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COLORING SOME FINITE SETS IN \mathbb{R}^n

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Abstract

This note relates to bounds on the chromatic number $\chi(\mathbb{R}^n)$ of the Euclidean space, which is the minimum number of colors needed to color all the points in \mathbb{R}^n so that any two points at the distance 1 receive different colors. In [6] a sequence of graphs G_n in \mathbb{R}^n was introduced showing that $\chi(\mathbb{R}^n) \geq \chi(G_n) \geq (1 + o(1))\frac{n^2}{6}$. For many years, this bound has been remaining the best known bound for the chromatic numbers of some low-dimensional spaces. Here we prove that $\chi(G_n) \sim \frac{n^2}{6}$ and find an exact formula for the chromatic number in the case of $n = 2^k$ and $n = 2^k - 1$.

Keywords: chromatic number, independence number, distance graph.

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1. INTRODUCTION

In this note, we study the classical chromatic number $\chi(\mathbb{R}^n)$ of the Euclidean space. The quantity $\chi(\mathbb{R}^n)$ is the minimum number of colors needed to color all the points in \mathbb{R}^n so that any two points at a given distance a receive different colors. By a well-known compactness result of Erdős and de Bruijn (see [1]), the value of $\chi(\mathbb{R}^n)$ is equal to the chromatic number of a *finite* distance graph $G = (V, E)$, where $V \subset \mathbb{R}^n$ and $E = \{\{\mathbf{x}, \mathbf{y}\} : |\mathbf{x} - \mathbf{y}| = a\}$.

Now we know that

$$(1.239 \dots + o(1))^n \leq \chi(\mathbb{R}^n) \leq (3 + o(1))^n,$$

where the lower bound is due to the third author of this paper (see [8]) and the upper bound is due to Larman and Rogers (see [6]). Also, in [3] one can find an up-to-date table of estimates obtained for the dimensions $n \leq 12$.

It is worth noting that the linear bound $\chi(\mathbb{R}^n) \geq n + 2$ is quite simple, and the first superlinear bound was discovered by Larman, Rogers, Erdős, and Sós in [6]. They considered a family of graphs $G_n = (V_n, E_n)$ with

$$V_n = \{\mathbf{x} = (x_1, \dots, x_n) : x_i \in \{0, 1\}, x_1 + \dots + x_n = 3\},$$

$$E_n = \{\{\mathbf{x}, \mathbf{y}\} : |\mathbf{x} - \mathbf{y}| = 2\}.$$

In other words, the vertices of G_n are all the 3-subsets of the set $[n] = \{1, \dots, n\}$ and two vertices A, B are connected with an edge iff $|A \cap B| = 1$. Larman and Rogers [6] used an earlier result by Zs. Nagy who proved the following theorem.

Theorem 1 [6]. *Let s and $t \leq 3$ be nonnegative integers and let $n = 4s + t$. Then*

$$\alpha(G_n) = \begin{cases} n, & \text{if } t = 0, \\ n - 1, & \text{if } t = 1, \\ n - 2, & \text{if } t = 2 \text{ or } t = 3. \end{cases}$$

The standard inequality $\chi(G_n) \geq \frac{|V_n|}{\alpha(G_n)}$ combined with the above theorem gives an obvious corollary.

Corollary 2 [6]. *Let s and $t \leq 3$ be nonnegative integers and let $n = 4s + t$. Then*

$$\chi(G_n) \geq \begin{cases} \frac{(n-1)(n-2)}{6}, & \text{if } t = 0, \\ \frac{n(n-2)}{6}, & \text{if } t = 1, \\ \frac{n(n-1)}{6}, & \text{if } t = 2 \text{ or } t = 3. \end{cases}$$

The bounds from the corollary are applied to estimate from below the chromatic numbers $\chi(\mathbb{R}^{n-1})$, since the vertices of G_n lie in the hyperplane $x_1 + \dots + x_n = 3$. Now all these bounds are surpassed due to the consideration of some other distance graphs (see [3]). However, it could happen that actually $\chi(G_n)$ is much bigger than the ratio $\frac{|V_n|}{\alpha(G_n)}$. It turns out that this is not the case, and the main result of this note is as follows.

Theorem 3. *If $n = 2^k$ for some integer $k \geq 2$, then*

$$\chi(G_n) = \frac{(n-1)(n-2)}{6}.$$

Additionally, if $n = 2^k - 1$ for some integer $k \geq 2$, then

$$\chi(G_n) = \frac{n(n-1)}{6}.$$

Finally, there is a constant c such that for every n ,

$$\chi(G_n) \leq \frac{(n-1)(n-2)}{6} + cn.$$

Our proof yields that $c \leq 5.5$. With some more work we could prove that $c \leq 4.5$. On the other hand, since $n(n-1)/6 - (n-1)(n-2)/6 = (n-1)/3$, we have $c \geq 1/3$.

In the next section, we prove Theorem 3.

2. PROOF OF THEOREM 3

Easily, $\chi(G_3) = 1$, $\chi(G_4) = 1$, $\chi(G_5) = 3$.

Let $f(n) := \frac{(n-1)(n-2)}{6}$. We show by induction on k that $\chi(G_{2^k}) = f(2^k)$. For $k = 2$ it is trivial. Assume that for some k we established the induction hypothesis. Partition the set $[n] = [2^{k+1}]$ into the equal parts $A_1 = [\frac{n}{2}]$, $A_2 = [n] \setminus [\frac{n}{2}]$ of size 2^k . Denote by U_1 and U_2 the sets of vertices of $G = G_{2^{k+1}}$ lying in the sets

A_1 and A_2 respectively. By the induction assumption, each of the induced subgraphs $G[U_1]$ and $G[U_2]$ can be properly colored with at most $f(2^k)$ colors. Cover all pairs of the 2^k elements of A_1 with disjoint perfect matchings $N_1, \dots, N_{2^{k-1}}$ and all pairs of the 2^k elements of A_2 with matchings $M_1, \dots, M_{2^{k-1}}$. We form a color class $C(i, j)$ for $1 \leq i \leq 2^k - 1, 1 \leq j \leq 2^{k-1}$ as follows. Consider the matchings N_i, M_i and assume that the edges are $\{u_1, u_2\}, \{u_3, u_4\}, \dots$ in N_i and $\{v_1, v_2\}, \{v_3, v_4\}, \dots$ in M_i . For $j = 1, \dots, 2^{k-1}$ let $D(i, j)$ denote the following set of 4-tuples (indices are considered modulo 2^k):

$$\{u_1, u_2, v_{2j-1}, v_{2j}\}, \{u_3, u_4, v_{2j+1}, v_{2j+2}\}, \dots, \{u_{2^k-1}, u_{2^k}, v_{2j-3}, v_{2j-2}\}.$$

For $i = 1, \dots, 2^k - 1$ and $j = 1, \dots, 2^{k-1}$, the color class $C(i, j)$ is formed by the collection of triples contained in the members of $D(i, j)$. The intersection sizes are all 0 or 2, so the triples in $C(i, j)$ form an independent set in G . Moreover, each triple is contained in a member of some $D(i, j)$. The total number of used colors is

$$2^{k-1}(2^k - 1) + f(2^k) = 2^{2k-1} - 2^{k-1} + \frac{(2^k-1)(2^k-2)}{6} = f(2^{k+1}).$$

This proves the first statement of the theorem. Since $\chi(G_n) \leq \chi(G_{n+1})$, this together with Corollary 2 also implies the statement of the theorem for $n = 2^k - 1$.

It remains to show that there exists a constant c such that $\chi(G_n) \leq \frac{n^2}{6} + cn$ for every n . Consider our coloring in steps.

Step 1: Let $n = 4s_1 + t_1$ where $t_1 \leq 3$. First, color all triples containing the elements $4s_1 + 1, \dots, 4s_1 + t_1$ with at most $t_1(n - 1) < 3n$ colors. Now consider the set $[4s_1]$ and all the triples in this set. Partition $[4s_1]$ into $A_1 = [2s_1]$ and $A_2 = [4s_1] - [2s_1]$ and color the triples intersecting both A_1 and A_2 with $s_1(2s_1 - 1) < \frac{n}{4}(\frac{n}{2} - 1)$ colors as above.

Step 2: Since the triples contained in A_1 are disjoint from the triples contained in A_2 , we will use for coloring the triples contained in A_2 the same colors and the same procedure as for the triples contained in A_1 . Consider A_1 . Let $n_1 = |A_1| = 2s_1 = 4s_2 + t_2$ where $t_2 \leq 3$. Since $2s_1$ is even, $t_2 \leq 2$. By construction, $n_1 \leq \frac{n}{2}$. Similarly to Step 1, color all triples containing the elements $4s_2 + 1, \dots, 4s_2 + t_2$ with at most $t_2(n_1 - 1) < 2n_1$ colors. Partition $[4s_2]$ into $A_{1,1} = [2s_2]$ and $A_{1,2} = [4s_2] - [2s_2]$ and color the triples intersecting both $A_{1,1}$ and $A_{1,2}$ with at most $\frac{n}{8}(\frac{n}{4} - 1)$ new colors.

Step i (for $i \geq 3$): If $2s_{i-1} \leq 2$, then Stop. Otherwise, repeat Step 2 with $[2s_{i-1}]$ in place of $[2s_1]$.

Altogether, we use at most $(3n + \frac{n}{4}(\frac{n}{2} - 1)) + (\frac{2n}{2} + \frac{n}{8}(\frac{n}{4} - 1)) + (\frac{2n}{4} + \frac{n}{16}(\frac{n}{8} - 1)) + \dots < 5n + \frac{n^2}{8} \cdot \frac{4}{3} = \frac{n^2}{6} + 5n = \frac{(n-1)(n-2)}{6} + 5.5n - 1/3$ colors. The theorem is proved. \blacksquare

3. DISCUSSION

As we have already said, the constant 5 in the bound $\chi(G_n) \leq \frac{n^2}{6} + 5n$ is not the best possible and can be improved. However, to find the exact value of the chromatic number is still interesting. For example, we know that $\chi(\mathbb{R}^{12}) \geq 27$ (see [3]). At the same time, $\chi(G_{13}) \geq \left\lceil \frac{\binom{13}{3}}{12} \right\rceil = 24$ (due to Corollary 2), and the proof of Theorem 3 applied for $n = 13$ yields a bound $\chi(G_{13}) \leq 31$.

It would be quite interesting to study more general graphs. Let $G(n, r, s)$ be the graph whose set of vertices consists of all the r -subsets of the set $[n]$ and whose set of edges is formed by all possible pairs of vertices A, B with $|A \cap B| = s$. Larman proved in [5] that

$$\chi(\mathbb{R}^n) \geq \chi(G(n, 5, 2)) \geq \frac{\binom{n}{5}}{\alpha(G(n, 5, 2))} \geq (1 + o(1)) \frac{\binom{n}{5}}{1485n^2} \sim \frac{n^3}{178200}.$$

Thus, the main result of Larman was in finding the bound $\alpha(G(n, 5, 2)) \leq (1 + o(1))1485n^2$. However, the so-called linear algebra method ([2], see also [8]) can be directly applied here to show that $\alpha(G(n, 5, 2)) \leq (1 + o(1))\binom{n}{2} \sim \frac{n^2}{2}$. This substantially improves Larman's estimate and gives $\chi(G(n, 5, 2)) \geq (1 + o(1))\frac{n^3}{60}$. We do not know any further improvements on this result. On the other hand, observe that for any 3-set A , the collection of 5-sets containing A forms an independent set in $G(n, 5, 2)$, yielding $\chi(G(n, 5, 2)) \leq \binom{n}{3} \sim \frac{n^3}{6}$. It is plausible that $\chi(G(n, 5, 2)) \sim cn^3$ with a constant $c \in [1/60, 1/6]$, but this constant is not yet found and even no better bounds for c have been published.

Furthermore, the graphs $G(n, 5, 3)$ have been studied, since the best known lower bound $\chi(\mathbb{R}^9) \geq 21$ is due to the fact that $\chi(G(10, 5, 3)) = 21$ (see [4]). No related results concerning the case of $n \rightarrow \infty$ have apparently been published.

Although for combinatorial geometry small values of n are of greater interest, we see that the consideration of graphs $G(n, r, s)$ with small r, s and *growing* n is of its intrinsic interest, too. So assume that r, s are fixed and $n \rightarrow \infty$. We have

$$\chi(G(n, r, s)) \leq \min\{O(n^{r-s}), O(n^{s+1})\}.$$

The first bound follows from Brooks' theorem, since the maximum degree of $G(n, r, s)$ is

$$\binom{r}{s} \binom{n-r}{r-s} = (1 + o(1)) \frac{r!}{s!(r-s)!(r-s)!} n^{r-s}.$$

The second bound is a simple generalization of the above-mentioned bound $\chi(G(n, 5, 2)) \leq (1 + o(1))n^3/6$.

Note that the second bound can be somewhat improved. Assume $s < r/2$, so $q := \lceil (r-1)/s \rceil$ is at least 2. Assuming that q divides n , partition $[n]$ into q

equal classes, A_1, \dots, A_q . Let \mathcal{C} be the family of $(s+1)$ -sets that are subsets of some A_i . For each $B \in \mathcal{C}$, the r -sets containing B form an independent set in $G(n, r, s)$, and by the pigeonhole principle every r -set contains such B , hence

$$\chi(G(n, r, s)) \leq |\mathcal{C}| = q \binom{n/q}{s+1} = (1 + o(1)) \frac{n^{s+1}}{q^s (s+1)!}.$$

In particular, $\chi(G(n, 5, 2)) \leq (1 + o(1)) \frac{n^3}{24}$, which improves the previous bound $\frac{n^3}{6}$.

It is worthwhile to look at the construction in Section 2 from a different point of view. For $n = 2^k$ we constructed a 4-uniform hypergraph \mathcal{H} with the property that every 3-subset of vertices is covered exactly once. Note that $e(\mathcal{H}) = \binom{n}{3}/4$. Then we decomposed $E(\mathcal{H})$ into $\binom{n}{3}$ perfect matchings. Each matching gives a color class of our coloring. Note that instead of providing the explicit decomposition, we could have used a classical theorem of Pippenger and Spencer [7], which claims the existence of $(1 + o(1)) \binom{n}{3}$ covering matchings.

This motivates the following possible approach to the case $r = 2s + 1$. The discussion here is not a proof, it is just a sketch of a possible way to generalize our argument. Assume that we managed to construct an $(r+s)$ -uniform hypergraph \mathcal{H} that covers every r -set exactly once. Then $e(\mathcal{H}) = \binom{n}{r} / \binom{n}{r+s}$. Assume that \mathcal{H} can be decomposed into t hypergraphs, $\mathcal{N}_1, \dots, \mathcal{N}_t$, such that for every i and every $A, B \in \mathcal{N}_i$ we have $|A \cap B| \leq s - 1$. Then the r -sets covered by sets in \mathcal{N}_i form an independent set, yielding $\chi(G(n, r, s)) \leq t$. Probably a generalization of the theorem of Pippenger and Spencer [7] would give $t \leq (1 + o(1)) \binom{n}{r} / \binom{n}{s} = (1 + o(1)) (s!/r!) n^{r-s}$. This bound, if true, would be asymptotically best possible, since the already mentioned linear algebra method (see [2, 8]) ensures that $\alpha(G(n, 2s+1, s)) \leq (1 + o(1)) \binom{n}{s}$ and so $\chi(G(n, 2s+1, s)) \geq (1 + o(1)) \binom{n}{r} / \binom{n}{s}$, provided $s+1$ is a prime power. In particular, we would get $\chi(G(n, 5, 2)) \sim \frac{n^2}{60}$.

The case of simultaneously growing n, r, s has also been studied. Namely, $r \sim r'n$ and $s \sim s'n$ with any $r' \in (0, 1)$ and $s' \in (0, r')$ have been considered. This is due to the fact that the first exponential estimate to the quantity $\chi(\mathbb{R}^n)$, $\chi(\mathbb{R}^n) \geq (1.207 \dots + o(1))^n$, was obtained by Frankl and Wilson in [2] with the help of some graphs $G(n, r, s)$ having $r \sim r'n$ and $s \sim \frac{r'}{2}n$. Lower bounds are usually based on the linear algebra (see [8]) and upper bounds can be found in [9].

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