq j \leq n$ . Then  $u w v_{j+1} v_{j+2} \cdots v_{j+m-2} u$  (if  $i \neq j + m - 2$ ) or  $u w v_{j+2} v_{j+3} \cdots v_{j+m-1} u$  (if  $i = j + m - 2$ ) is an  $m$ -cycle different from  $C_1, C_2, \dots, C_{n-m+1}$ . Note that  $D$  contains exactly  $n - m + 1$  cycles of length  $m$ . This is a contradiction.  $\square$

By Claim 10, there are two vertices  $v_a, v_b$  ( $1 \leq a, b \leq n$ ), such that  $v_a \rightarrow y \rightarrow v_b$ . Assume that  $v_k$  is the first vertex from  $v_1$  to  $v_n$  dominating  $y$ .

**Claim 11.**  $v_i \Rightarrow y$  for all  $k \leq i \leq n$ .

**Proof.** Otherwise, there exists some index  $t$  such that either  $v_t \rightarrow y \rightarrow v_{t+1}$  ( $k \leq t \leq n - 1$ ) or  $y, v_{t+1} \in V_{t+1}$  but  $v_t \rightarrow y \rightarrow v_{t+2}$  ( $k \leq t \leq n - 2$ ). We still assume that  $t$  is such a minimum index.

If  $t \leq n - m + 1$ , then either  $v_t y v_{t+1} \cdots v_{t+m-2} v_t$  or  $v_t y v_{t+2} \cdots v_{t+m-1} v_t$  is an  $m$ -cycle different from  $C_1, C_2, \dots, C_{n-m+1}$ , a contradiction.

If  $n - m + 2 \leq t \leq n - 2$ , then either  $v_t y v_{t+1} \cdots v_n v_{n-m+2} \cdots v_t$  or  $v_t y v_{t+2} \cdots v_n v_{n-m+1} \cdots v_t$  is an  $m$ -cycle different from  $C_1, C_2, \dots, C_{n-m+1}$ , a contradiction.

If  $t = n - 1$ , then  $y \rightarrow v_n$  and  $v_{n-1} y v_n v_{n-m+2} \cdots v_{n-1}$  is an  $m$ -cycle different from  $C_1, C_2, \dots, C_{n-m+1}$ , a contradiction.  $\square$

**Claim 12.**  $y \rightarrow v_1$ .

**Proof.** If  $y \in V_1$ , then  $y \rightarrow v_2$  (otherwise,  $k = 2$  and  $\{v_2, v_3, \dots, v_n\} \rightarrow y$  by Claim 11, which contradicts the assumption that  $D\langle V(C') \cup \{y\} \rangle$  is strong). By Claim 11, we have  $v_n \rightarrow y$ . Therefore,  $D\langle v_2, \dots, v_n, y \rangle$  is a strong tournament. Then  $y$  is in an  $m$ -cycle of  $D\langle v_2, \dots, v_n, y \rangle$ , which is different from  $C_1, C_2, \dots, C_{n-m+1}$ , a contradiction.

Therefore,  $y \notin V_1$ . If  $v_1 \rightarrow y$ , then  $\{v_2, v_3, \dots, v_n\} \Rightarrow y$  by Claim 11, which contradicts the assumption that  $D\langle V(C') \cup \{y\} \rangle$  is strong. So we have  $y \rightarrow v_1$ .  $\square$

By Claim 12, we have that  $2 \leq k \leq n$  and  $y \Rightarrow v_{m-1}$ . Otherwise,  $y v_1 v_2 \cdots v_{m-1} y$  is an  $m$ -cycle different from  $C_1, C_2, \dots, C_{n-m+1}$ , a contradiction.

If  $k = 2$ , then by  $m \geq 4$  and Claim 11, we have  $y \in V_{m-1}$ , and hence,  $\{v_2, v_3, \dots, v_n\} \Rightarrow y \rightarrow v_1$ . Now,  $D$  contains  $D_n$  as its subdigraph.

If  $2 < k < m - 1$ , then  $y \in V_{m-1}$ ,  $y \rightarrow v_2$  and  $v_m \rightarrow y$ . Thus,  $y v_2 \cdots v_m y$  is an  $m$ -cycle different from  $C_1, C_2, \dots, C_{n-m+1}$ , a contradiction.

If  $m - 1 \leq k \leq n - 1$ , then  $1 \leq k - m + 2 \leq k - 2$  and  $y \Rightarrow v_{k-m+2}$ . Now  $v_k y v_{k-m+2} \cdots v_k$  (if  $y \rightarrow v_{k-m+2}$ ) or  $v_{k+1} y v_{k-m+3} \cdots v_{k+1}$  (if  $y \in V_{k-m+2}$ ) is an  $m$ -cycle different from  $C_1, C_2, \dots, C_{n-m+1}$ , a contradiction.

If  $k = n$ , then  $v_n \rightarrow y \Rightarrow \{v_1, v_2 \cdots v_{n-1}\}$  by the choice of  $k$ . It is easy to see that  $y \in V_{n-m+2}$ , as otherwise  $yv_{n-m+2} \cdots v_n y$  is an  $m$ -cycle different from  $C_1, C_2, \dots, C_{n-m+1}$ , a contradiction. Now  $D$  contains  $D_n^{-1}$  as its subdigraph. ■

**Theorem 13.** *Let  $D$  be a strong  $n$ -partite tournament,  $n \geq 5$ , which is not itself a tournament. If  $D$  contains an  $(n - 1)$ -cycle with no pair of vertices from the same partite set, then  $D$  does not contain exactly  $n - m + 1$  cycles of length  $m$  for two values of  $m \in \{4, 5, \dots, n - 1\}$ .*

**Proof.** Let  $m$  and  $m_1$  be two distinct values from the set  $\{4, 5, \dots, n - 1\}$  and assume that  $D$  has exactly  $n - m + 1$  cycles of length  $m$ . Let  $V_1, V_2, \dots, V_n$  be the partite sets of  $D$ . By Theorem 9,  $D$  contains some  $D_i$  or  $D_i^{-1}$  as its subdigraph for  $i \in \{n - 1, n\}$ .

If  $D$  contains  $D_{n-1}$  ( $D_{n-1}^{-1}$ ) as its subdigraph, then let  $C = v_1 v_2 \cdots v_{n-1} v_1$  be an  $(n - 1)$ -cycle of  $D_{n-1}$  ( $D_{n-1}^{-1}$ ) with  $v_i \in V_i$  ( $i = 1, 2, \dots, n - 1$ ),  $v_i \rightarrow v_j$  for all  $1 < j + 1 < i \leq n - 1$ ,  $y \in V_{m-1}$ ,  $\{v_2, v_3, \dots, v_{n-1}\} \Rightarrow y \rightarrow v_1$  ( $y \in V_{n-m+1}$  and  $v_{n-1} \rightarrow y \Rightarrow \{v_1, v_2, \dots, v_{n-2}\}$ ). By Theorem 1,  $D\langle V(C) \rangle$  contains at least  $(n - 1) - m_1 + 1 = n - m_1$  cycles of length  $m_1$ . Note that  $yv_1 v_2 \cdots v_{m_1-1} y$  ( $v_{n-1} y v_{n-(m_1-1)} v_{n-(m_1-2)} \cdots v_{n-1}$ ) is another  $m_1$ -cycle of  $D_{n-1}$  ( $D_{n-1}^{-1}$ ). In addition, there exists a vertex in  $V_n$ , say  $x$ , which is in an  $m_1$ -cycle of  $D$  different from the above  $m_1$ -cycles. Thus,  $D$  contains at least  $n - m_1 + 2$  cycles of length  $m_1$ .

If  $D$  contains  $D_n$  ( $D_n^{-1}$ ) as its subdigraph, then let  $C = v_1 v_2 \cdots v_n v_1$  be an  $n$ -cycle of  $D_n$  ( $D_n^{-1}$ ) with  $v_i \in V_i$  ( $i = 1, 2, \dots, n$ ),  $v_i \rightarrow v_j$  for all  $1 < j + 1 < i \leq n$ ,  $y \in V_{m-1}$ ,  $\{v_2, v_3, \dots, v_n\} \Rightarrow y \rightarrow v_1$  ( $y \in V_{n-m+2}$ ,  $v_n \rightarrow y \Rightarrow \{v_1, v_2, \dots, v_{n-1}\}$ ). By Theorem 1,  $D\langle V(C) \rangle$  contains at least  $n - m_1 + 1$  cycles of length  $m_1$ . It is easy to see that  $yv_1 v_2 \cdots v_{m_1-1} y$  ( $v_n y v_{n-(m_1-2)} v_{n-(m_1-3)} \cdots v_n$ ) is another  $m_1$ -cycle of  $D_n$  ( $D_n^{-1}$ ). Then  $D$  contains at least  $n - m_1 + 2$  cycles of length  $m_1$ . The theorem is complete. ■

In 2004, Winzen [11] showed that an  $n$ -partite tournament  $D$  with  $n \geq 4$  and  $i_g(D) \leq 2$  contains a strong subtournament of order  $p$  for every  $p \in \{3, 4, \dots, n - 1\}$ . So  $D$  contains an  $(n - 1)$ -cycle with no pair of vertices from the same partite set, which yields the following result.

**Corollary 14.** *If  $D$  is a strong  $n$ -partite tournament with  $n \geq 5$  and  $i_g(D) \leq 2$ , which is not itself a tournament, then  $D$  does not contain exactly  $n - m + 1$  cycles of length  $m$  for two values of  $m \in \{4, 5, \dots, n - 1\}$ .*

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