

## A NOTE ON FORCING 3-REPETITIONS IN DEGREE SEQUENCES

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### Abstract

In Caro, Shapira and Yuster [*Forcing  $k$ -repetitions in degree sequences*, Electron. J. Combin. 21 (2014) #P1.24] it is proven that for any graph  $G$  with at least 5 vertices, one can delete at most 6 vertices such that the subgraph obtained has at least three vertices with the same degree. Furthermore they show that for certain graphs one needs to remove at least 3 vertices in order that the resulting graph has at least 3 vertices of the same degree.

In this note we prove that for any graph  $G$  with at least 5 vertices, one can delete at most 5 vertices such that the subgraph obtained has at least three vertices with the same degree. We also show that for any triangle-free graph  $G$  with at least 6 vertices, one can delete at most one vertex such that the subgraph obtained has at least three vertices with the same degree and this result is tight for triangle-free graphs.

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### 1. INTRODUCTION

All the graphs in this paper are simple, that is they have no loops and no multiple edges. A well-known elementary exercise states that every simple graph has two vertices with the same degree. Since there are graphs without 3 vertices of the same degree, it is natural to ask how many vertices one needs to remove from a graph in order that the resulting graph has at least 3 vertices of the same degree. The following theorem was proven in [1].

**Theorem 1.1.** *For any graph  $G$  with at least 5 vertices, one can delete at most 6 vertices such that the subgraph obtained has at least three vertices with the same degree.*

Furthermore it is shown in [1] that for certain graphs one needs to remove at least 3 vertices in order that the resulting graph has at least 3 vertices of the same degree.

More generally the following was shown in [1]. For a given positive integer  $k$ , let  $C = C(k)$  denote the least integer such that any graph with  $n$  vertices has an induced subgraph with at least  $n - C$  vertices, such that at least  $\min\{k, n - C\}$  vertices in this subgraph are of the same degree, then  $\Omega(k \log k) \leq C(k) \leq (8k)^k$ . For  $k = 3$  the bounds in [1] are  $3 \leq C(3) \leq 6$  and the exact value was left as an open question in that paper. In this note we slightly improve Theorem 1.1, by showing that  $C(3) \leq 5$ . Formally we prove the following.

**Theorem 1.2.** *For any graph  $G$  with at least 5 vertices, one can delete at most 5 vertices such that the subgraph obtained has at least three vertices with the same degree.*

We also prove the following result for triangle-free graphs.

**Theorem 1.3.** *For any triangle-free graph  $G$  with at least 6 vertices, one can delete at most one vertex such that the subgraph obtained has at least three vertices with the same degree.*

Theorem 1.3 is tight as a path on 4 vertices together with an isolated vertex is a graph on 5 vertices with no three vertices of the same degree, for which no removal of a vertex creates three vertices of the same degree. Furthermore for any  $n \geq 1$  the bipartite half-graph (defined below) contains  $2n$  vertices and no three of them of the same degree.

**Definition 1.1.** The half-graph on  $2n$  vertices is a bipartite graph with vertex set  $\{x_1, \dots, x_n\} \cup \{y_1, \dots, y_n\}$  such that  $x_i$  is adjacent to  $y_j$  if and only if  $i + j > n$ . Note that  $x_i$  and  $y_i$  have degree  $i$ , for  $1 \leq i \leq n$ .

Note that for an odd number of vertices we can take a half-graph and one isolated vertex. The resulting graph has no three vertices of the same degree.

## 2. PROOF OF THEOREM 1.2

Let  $G(V, E)$  be a graph on  $n \geq 5$  vertices. Denote by  $d(v)$  the degree of vertex  $v$  in graph  $G$ . Let  $H \subset V$  be a set of 3 vertices  $x, y, z$  where  $d(x) \leq d(y) \leq d(z)$ . Set  $H$  is called *balanceable* if one of the following conditions hold.

1.  $G[H]$  is an independent set.
2.  $G[H]$  is a clique.
3.  $G[H]$  contains only the edge  $(x, y)$ .
4.  $G[H]$  contains only the edges  $(x, y)$  and  $(x, z)$ .

Let  $p = d(z) - d(y)$  and  $q = d(y) - d(x)$ . We will need the following lemma whose proof is almost identical to the proof of Theorem 1.1 and is given only for completeness.

**Lemma 2.1.** *If graph  $G$  contains a balanceable set  $H$ , then one can delete at most  $p + q + \max(p, q)$  vertices from  $G$  such that the subgraph obtained has at least three vertices with the same degree.*

**Proof.** Let  $H$  be a balanceable set in graph  $G$  with vertices  $x, y, z$  such that  $d(x) \leq d(y) \leq d(z)$  and recall that  $p = d(z) - d(y)$  and  $q = d(y) - d(x)$ . Throughout the proof we denote by  $N(\cdot)$  the set of neighbors of a vertex in the current  $G$  (that is, in the graph  $G$  after some vertices have possibly been deleted). Similarly, we denote by  $d(\cdot)$  the degree of a vertex in the current  $G$ . We can assume without loss of generality that set  $H$  satisfies conditions 1 or 3 of a balanceable set (otherwise we just take the complement of graph  $G$  as our graph). We will consider two cases.

*Case 1.*  $d(x) < d(y) = d(z)$ . In this case  $p = 0$ . If  $(N(z) \setminus N(x)) \cap (N(y) \setminus N(x)) \neq \emptyset$ , we can delete a vertex of this intersection and decrease the degrees of  $y$  and  $z$  by 1 without affecting the degree of  $x$  (notice that  $x \notin (N(z) \setminus N(x)) \cap (N(y) \setminus N(x))$  as  $H$  satisfies conditions 1 or 3 of a balanceable set). Otherwise, if  $(N(z) \setminus N(x)) \cap (N(y) \setminus N(x)) = \emptyset$  we can delete a vertex of  $N(z) \setminus N(x)$  and a vertex of  $(N(y) \setminus N(x)) \setminus \{x\}$  and decrease the degrees of  $y$  and  $z$  by 1 without affecting the degree of  $x$ . Observe that in any case we delete at most two vertices. Repeating this process at most  $q$  times we eventually obtain  $d(x) = d(y) = d(z)$ . Overall we have deleted at most  $2q \leq p + q + \max(p, q)$  vertices.

*Case 2.*  $d(x) \leq d(y) < d(z)$ . Let us first equate  $d(z)$  and  $d(y)$  by deleting some  $u \in N(z) \setminus N(y)$ . Observe that  $u \neq x$ . We always prefer to delete a vertex  $u$  that is non-adjacent to  $x$ , as long as there is such a vertex  $u$ . Overall, we have deleted  $p$  vertices. The problem is that in the current graph we may have that  $d(x)$  also decreased by some amount  $r \leq p$ . Suppose first that  $r = 0$ . As in the previous case, we may need to delete  $2q$  additional vertices to equate the degrees of  $y$  and  $z$  to that of  $x$  so in total we removed  $p + 2q \leq p + q + \max(p, q)$  vertices. If  $r > 0$ , then this means that at some point, when we deleted a vertex  $u$ , that vertex also had to be adjacent to  $x$ . Hence, in the current graph  $(N(z) \setminus N(x)) \subseteq (N(y) \setminus N(x))$ . So, we may simply delete  $r + q$  additional vertices

(all of them from  $N(z) \setminus N(x)$ ) to equate the degrees of  $y$  and  $z$  to that of  $x$ . Overall, we deleted  $p + r + q \leq 2p + q \leq p + q + \max(p, q)$  vertices and the proof follows. ■

We call a balanceable set  $H$  good if  $p + q \leq 3$  and  $\max(p, q) \leq 2$ . Hence by Lemma 2.1 it is sufficient to find a good set in  $G$  in order to finish our proof, as in this case we have  $p + q + \max(p, q) \leq 5$ .

We will need the following lemma proven in [2].

**Lemma 2.2.** *Let  $G$  be any graph of order  $n$  and let  $m$  be any positive integer less than  $n$ . Then  $G$  contains a set  $X$  of  $m+1$  vertices such that  $|d(u) - d(v)| \leq m-1$  for all  $u, v \in X$*

Hence by Lemma 2.2 graph  $G$  has a set  $X$  of 5 vertices such that  $|d(u) - d(v)| \leq 3$  for all  $u, v \in X$ . Assume that the vertices of  $X$  are  $v_1, v_2, v_3, v_4, v_5$ . Furthermore assume that  $d(v_1) \leq d(v_2) \leq \dots \leq d(v_5)$ . We also may assume that  $G[X]$  contains at least two of the edges in the set  $\{(v_1, v_2), (v_2, v_3), (v_3, v_4), (v_4, v_5)\}$ , for otherwise this would be true in the complement of the graph  $G$ . Now let  $d(v_1) = d$ . By the definition of  $X$  we know that  $d(v_5) \leq d + 3$ . We also may assume the following.

**Lemma 2.3.**  $d + 1 \leq d(v_3) \leq d + 2$ .

**Proof.** If  $d(v_3) = d$ , then  $d(v_1) = d(v_2) = d(v_3) = d$  and we are done. Furthermore if  $d(v_3) = d + 3$ , then  $d(v_3) = d(v_4) = d(v_5) = d + 3$  and we are done. ■

**Corollary 1.** *If  $H \subset X$  is a balanceable set such that  $v_3 \in H$ , then  $H$  is a good set.*

**Proof.** In this case  $\max(p, q) \leq 2$ . Furthermore  $p + q \leq 3$  by the definition of  $X$ . ■

Finally we will need the following simple lemma.

**Lemma 2.4.** *For any  $2 \leq i \leq 4$ , if  $(v_i, v_{i+1}) \notin E$  and  $(v_{i-1}, v_i) \in E$ , then the set  $\{v_{i-1}, v_i, v_{i+1}\}$  is a good set.*

**Proof.** Each such set for  $2 \leq i \leq 4$  is a balanceable set by definition, furthermore each such set contains the vertex  $v_3$  and hence is good by Corollary 1. ■

Let  $E' = \{(v_1, v_2), (v_2, v_3), (v_3, v_4), (v_4, v_5)\}$ . Recall that we may assume that  $G[X]$  contains at least two of the edges of  $E'$ . We will consider two cases.

*Case 1.*  $G[X]$  contains at least three of the edges of  $E'$ . By Lemma 2.4 we may assume that  $G[X]$  contains the edges  $(v_2, v_3), (v_3, v_4), (v_4, v_5)$ . Now if  $(v_3, v_5) \in E$ , then  $\{v_3, v_4, v_5\}$  is a good set, while if  $(v_3, v_5) \notin E$ , then  $\{v_2, v_3, v_5\}$  is a good set.

*Case 2.*  $G[X]$  contains exactly two edges of  $E'$ . By Lemma 2.4 we may assume that  $G[X]$  contains the edges  $(v_3, v_4), (v_4, v_5)$  but not the edges  $(v_1, v_2), (v_2, v_3)$ . Now if  $(v_3, v_5) \in E$ , then  $\{v_3, v_4, v_5\}$  is a good set. Hence we may assume that  $(v_3, v_5) \notin E$ . Furthermore if  $(v_1, v_3) \notin E$ , then  $\{v_1, v_2, v_3\}$  is a good set. Hence we may assume that  $(v_1, v_3) \in E$ . Now since  $(v_1, v_3) \in E$  and  $(v_3, v_5) \notin E$  we conclude that  $\{v_1, v_3, v_5\}$  is a good set.

This finishes the proof.

### 3. PROOF OF THEOREM 1.3

We will need the following theorem proven in [3].

**Theorem 3.1.** *Let  $G$  be a triangle-free graphs without three vertices of the same degree and without isolated vertices. Let the vertices of  $G$  be  $v_1, \dots, v_n$  such that  $d(v_1) \leq d(v_2) \leq \dots \leq d(v_n)$ . Then  $d(v_1) = 1$  and for all  $1 \leq i \leq n - 2$  we have  $d(v_{i+2}) - d(v_i) = 1$ .*

We start by proving the following.

**Theorem 3.2.** *For any triangle-free graph  $G$  with at least 5 vertices and no isolated vertices, one can delete at most one vertex such that the subgraph obtained has at least three vertices with the same degree.*

**Proof.** Suppose that  $G$  contains  $n \geq 5$  vertices and the vertices of  $G$  are  $v_1, \dots, v_n$  such that  $d(v_1) \leq d(v_2) \leq \dots \leq d(v_n)$ . By Theorem 3.1 we may consider two cases.

*Case 1.*  $d(v_1) = 1$  and  $d(v_2) = 1$  and for all  $1 \leq i \leq n - 2$  we have  $d(v_{i+2}) - d(v_i) = 1$ . In this case the degree sequence of  $G$  starts with 1, 1, 2, 2, 3.

If  $v_1$  is not adjacent to  $v_i$  for some  $2 \leq i \leq 4$ , we can remove  $v_1$  and obtain a graph with three vertices of the same degree. If  $v_2$  is not adjacent to  $v_i$  for some  $3 \leq i \leq 4$  or  $i = 1$ , we can remove  $v_2$  and obtain a graph with three vertices of the same degree. Hence we need to only consider the case where  $v_n$  is not adjacent to both  $v_1$  and  $v_2$ . Thus the neighbor  $v_j$  of  $v_n$  with the smallest index  $j$  satisfies  $j \geq 3$ , so we can remove  $v_n$  and obtain a graph with three vertices of the same degree.

*Case 2.*  $d(v_1) = 1$  and  $d(v_2) = 2$  and for all  $1 \leq i \leq n - 2$  we have  $d(v_{i+2}) - d(v_i) = 1$ . As the sum of degrees in a graph is even we may assume that  $n \geq 8$  in this case. Hence the degree sequence of  $G$  starts with 1, 2, 2, 3, 3, 4, 4, 5.

If  $v_1$  is not adjacent to  $v_i$  for some  $2 \leq i \leq 3$ , we can remove  $v_1$  and obtain a graph with three vertices of the same degree. Assume without loss of generality that  $v_1$  is adjacent to  $v_2$  (as  $d(v_3) = d(v_2) = 2$ ). Removing  $v_1$  we obtain a graph

$G'$  in which the degree sequence starts with 1, 2, 3 and this graph must contain three vertices of the same degree by Theorem 3.1. ■

Theorem 1.3 follows as if graph  $G$  has two isolated vertices we can remove a neighbor of a vertex of degree 1 to obtain three vertices of degree 0.

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