DÜRER POLYHEDRA: THE DARK SIDE OF MELANCHOLIA

PATRICK W. FOWLER
School of Chemistry, University of Exeter
Stocker Rd, Exeter EX4 4QD, UK

AND

PETER E. JOHN
Institute of Mathematics
Technical University of Ilmenau
Ilmenau, Germany

Abstract

Dürer's engraving Melencolia I famously includes a perspective view of a solid polyhedral block of which the visible portion is an 8-circuit bounding a pentagon-triangle patch. The polyhedron is usually taken to be a cube truncated on antipodal corners, but an infinity of others are compatible with the visible patch. Construction of all cubic polyhedra compatible with the visible portion (i.e., Dürer Polyhedra) is discussed, explicit graphs and symmetries are listed for small cases (≤ 18 vertices) and total counts are given for 10 ≤ vertices ≤ 26.

Keywords: graph theory, geometry, chemistry.
2000 Mathematics Subject Classification: 05C30, 05C90, 52B10, 92E10.

1. Introduction

How well can we reconstruct a simple geometrical object from a single snapshot in which one hemisphere is lit and the other is in darkness? In Albrecht Dürer's famous engraving of 1514, Melencolia I, a winged figure broods over a collection of objects that includes a large, strikingly irregular, polyhedral stone block described variously by art historians as a rhomboid [1],
a truncated rhomboid or truncated rhombohedron [2] (‘abgestumpfter Rhomboeder’ [3]). Four faces only of the block are visible (a triangle and three pentagons). It is suggested [4] that the whole is a cube truncated on opposite vertices, a claim supported by a view of a truncated cube in Dürer’s writings on perspective geometry [5], and a preliminary drawing in the Dresden Sketchbook [6] with hidden lines marked. The antipodally bi-truncated cube is sometimes called ‘Dürer’s octahedron’ in the mathematical literature [7, 8], but this identification is clearly combinatorial rather than affine in nature, as the angles of the figure as drawn are famously irregular, even after allowance is made for perspective effects [9, 10, 11].

Here we ask another combinatorial question: what can we deduce with certainty about the hidden dark side of the mysterious polyhedron, simply from what we can see in the engraving? An infinite set of completions, each equally deserving the appellation of Dürer Polyhedron, is found. As the reasoning has applications to problems of blind-side reconstruction of polyhedra in other contexts, and uses only elementary graph theory, it is briefly summarised here along with the specific solutions for Dürer polyhedra.

2. Method

Assume that the object to be reconstructed is a cubic polyhedron $P$ without holes or handles (i.e., that the graph of vertices and edges has three edges meeting at every vertex, is three-connected and can be drawn in the plane without edge crossings [12]). Note that in [12, p. 235] is also given a proof of the famous Steinitz theorem [13, 14]:

A graph $G$ is (isomorphic to) the graph $G(P)$ of a cubic polyhedron $P$ if and only if it is planar and three-connected;

which is important for our investigation.

$P^0$ and $P^*$ are respectively the visible and invisible parts of $P$, separated by a ‘terminator’ cycle $C^0 = C(P^0)$. Let $P^*$ denote the union of $P^1$ and $C^0$. The reconstruction problem is converted to one of finding suitable in-fillings of a Schlegel diagram of $P^0$ by using graphs embedded in the plane i.e., $G = G(P), G^0 = G(P^0), G^* = G(P^*)$. The terminator becomes the circuit $C^0 = C(P^0)$ bounding a finite face of $G^0$. Symbols $v, e, f$ with appropriate superscripts denote the numbers of vertices, edges and faces of $G, G^0$ and $G^*$. Clearly, $v^* = v - v^0, e^* = e - e^0$ and $f^* = f - f^0$ and, as $P$ is cubic,
\[ e^* = \frac{3v^*}{2} + \frac{3v^0}{2} - e^0, \]
\[ f^* = \frac{v^*}{2} + \frac{v^0}{2} + 2 - f^0. \]

In our case, \( P \) is a Dürer polyhedron, \( P = D \) (i.e., one that contains the triangle + triple-pentagon patch seen in the engraving). \( G(D^0) \) has \( v^0 = 10 \), \( e^0 = 13 \), \( f^0 = 4 \), with four divalent vertices in the circuit \( C^1 \), and hence \( e^* \geq 2 \), \( f^* \geq 3 \) for \( v^* \geq 0 \), and, if \( \tilde{f}_r \) denotes the number of faces of size \( r \), \( \sum_r (6 - r)\tilde{f}_r = 12 \), thus \( \sum_r (6 - r)f^*_r = 6 \). The in-filling problem is sketched in Figure 1.

Figure 1. Dürer polyhedron and its Schlegel diagram. Regions \( \text{III} \) and \( V \) denote the triangular and pentagonal faces of the visible hemisphere of the polyhedron. ● are saturated trivalent vertices; ○ have two valencies in the visible portion, requiring incidence with an edge within the hidden hemisphere (the shaded circular disk).

A systematic procedure for generating all solutions is now described.

Let \( G = G(P) \) with \( C^0 = C(G^0) \) be a solution of our problem. Delete all edges that are incident with vertices of degree 3 in \( G^0 \), and after that, all isolated vertices. The remaining graph \( B(G, C^0) \) has \( c = c(B(G, C^0)) \) components \( B_i \), where \( 1 \leq c \leq 2 \). Each component \( B_i \) is an internal bridge of \( G \) with respect to \( C^0 \). Let \( B_j = B_j(G, C^0) \) denote the set of all bridges \( B \) for which the number of vertices of degree 1 is \( b^0 = b^0(B) = j \). Clearly, in our case is \( j = 2 \) or 4.
For \( j = 2 \) there exists exactly one solution (two bridges \( B_1, B_2 \) isomorphic to the complete graph \( K_2 \)). See Figure 2.

\[
\begin{array}{|c|c|c|}
\hline
\text{two bridges} & \text{one bridge} & \text{one bridge} \\
\hline
s^* = -1 & s^* = 0 & s^* = +1 \\
\hline
\end{array}
\]

Figure 2. Possible bridges for \( j = 4 \). The complete set of possibilities with \( s^* \leq 1 \).

For \( j = 4 \) there exist one-bridge solutions \( B \), which have \( v^* = v^*(B) \) vertices of degree 3, \( e^* = e^*(B) \) edges, and the number \( f^* = f^*(B) \) of faces of \( B \) is defined by \( f^* := f(G^*) - 1 \). Clearly, \( 3v^* = 2e^* - b^0 \). The general solutions are therefore, for integer \( s^* \geq -b^0/2 \):

\[
\begin{align*}
e^* & = 3s^* + 2b^0, \\
v^* & = 2s^* + b^0, \\
f^* & = s^* + b^0 + 1.
\end{align*}
\]

For \( j = 4 \) and small integers \( s^* \), the possible elements of \( \mathcal{B}_4 \) are listed in Figure 2.

Small cases are easily checked by working through published complete lists of graphs sorted by edge or vertex count [15] or generating new lists [16].

To solve our problem of finding polyhedra, we must take into account the sequence of vertices in \( C^0 \) in order to decide on the number and placing of the bridges within \( C^0 \).
The difficulty with the present approach is in assigning symmetry and checking isomorphism for the polyhedra generated by the distinct placements of the bridges within the Schegel diagram. The graph $G^0 = G(D^0)$ has $C_4$ symmetry, which may be broken, enhanced or reduced on filling in the bridge(s). Hand enumeration of possibilities for bridges gives the set of the first 78 Dürer polyhedra shown in Figure 3, listed in order of increasing $v^*$. Figure 4 shows the first 20 members of the series, drawn from the ‘topological coordinates’ [17], i.e., from scaled eigenvectors of their adjacency matrices, and assigns them to their maximal point group symmetries; affine realisations of any particular polyhedron may belong to subgroups of this maximal group.

3. Discussion

The smallest Dürer polyhedron is the $C_{3v}$ object obtained by truncation of a tetrahedron on three of its four vertices. Although it is a symmetrical, uniquely minimal solution without hidden vertices, it would be balanced on an edge when projected as in Melencolia I, and so can be rejected on physical grounds as a candidate for the polyhedron that Dürer had in mind when making the engraving, unless he was prepared to invoke hidden supports or ignore impending collapse.

The two distinct solutions for $v^* = 2$ are related by a Stone-Wales transformation [18] i.e., a 90° rotation of the bridge $B$. One is the bi-truncated cube ($D_{3d}$) which in an unsymmetrical realisation is the traditional explanation for the solid in the engraving; the other has only $C_4$ symmetry.

Clearly the class of Dürer polyhedra is infinite, as two distinct polyhedra at any even value of $v^*$ can be obtained from the $v^* = 2$ solutions by converting the central hidden edge to a ladder of $(v^* - 2)/2$ fused squares. Each solution for $v^* > 2$ outside this series also generates an infinite set of larger solutions by ladder extension. Some statistics on the initial growth of the series of Dürer polyhedra can be obtained by filtering the Dürer polyhedra from the lists of general cubic polyhedra available from the plantri program of Brinkmann and Mckay [19]. As Table 1 shows, the Dürer polyhedra show the expected combinatorial rise in absolute number as the number of vertices increases, but appear to form a decreasing fraction of the population of cubic polyhedra. Post hoc filtering is not an efficient way to generate subsets of the cubic polyhedra, but the program could be adapted to generate only Dürer polyhedra if it became important to have more complete lists.
Figure 3. Dürer polyhedra on $v = 10 + 2k$ vertices, $(k = 0, 1, 2, 3, 4)$, shown as bridges of the terminator circuit $C(D^0)$, each labelled with a multiplicity factor i.e., the number of non-isomorphic polyhedra derived by setting the given bridge in $C$. 
Figure 4. The first 20 Dürer polyhedra, drawn from eigenvectors of the adjacency matrix [17]. Polyhedra are labelled \( v : f_3 f_4 f_5 \ldots : S \), where \( v \) is the vertex count, \( f_r \) the number of faces of size \( r \) and \( S \) the maximal point group symmetry.
Table 1. Comparison of counts for Dürer and general cubic polyhedra. \( N(D) \) and \( N(c) \) are the respective numbers of Dürer and cubic polyhedra on \( v \) vertices, and \( \%D \) is their ratio, expressed as a percentage. The counts are obtained by filtering the lists of distinct cubic polyhedra produced by the plantri program [19].

<table>
<thead>
<tr>
<th>( v )</th>
<th>4</th>
<th>6</th>
<th>8</th>
<th>10</th>
<th>12</th>
<th>14</th>
<th>16</th>
<th>18</th>
<th>20</th>
<th>22</th>
<th>24</th>
<th>26</th>
</tr>
</thead>
<tbody>
<tr>
<td>( N(D) )</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>1</td>
<td>2</td>
<td>4</td>
<td>13</td>
<td>58</td>
<td>308</td>
<td>1826</td>
<td>11810</td>
<td>80036</td>
</tr>
<tr>
<td>( N(c) )</td>
<td>1</td>
<td>1</td>
<td>2</td>
<td>5</td>
<td>14</td>
<td>50</td>
<td>233</td>
<td>1249</td>
<td>7595</td>
<td>49566</td>
<td>339722</td>
<td>2406841</td>
</tr>
<tr>
<td>( %D )</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>20</td>
<td>14</td>
<td>8</td>
<td>6</td>
<td>5</td>
<td>4</td>
<td>3.7</td>
<td>3.5</td>
<td>3.3</td>
</tr>
</tbody>
</table>

All these polyhedra are combinatorially compatible with what can be seen in the engraving, though as the number of hidden vertices increases, the implied metric and physical constraints become more severe. The identity of the polyhedron in the engraving is a historical fact, recoverable or not, and external evidence seems to point to the second-smallest Dürer polyhedron, the bi-truncated cube, as the illustrated object. The purely internal evidence is open to wider interpretation. Without making the anachronistic claim that Dürer would have been able to find these polyhedra by the Euler theorem, we note that any one of the smaller members of the series could have been found empirically by truncating the cube.

Acknowledgement

The authors thank the British Council and DAAD for financial support, PWF thanks Kreiskrankenhaus Ilmenau for medical support, and Prof. H. Sachs for helpful discussion.

References


DÜRER POLYHEDRA: THE DARK SIDE OF MELANCHOLIA

[5] The work by Dürer cited in Ref. [3] as containing such a drawing is: Un-
terweisung der Messung mit dem Zirkel und Richtscheit in Linien Ebne
und Gantzen Corporen (Instruction in the Measurement with Compasses and
Straight-Edge of Lines, Planes and Solid Bodies), Nuremberg, 1525, which
is available as The Painter’s Manual, A. Dürer, trans. W.S. Strauss, Abaris


20–36.

Combinatorics. An International Congress, ICAC ’97, Hong Kong. ed., Kar-Ping

Years Old Puzzle Solved?, Netherlands Akad. Wetensch. Proc. (B) 84 (3)


(Sept. 1984).


Part 3 AB 12, (1922) 1–139.

(Berlin, 1934).


[16] The nauty program written by B.D. McKay, including the graph generators
makeg and makebg is available from http://www.anu.edu.au/people/bdm/
index.html

[17] D.E. Manolopoulos and P.W. Fowler, Molecular graphs, point groups and

[18] A.J. Stone and D.J. Wales, Theoretical studies of icosahedral C60 and some

[19] G. Brinkmann and B. McKay, Version 1.0 of the programme plantri.c and its
documentation are obtainable at http://cs.anu.edu.au/people/bdm.

Received 22 August 2000
Revised 14 May 2001