

## ON THE OPTIMALITY OF 3-RESTRICTED ARC CONNECTIVITY FOR DIGRAPHS AND BIPARTITE DIGRAPHS

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### Abstract

Let  $D$  be a strong digraph. An arc subset  $S$  is a  $k$ -restricted arc cut of  $D$  if  $D - S$  has a strong component  $D'$  with order at least  $k$  such that  $D \setminus V(D')$  contains a connected subdigraph with order at least  $k$ . If such a  $k$ -restricted arc cut exists in  $D$ , then  $D$  is called  $\lambda^k$ -connected. For a  $\lambda^k$ -connected digraph  $D$ , the  $k$ -restricted arc connectivity, denoted by  $\lambda^k(D)$ , is the minimum cardinality over all  $k$ -restricted arc cuts of  $D$ . It is known that for many digraphs  $\lambda^k(D) \leq \xi^k(D)$ , where  $\xi^k(D)$  denotes the minimum  $k$ -degree of  $D$ .  $D$  is called  $\lambda^k$ -optimal if  $\lambda^k(D) = \xi^k(D)$ . In this paper, we will give some sufficient conditions for digraphs and bipartite digraphs to be  $\lambda^3$ -optimal.

**Keywords:** restricted arc-connectivity, bipartite digraph, optimality, digraph, network.

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### 1. INTRODUCTION

It is well-known that the network can be modelled as a digraph  $D$  with vertices  $V(D)$  representing sites and arcs  $A(D)$  representing links between sites of the network. Let  $v \in V(D)$ , the *out-neighborhood* of  $v$  is the set  $N^+(v) = \{x \in V(D) : vx \in A(D)\}$  and the *out-degree* of  $v$  is  $d^+(v) = |N^+(v)|$ . The *in-neighborhood* of  $v$  is the set  $N^-(v) = \{x \in V(D) : xv \in A(D)\}$  and the *in-degree* of  $v$  is  $d^-(v) = |N^-(v)|$ . The *neighborhood* of  $v$  is  $N(v) = N^+(v) \cup N^-(v)$ .

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Let  $\delta^+(D)$ ,  $\delta^-(D)$  and  $\delta(D)$  denote, respectively, the *minimum out-degree*, the *minimum in-degree* and the *minimum degree* of  $D$ .

For a pair nonempty vertex sets  $X$  and  $Y$  of  $D$ ,  $[X, Y] = \{xy \in A(D) : x \in X, y \in Y\}$ . Specially, if  $Y = \bar{X}$ , where  $\bar{X} = V(D) \setminus X$ , then we write  $\partial^+(X)$  or  $\partial^-(Y)$  instead of  $[X, Y]$ . For  $X \subseteq V(D)$ , the *subdigraph* of  $D$  induced by  $X$  is denoted by  $D[X]$ . The *underlying graph*  $U(D)$  of  $D$  is the unique graph obtained from  $D$  by deleting the orientation of all arcs and keeping one edge of a pair of multiple edges.  $D$  is *connected* if  $U(D)$  is connected and  $D$  is *strongly connected* (or, just, *strong*) if there exists a directed  $(x, y)$ -path and a directed  $(y, x)$ -path for any  $x, y \in V(D)$ . We define a digraph with one vertex to be strong. A *connected (strong) component* of  $D$  is a maximal induced subdigraph of  $D$  which is connected (strong). If  $D$  has  $p$  strong components, then these strong components can be labeled  $D_1, \dots, D_p$  such that there is no arc from  $D_j$  to  $D_i$  unless  $j < i$ . We call such an ordering an *acyclic ordering of the strong components* of  $D$ .

In a strong digraph  $D$ , we often use *arc connectivity* of  $D$  to measure the reliability. An arc set  $S$  is a *arc cut* of  $D$  if  $D - S$  is not strong. The arc connectivity  $\lambda(D)$  is the minimum cardinality over all arc cuts of  $D$ . The arc cut  $S$  of  $D$  with cardinality  $\lambda(D)$  is called a  $\lambda$ -*cut*. Whitney's inequality shows  $\lambda(D) \leq \delta(D)$ . A strong digraph  $D$  with  $\lambda(D) = \delta(D)$  is called  $\lambda$ -*optimal*. However, only using arc connectivity to measure the reliability is not enough. In [12], Volkmann introduced the concept of *restricted arc connectivity*. An arc subset  $S$  of  $D$  is a *restricted arc cut* if  $D - S$  has a strong component  $D'$  with order at least 2 such that  $D \setminus V(D')$  contains an arc. If such an arc cut exists in  $D$ , then  $D$  is called  $\lambda'$ -*connected*. For a  $\lambda'$ -connected digraph  $D$ , the restricted arc connectivity, denoted by  $\lambda'(D)$ , is the minimum cardinality over all restricted arc cuts of  $D$ . The restricted arc cut  $S$  of  $D$  with cardinality  $\lambda'(D)$  is called a  $\lambda'$ -*cut*. In [13], Wang and Lin introduced the notion of *minimum arc degree*. Let  $xy \in A(D)$ . Then

$$\Omega(\{x, y\}) = \{\partial^+(\{x, y\}), \partial^-(\{x, y\}), \partial^+(\{x\}) \cup \partial^-(\{y\}), \partial^+(\{y\}) \cup \partial^-(\{x\})\}.$$

The *arc degree* of  $xy$  is  $\xi'(xy) = \min\{|S| : S \in \Omega(\{x, y\})\}$  and the minimum arc degree of  $D$  is  $\xi'(D) = \min\{\xi'(xy) : xy \in A(D)\}$ .

It was proved in [3, 13] that for many  $\lambda'$ -connected digraphs,  $\xi'(D)$  is an upper bound of  $\lambda'(D)$ . In [13], Wang and Lin introduced the concept of  $\lambda'$ -optimality. A  $\lambda'$ -connected digraph  $D$  with  $\xi'(D) = \lambda'(D)$  is called  $\lambda'$ -*optimal*. As a generalization of restricted arc connectivity, in [10], Lin *et al.* introduced the concept of *k-restricted arc connectivity*.

**Definition** [10]. Let  $D$  be a strong digraph. An arc subset  $S$  is a *k-restricted arc cut* of  $D$  if  $D - S$  has a strong component  $D'$  with order at least  $k$  such that

$D \setminus V(D')$  contains a connected subdigraph with order at least  $k$ . If such a  $k$ -restricted arc cut exists in  $D$ , then  $D$  is called  $\lambda^k$ -connected. For a  $\lambda^k$ -connected digraph  $D$ , the  $k$ -restricted arc connectivity, denoted by  $\lambda^k(D)$ , is the minimum cardinality over all  $k$ -restricted arc cuts of  $D$ . The  $k$ -restricted arc cut  $S$  of  $D$  with cardinality  $\lambda^k(D)$  is called a  $\lambda^k$ -cut.

**Definition** [10]. Let  $D$  be a strong digraph. For any  $X \subseteq V(D)$ , let  $\Omega(X) = \{\partial^+(X_1) \cup \partial^-(X \setminus X_1) : X_1 \subseteq X\}$  and  $\xi(X) = \min\{|S| : S \in \Omega(X)\}$ . Define the *minimum  $k$ -degree* of  $D$  to be

$$\xi^k(D) = \min\{\xi(X) : X \subseteq V(D), |X| = k, D[X] \text{ is connected}\}.$$

Clearly,  $\lambda^1(D) = \lambda(D)$ ,  $\lambda^2(D) = \lambda'(D)$ ,  $\xi^1(D) = \delta(D)$  and  $\xi^2(D) = \xi'(D)$ . Let  $D$  be a  $\lambda^k$ -connected digraph, where  $k \geq 2$ . Then  $D$  is  $\lambda^{k-1}$ -connected and  $\lambda^{k-1}(D) \leq \lambda^k(D)$ . It was shown in [10] that  $\xi^k(D)$  is an upper bound of  $\lambda^k(D)$  for many digraphs. And a  $\lambda^k$ -connected digraph  $D$  with  $\lambda^k(D) = \xi^k(D)$  is called  $\lambda^k$ -optimal.

The research on the  $\lambda^k$ -optimality of digraph  $D$  is considered to be a hot issue. In [11], Hellwig and Volkmann concluded many sufficient conditions for digraphs to be  $\lambda$ -optimal. Besides, sufficient conditions for digraphs to be  $\lambda'$ -optimal were also given by several authors, for example by Balbuena *et al.* [1–4], Chen *et al.* [5, 6], Grüter and Guo [7, 8], Liu and Zhang [9], Volkmann [12] and Wang and Lin [13]. However, closely related conditions for  $\lambda^3$ -optimal digraphs have received little attention until recently. In [10], Lin *et al.* gave some sufficient conditions for digraphs to be  $\lambda^3$ -optimal. In this paper, we will give some sufficient conditions for digraphs to be  $\lambda^3$ -optimal. As corollaries, degree conditions or degree sum conditions for a digraph or a bipartite digraph to be  $\lambda^3$ -optimal are given. The main contributions in this paper are as following.

**Theorem 1.** *Let  $D$  be a digraph with  $|V(D)| \geq 6$ . If  $|N^+(u) \cap N^-(v)| \geq 5$  for any  $u, v \in V(D)$  with  $uv \notin A(D)$ , then  $D$  is  $\lambda^3$ -optimal.*

**Theorem 2.** *Let  $D = (X, Y, A(D))$  be a bipartite digraph with  $|V(D)| \geq 6$ . If  $|N^+(u) \cap N^-(v)| \geq \left\lceil \frac{|V(D)|}{4} \right\rceil + 1$  for any  $u, v \in V(D)$  in the same partite, then  $D$  is  $\lambda^3$ -optimal.*

## 2. PROOF OF THEOREM 1

We first introduce three useful lemmas.

**Lemma 3** (Theorem 1.4 in [10]). *Let  $D$  be a strong digraph with  $\delta^+(D) \geq 2k - 1$  or  $\delta^-(D) \geq 2k - 1$ . Then  $D$  is  $\lambda^k$ -connected and  $\lambda^k(D) \leq \xi^k(D)$ .*

**Lemma 4.** *Let  $D$  be a strong digraph with  $\delta^+(D) \geq 2k - 1$  or  $\delta^-(D) \geq 2k - 1$ , and let  $S = \partial^+(X)$  be a  $\lambda^k$ -cut of  $D$ , where  $X$  is a subset of  $V(D)$ . If  $D[X]$  contains a connected subdigraph  $B$  with order  $k$  such that  $|N^+(x) \cap \bar{X}| \geq k$  for any  $x \in X \setminus V(B)$  or  $D[\bar{X}]$  contains a connected subdigraph  $C$  with order  $k$  such that  $|N^-(y) \cap X| \geq k$  for any  $y \in \bar{X} \setminus V(C)$ , then  $D$  is  $\lambda^k$ -optimal.*

**Proof.** By Lemma 3,  $D$  is  $\lambda^k$ -connected and  $\lambda^k(D) \leq \xi^k(D)$ . By symmetry, we only prove the case that  $D[X]$  contains a connected subdigraph  $B$  with order  $k$  such that  $|N^+(x) \cap \bar{X}| \geq k$  for any  $x \in X \setminus V(B)$ . The hypotheses imply that

$$\begin{aligned} \xi^k(D) &\leq |\partial^+(V(B))| = |[V(B), X \setminus V(B)]| + |[V(B), \bar{X}]| \\ &\leq k|X \setminus V(B)| + |[V(B), \bar{X}]| \leq \sum_{x \in X \setminus V(B)} |N^+(x) \cap \bar{X}| + |[V(B), \bar{X}]| \\ &= |[X \setminus V(B), \bar{X}]| + |[V(B), \bar{X}]| = |[X, \bar{X}]| = |S| = \lambda^k(D). \end{aligned}$$

Thus  $\lambda^k(D) = \xi^k(D)$  and  $D$  is  $\lambda^k$ -optimal. ■

**Lemma 5** (Lemma 4.1 in [10]). *Let  $D$  be a strong digraph with  $|V(D)| \geq 6$  and  $\delta(D) \geq 4$ , and let  $S$  be a  $\lambda^3$ -cut of  $D$ . If  $D$  is not  $\lambda^3$ -optimal, then there exists a subset of vertices  $X \subset V(D)$  such that  $S = \partial^+(X)$  and both induced subdigraphs  $D[X]$  and  $D[\bar{X}]$  contain a connected subdigraph with order 3.*

**Proof of Theorem 1.** Clearly,  $D$  is a strong digraph with  $\delta(D) \geq 5$ . By Lemma 3,  $D$  is  $\lambda^3$ -connected and  $\lambda^3(D) \leq \xi^3(D)$ . Suppose, on the contrary, that  $D$  is not  $\lambda^3$ -optimal, that is,  $\lambda^3(D) < \xi^3(D)$ . Let  $S$  be a  $\lambda^3$ -cut of  $D$ . By Lemma 5, there exists a subset of vertices  $X \subset V(D)$  such that  $S = \partial^+(X)$  and both induced subdigraphs  $D[X]$  and  $D[\bar{X}]$  contain a connected subdigraph with order 3.

Let  $Y = \bar{X}$ , and let  $X_i = \{x \in X : |N^+(x) \cap Y| = i\}$ ,  $Y_i = \{y \in Y : |N^-(y) \cap X| = i\}$ ,  $i = 0, 1, 2$ , and let  $X_3 = \{x \in X : |N^+(x) \cap Y| \geq 3\}$ ,  $Y_3 = \{y \in Y : |N^-(y) \cap X| \geq 3\}$ .

**Claim 1.**  $\min\{|X|, |Y|\} \geq 4$ .

**Proof.** Suppose that  $|X| = 3$ . Then  $\lambda^3(D) = |S| = |\partial^+(X)| \geq \xi(X) \geq \xi^3(D)$ , contrary to the assumption. Suppose that  $|Y| = 3$ . Then  $\lambda^3(D) = |S| = |\partial^-(Y)| \geq \xi(Y) \geq \xi^3(D)$ , contrary to the assumption. Claim 1 follows. □

**Claim 2.**  $X_0 = Y_0 = \emptyset$ .

**Proof.** By symmetry, we only prove that  $X_0 = \emptyset$  by contradiction. Suppose  $X_0 \neq \emptyset$  and let  $x \in X_0$ . Then for any  $\bar{x} \in Y$ ,  $x\bar{x} \notin A(D)$  and we have that  $5 \leq |N^+(x) \cap N^-(\bar{x})| = |N^+(x) \cap N^-(\bar{x}) \cap X| + |N^+(x) \cap N^-(\bar{x}) \cap Y| \leq |N^-(\bar{x}) \cap X| + |N^+(x) \cap Y| = |N^-(\bar{x}) \cap X|$ . It implies that  $|N^-(\bar{x}) \cap X| \geq 5$ . Therefore  $Y \subseteq Y_3$ . So  $D$  is  $\lambda^3$ -optimal by Lemma 4, a contradiction to our assumption. □

Combining Claim 2 with Lemma 4, we have that  $Y_1 \cup Y_2 \neq \emptyset$  and  $X_1 \cup X_2 \neq \emptyset$ . Otherwise we will obtain that  $D$  is  $\lambda^3$ -optimal, which is a contradiction. Next, we consider two cases.

*Case 1.*  $X_1 \neq \emptyset$ . Let  $x' \in X_1$  and suppose  $N^+(x') \cap Y = \{y'\}$ . Then for any  $y \in Y \setminus \{y'\}$ ,  $x'y \notin A(D)$ , so we have that  $5 \leq |N^+(x') \cap N^-(y)| = |N^+(x') \cap N^-(y) \cap X| + |N^+(x') \cap N^-(y) \cap Y| \leq |N^-(y) \cap X| + |N^+(x') \cap Y| = |N^-(y) \cap X| + 1$ . So  $|N^-(y) \cap X| \geq 4$  and  $Y \setminus \{y'\} \subseteq Y_3$ . On the other hand, since  $Y_1 \cup Y_2 \neq \emptyset$ , so  $y' \in Y_1 \cup Y_2$ . Besides,  $5 \leq \delta(D) \leq \delta^-(y') = |N^-(y')| = |N^-(y') \cap Y| + |N^-(y') \cap X| \leq |N^-(y') \cap Y| + 2$ , thus  $|N^-(y') \cap Y| \geq 3$ . Let  $y_1, y_2 \in N^-(y') \cap Y$ , then  $D[y', y_1, y_2]$  is connected and  $|N^-(y) \cap X| \geq 4$  for any  $y \in Y \setminus \{y', y_1, y_2\}$ . By Lemma 4, we have that  $D$  is  $\lambda^3$ -optimal, a contradiction.

*Case 2.*  $X_2 \neq \emptyset$ . Let  $x' \in X_2$  and suppose  $N^+(x') \cap Y = \{y', y''\}$ . Then for any  $y \in Y \setminus \{y', y''\}$ ,  $x'y \notin A(D)$ , thus  $5 \leq |N^+(x') \cap N^-(y)| = |N^+(x') \cap N^-(y) \cap X| + |N^+(x') \cap N^-(y) \cap Y| \leq |N^-(y) \cap X| + |N^+(x') \cap Y| = |N^-(y) \cap X| + 2$ . So  $|N^-(y) \cap X| \geq 3$  and  $Y \setminus \{y', y''\} \subseteq Y_3$ . On the other hand, since  $Y_1 \cup Y_2 \neq \emptyset$ ,  $y' \in Y_1 \cup Y_2$  or  $y'' \in Y_1 \cup Y_2$ . If  $|Y_1 \cup Y_2| = 1$ , then we can prove that  $D$  is  $\lambda^3$ -optimal by a proof similar to Case 1, which is a contradiction. If  $Y_1 \cup Y_2 = \{y', y''\}$ , then we consider two subcases.

*Subcase 2.1.*  $y'y'' \in A(D)$  or  $y''y' \in A(D)$ . Since  $y'' \in Y_1 \cup Y_2$  and  $\delta(D) \geq 5$ , then there exists  $y_1 \in N^-(y'') \cap Y$  such that  $y_1 \neq y'$ . Therefore  $D[y', y'', y_1]$  is connected and  $|N^-(y) \cap X| \geq 3$  for any  $y \in Y \setminus \{y', y'', y_1\}$ . By Lemma 4, we have that  $D$  is  $\lambda^3$ -optimal, a contradiction.

*Subcase 2.2.*  $y'y'' \notin A(D)$  and  $y''y' \notin A(D)$ . Since  $y'y'' \notin A(D)$  and  $y''y' \notin A(D)$ , then  $5 \leq |N^+(y') \cap N^-(y'')| = |N^+(y') \cap N^-(y'') \cap X| + |N^+(y') \cap N^-(y'') \cap Y| \leq |N^-(y'') \cap X| + |N^+(y') \cap N^-(y'') \cap Y| \leq 2 + |N^+(y') \cap N^-(y'') \cap Y|$ . Therefore  $|N^+(y') \cap N^-(y'') \cap Y| \geq 3$ . Let  $y_1 \in N^+(y') \cap N^-(y'') \cap Y$ . Then  $D[y', y'', y_1]$  is connected and  $|N^-(y) \cap X| \geq 3$  for any  $y \in Y \setminus \{y', y'', y_1\}$ . By Lemma 4, we have that  $D$  is  $\lambda^3$ -optimal, a contradiction.

The proof is complete. ■

From Theorem 1, we have following corollaries.

**Corollary 6.** *Let  $D$  be a digraph with  $|V(D)| \geq 6$ . If  $d^+(u) + d^-(v) \geq |V(D)| + 3$  for any  $u, v \in V(D)$  with  $uv \notin A(D)$ , then  $D$  is  $\lambda^3$ -optimal.*

**Corollary 7** (Theorem 1.7 in [10]). *Let  $D$  be a digraph with  $|V(D)| \geq 6$ . If  $\delta(D) \geq \frac{|V(D)|+3}{2}$ , then  $D$  is  $\lambda^3$ -optimal.*

**Remark 8.** To show the condition that “ $|N^+(u) \cap N^-(v)| \geq 5$  for any  $u, v \in V(D)$  with  $uv \notin A(D)$ ” in Theorem 1 is sharp, we give a class of digraphs. Let  $m, k$  be positive integers with  $m \geq 3$ , and let  $D$  be a digraph with  $|V(D)| = 4m + 4$ .

Define the vertex set of  $D$  as  $V(D) = B \cup C$ , where  $B = \{x_0, \dots, x_m, w_0, \dots, w_m\}$  and  $C = \{y_0, \dots, y_m, z_0, \dots, z_m\}$ . And define the arc set of  $D$  as  $A(D) = A(D[B]) \cup A(D[C]) \cup M_1 \cup M_2 \cup M_3 \cup M_4$ , where  $A(D[B]) \cup A(D[C]) = \{uv : \text{for any } u, v \in B \text{ or } C\}$ ,  $M_1 = \{x_i y_{k(\bmod m+1)} : 0 \leq i \leq m \text{ and } 0 \leq k - i \leq 1\}$ ,  $M_2 = \{w_i z_{k(\bmod m+1)} : 0 \leq i \leq m \text{ and } 0 \leq k - i \leq 2\}$ ,  $M_3 = \{y_i x_{k(\bmod m+1)} : 0 \leq i \leq m \text{ and } 0 \leq k - i \leq 2\}$  and  $M_4 = \{z_i w_{k(\bmod m+1)} : 0 \leq i \leq m \text{ and } 0 \leq k - i \leq 2\}$ .

Clearly,  $D$  is strong and there exists  $0 \leq i, j \leq m$  such that  $|N^+(x_i) \cap N^-(y_j)| = 4$  and  $x_i y_j \notin A(D)$ . And  $\partial^+(B)$  is a 3-restricted edge cut with  $|\partial^+(B)| = (2+3) \cdot (m+1) = 5m+5$ . On the other hand,  $\xi^3(D) = \xi(\{x_l, x_p, x_q\}) = |\partial^+(\{x_l, x_p, x_q\})| = 3 \cdot (2m+3) - 6 = 6m+3$ , where  $0 \leq l, p, q \leq m$ . So  $\lambda^3(D) \leq |\partial^+(B)| = 5m+5 < 6m+3 = \xi^3(D)$  for  $m \geq 3$ . Thus  $D$  is not  $\lambda^3$ -optimal.

Besides, in  $D$ , there exists  $0 \leq i, j \leq m$  such that  $x_i y_j \notin A(D)$  and  $d^+(x_i) + d^-(y_j) = 2 \cdot (2m+3) = |V(D)| + 2 < |V(D)| + 3$ , and  $\delta(D) = 2m+3 = \frac{|V(D)|}{2} + 1 < \frac{|V(D)|+3}{2}$ . So this example also shows that the conditions that “ $d^+(u) + d^-(v) \geq |V(D)| + 3$  for any  $u, v \in V(D)$  with  $uv \notin A(D)$ ” in Corollary 6 and “ $\delta(D) \geq \frac{|V(D)|+3}{2}$ ” in Corollary 7 are sharp.

### 3. PROOF OF THEOREM 2

We first introduce several useful lemmas.

**Lemma 9** (Lemma 2.1 in [10]). *Let  $D$  be a strong digraph and  $X_1, Y_1$  disjoint subsets of  $V(D)$ . If  $D[X_1]$  contains a connected subdigraph with order at least  $k$  and  $D[Y_1]$  contains a strong subdigraph with order at least  $k$ , then  $D$  is  $\lambda^k$ -connected and each arc set in  $\{\partial^-(Y_1), \partial^+(Y_1)\} \cup \Omega(X_1)$  is a  $k$ -restricted arc cut of  $D$ .*

**Lemma 10.** *Let  $D = (X, Y, A(D))$  be a strong bipartite digraph with  $\delta^+(D) \geq 3$  or  $\delta^-(D) \geq 3$ . Then  $D$  is  $\lambda^3$ -connected and  $\lambda^3(D) \leq \xi^3(D)$ .*

**Proof.** By symmetry, we only consider the case that  $\delta^-(D) \geq 3$ . Let  $X'$  be a subset of  $V(D)$  with  $|X'| = 3$  such that  $D[X']$  is connected and  $\xi^3(D) = \xi(X')$ . Without loss of generality, assume that  $|X' \cap X| = 1$  and  $|X' \cap Y| = 2$ . Let  $X' \cap X = \{x\}$  and  $X' \cap Y = \{y, z\}$ . Let  $D_1, \dots, D_p$  be an acyclic ordering of the strong components of  $D \setminus X'$ .

First, we claim that  $V(D_1) \cap Y \neq \emptyset$ . Otherwise, we have that  $V(D_1) \subseteq X$  and  $|V(D_1)| = 1$ . Let  $V(D_1) = \{u\}$ . Then  $N^-(u) \subseteq \{y, z\}$ . So  $3 \leq \delta^-(D) \leq d^-(u) = |N^-(u)| \leq |\{y, z\}| = 2$ , a contradiction. Next, we aim to prove  $|V(D_1)| \geq 3$ .

Since  $N^-(v) \subseteq \{x\} \cup (V(D_1) \cap X)$  for any  $v \in V(D_1) \cap Y$ , we have  $3 \leq \delta^-(D) \leq d^-(v) = |N^-(v)| \leq |\{x\} \cup (V(D_1) \cap X)| = |\{x\}| + |V(D_1) \cap X| = 1 + |V(D_1) \cap X|$ . Thus  $|V(D_1) \cap X| \geq 2$  and  $|V(D_1)| = |V(D_1) \cap X| + |V(D_1) \cap Y| \geq 2 + 1 = 3$ . It follows that  $|V(D_1)| \geq 3$ . Since  $D[X']$  is connected and

$D[X'] \subseteq D \setminus V(D_1)$ , by Lemma 9, each arc set in  $\Omega(X')$  is a 3-restricted arc cut of  $D$ . Therefore,  $D$  is  $\lambda^3$ -connected and  $\lambda^3(D) \leq \xi(X') = \xi^3(D)$ . ■

**Lemma 11.** *Let  $D = (X, Y, A(D))$  be a strong bipartite digraph with  $\delta^+(D) \geq 3$  or  $\delta^-(D) \geq 3$ , and let  $S = \partial^+(X')$  be a  $\lambda^3$ -cut of  $D$ , where  $X'$  is a subset of  $V(D)$ . If  $D[X']$  contains a connected subdigraph  $B$  with order 3 such that  $|N^+(x) \cap \overline{X'}| \geq 2$  for any  $x \in X' \setminus V(B)$  or  $D[\overline{X'}]$  contains a connected subdigraph  $C$  with order 3 such that  $|N^-(y) \cap X'| \geq 2$  for any  $y \in \overline{X'} \setminus V(C)$ , then  $D$  is  $\lambda^3$ -optimal.*

**Proof.** By Lemma 10,  $D$  is  $\lambda^3$ -connected and  $\lambda^3(D) \leq \xi^3(D)$ . By symmetry, we only prove the case that  $D[X']$  contains a connected subdigraph  $B$  with order 3 such that  $|N^+(x) \cap \overline{X'}| \geq 2$  for any  $x \in X' \setminus V(B)$ . The hypotheses imply that

$$\begin{aligned} \xi^3(D) &\leq |\partial^+(V(B))| = |[V(B), X' \setminus V(B)]| + |[V(B), \overline{X'}]| \\ &\leq 2|X' \setminus V(B)| + |[V(B), \overline{X'}]| \leq \sum_{x \in X' \setminus V(B)} |N^+(x) \cap \overline{X'}| + |[V(B), \overline{X'}]| \\ &= |[X' \setminus V(B), \overline{X'}]| + |[V(B), \overline{X'}]| = |[X', \overline{X'}]| = |S| = \lambda^3(D). \end{aligned}$$

Thus  $\lambda^3(D) = \xi^3(D)$  and  $D$  is  $\lambda^3$ -optimal. ■

By a proof similar to that of Lemma 4.1 shown in [10], we can get the following Lemma 12.

**Lemma 12.** *Let  $D = (X, Y, A(D))$  be a strong bipartite digraph with  $\delta(D) \geq 3$ , and let  $S$  be a  $\lambda^3$ -cut of  $D$ . If  $D$  is not  $\lambda^3$ -optimal, then there exists a subset of vertices  $X' \subset V(D)$  such that  $S = \partial^+(X')$  and both induced subdigraphs  $D[X']$  and  $D[\overline{X'}]$  contain a connected subdigraph with order 3.*

**Proof of Theorem 2.** Since  $|V(D)| \geq 6$ , for any  $u, v \in V(D)$  in the same partite,  $|N^+(u) \cap N^-(v)| \geq \left\lceil \frac{|V(D)|}{4} \right\rceil + 1 \geq 3$ . Therefore  $D$  is strong and  $\delta(D) \geq 3$ . By Lemma 10,  $D$  is  $\lambda^3$ -connected and  $\lambda^3(D) \leq \xi^3(D)$ . Suppose, on the contrary, that  $D$  is not  $\lambda^3$ -optimal, that is,  $\lambda^3(D) < \xi^3(D)$ . Let  $S$  be a  $\lambda^3$ -cut of  $D$ . Then by Lemma 12, there exists a subset of vertices  $X' \subset V(D)$  such that  $S = \partial^+(X')$  and both induced subdigraphs  $D[X']$  and  $D[\overline{X'}]$  contain a connected subdigraph with order 3.

Let  $\overline{X'} = X''$ , and let  $X'_X = X' \cap X$ ,  $X'_Y = X' \cap Y$ ,  $X''_X = X'' \cap X$  and  $X''_Y = X'' \cap Y$ . And let  $X'_{Xi} = \{x \in X'_X : |N^+(x) \cap X''_Y| = i\}$ ,  $X'_{Yi} = \{y \in X'_Y : |N^+(y) \cap X''_X| = i\}$ ,  $X''_{Xi} = \{x \in X''_X : |N^-(x) \cap X'_Y| = i\}$ ,  $X''_{Yi} = \{y \in X''_Y : |N^-(y) \cap X'_X| = i\}$ ,  $i = 0, 1$ , and  $X'_{X2} = \{x \in X'_X : |N^+(x) \cap X''_Y| \geq 2\}$ ,  $X'_{Y2} = \{y \in X'_Y : |N^+(y) \cap X''_X| \geq 2\}$ ,  $X''_{X2} = \{x \in X''_X : |N^-(x) \cap X'_Y| \geq 2\}$ ,  $X''_{Y2} = \{y \in X''_Y : |N^-(y) \cap X'_X| \geq 2\}$ .

**Claim 1.**  $\min\{|X'_X|, |X'_Y|, |X''_X|, |X''_Y|\} \geq 2$ .

**Proof.** If, on the contrary  $|X'_X| = 1$ , let  $X'_X = \{v\}$ . Then  $|N(v) \cap X'_Y| \geq 2$  for  $D[X']$  contains a connected subdigraph with order 3. Let  $y_1, y_2 \in N(v) \cap X'_Y$ . Then  $D[v, y_1, y_2]$  is connected, and for any  $x' \in X' \setminus \{v, y_1, y_2\}$ ,  $N^+(x') \subseteq \{v\} \cup (N^+(x') \cap X'')$ , we have  $3 \leq \delta(D) \leq d^+(x') = |N^+(x')| \leq |\{v\}| + |N^+(x') \cap X''| = 1 + |N^+(x') \cap X''|$ . Therefore  $|N^+(x') \cap X''| \geq 2$ . By Lemma 11,  $D$  is  $\lambda^3$ -optimal, a contradiction to our assumption. Thus  $|X'_X| \geq 2$ . Similarly, we can prove that  $\min\{|X'_Y|, |X''_X|, |X''_Y|\} \geq 2$ .  $\square$

**Claim 2.** Either  $X'_{X_0} = \emptyset$  or  $X''_{X_0} = \emptyset$  and either  $X'_{Y_0} = \emptyset$  or  $X''_{Y_0} = \emptyset$ .

**Proof.** If  $X'_{X_0} \neq \emptyset$  and  $X''_{X_0} \neq \emptyset$ , then there exists  $x \in X'_{X_0} \subseteq X$  and  $\bar{x} \in X''_{X_0} \subseteq X$  such that  $|N^+(x) \cap N^-(\bar{x})| \geq \left\lceil \frac{|V(D)|}{4} \right\rceil + 1$ . On the other hand, since  $x \in X'_{X_0}$  and  $\bar{x} \in X''_{X_0}$ ,  $N^+(x) \subseteq X'_Y$  and  $N^-(\bar{x}) \subseteq X''_Y$ , which implies that  $N^+(x) \cap N^-(\bar{x}) = \emptyset$ , a contradiction. Thus either  $X'_{X_0} = \emptyset$  or  $X''_{X_0} = \emptyset$ . Similarly, we can obtain that either  $X'_{Y_0} = \emptyset$  or  $X''_{Y_0} = \emptyset$ .  $\square$

We consider the following two cases.

*Case 1.*  $X'_{X_0} = X'_{Y_0} = \emptyset$  or  $X''_{X_0} = X''_{Y_0} = \emptyset$ . By symmetry, we only prove the case that  $X'_{X_0} = X'_{Y_0} = \emptyset$ .

**Claim 1.1.** Either  $X'_{X_1} = \emptyset$  and  $X'_{Y_1} \neq \emptyset$  or  $X'_{X_1} \neq \emptyset$  and  $X'_{Y_1} = \emptyset$ .

**Proof.** Since  $D$  is not  $\lambda^3$ -optimal, by Lemma 11, we have that  $X'_{X_1} \cup X'_{Y_1} \neq \emptyset$ . Suppose  $X'_{X_1} \neq \emptyset$  and  $X'_{Y_1} \neq \emptyset$ . Take  $x_1 \in X'_{X_1}$ . Then for any  $\bar{x} \in X''_X$ , we have that  $\left\lceil \frac{|V(D)|}{4} \right\rceil + 1 \leq |N^+(x_1) \cap N^-(\bar{x})| = |N^+(x_1) \cap N^-(\bar{x}) \cap X'| + |N^+(x_1) \cap N^-(\bar{x}) \cap X''| \leq |N^-(\bar{x}) \cap X'| + |N^+(x_1) \cap X''| = |N^-(\bar{x}) \cap X'| + 1$ . It implies that  $|N^-(\bar{x}) \cap X'| \geq \left\lceil \frac{|V(D)|}{4} \right\rceil \geq 2$ . So  $X''_X \subseteq X''_{X_2}$ . By a similar proof, we can also prove that  $X''_Y \subseteq X''_{Y_2}$ . Therefore  $D$  is  $\lambda^3$ -optimal by Lemma 11, a contradiction. The proof of Claim 1.1 is complete.  $\square$

Without loss of generality, let  $X'_{X_1} \neq \emptyset$  and  $X'_{Y_1} = \emptyset$ .

*Case 1.1.*  $|X'_{X_1}| = 1$ . Let  $x_1 \in X'_{X_1}$ . Then  $3 \leq \delta(D) \leq d^+(x_1) = |N^+(x_1)| = |N^+(x_1) \cap X'_Y| + |N^+(x_1) \cap X''_Y| = |N^+(x_1) \cap X'_Y| + 1$ , therefore  $|N^+(x_1) \cap X'_Y| \geq 2$ . Let  $y_1, y_2 \in N^+(x_1) \cap X'_Y$ . Then  $D[x_1, y_1, y_2]$  is connected, and for any  $v \in X' \setminus \{x_1, y_1, y_2\}$ ,  $|N^+(v) \cap X''| \geq 2$ . By Lemma 11,  $D$  is  $\lambda^3$ -optimal, a contradiction.

*Case 1.2.*  $|X'_{X_1}| \geq 2$ . Let  $x_1, x_2 \in X'_{X_1}$ . Then  $\left\lceil \frac{|V(D)|}{4} \right\rceil + 1 \leq |N^+(x_1) \cap N^-(x_2)| = |N^+(x_1) \cap N^-(x_2) \cap X'_Y| + |N^+(x_1) \cap N^-(x_2) \cap X''_Y| \leq |N^+(x_1) \cap N^-(x_2) \cap X'_Y| + |N^+(x_1) \cap X''_Y| = |N^+(x_1) \cap N^-(x_2) \cap X'_Y| + 1$ . So  $|N^+(x_1) \cap$



$|N^-(x_2) \cap X'_Y| \geq \left\lceil \frac{|V(D)|}{4} \right\rceil \geq 2$ . Let  $y_1 \in N^+(x_1) \cap N^-(x_2) \cap X'_Y$ . Then

$$\begin{aligned} \xi^3(D) &\leq \xi(\{x_1, x_2, y_1\}) \leq |\partial^+(\{x_1, x_2, y_1\})| \\ &= |[\{x_1\}, X'_Y \setminus \{y_1\}]| + |[\{x_1\}, X''_Y]| + |[\{x_2\}, X'_Y \setminus \{y_1\}]| + |[\{x_2\}, X''_Y]| \\ &\quad + |[\{y_1\}, X'_X \setminus \{x_1, x_2\}]| + |[\{y_1\}, X''_X]| \\ &\leq 2 \cdot (|X'_Y| - 1) + 2 + |X'_X| - 2 + |[\{y_1\}, X''_X]| \leq |S| = \lambda^3(D). \end{aligned}$$

Thus  $D$  is  $\lambda^3$ -optimal, a contradiction.

*Case 2.*  $X'_{X0} = X''_{Y0} = \emptyset$  or  $X''_{X0} = X'_{Y0} = \emptyset$ . By symmetry, we only prove the case that  $X'_{X0} = X''_{Y0} = \emptyset$ . Without loss of generality, we may assume that  $X'_{Y0} \neq \emptyset$  and  $X''_{X0} \neq \emptyset$ . Otherwise, by Case 1,  $D$  is  $\lambda^3$ -optimal, a contradiction. On the other hand, since for any  $u \in X'_{Y0}$ ,  $N^+(u) \subseteq X'_X$ , we have  $\left\lceil \frac{|V(D)|}{4} \right\rceil + 1 \leq \delta(D) \leq d^+(u) = |N^+(u)| \leq |X'_X|$ . Therefore  $|X'_X| \geq \left\lceil \frac{|V(D)|}{4} \right\rceil + 1$ . Similarly, we can also prove that  $|X''_Y| \geq \left\lceil \frac{|V(D)|}{4} \right\rceil + 1$ . Thus

$$\begin{aligned} (1) \quad |X'_Y| + |X''_X| &= |V(D)| - |X'_X| - |X''_Y| \\ &\leq |V(D)| - 2 \cdot \left( \left\lceil \frac{|V(D)|}{4} \right\rceil + 1 \right) \leq \frac{|V(D)|}{2} - 2. \end{aligned}$$

**Claim 2.1.**  $|X'_X| \geq |X'_Y| + 1$  or  $|X''_Y| \geq |X''_X| + 1$ .

**Proof.** Otherwise, we have that  $|X'_Y| + |X''_X| \geq |X'_X| + |X''_Y| \geq 2 \cdot \left( \left\lceil \frac{|V(D)|}{4} \right\rceil + 1 \right) \geq \frac{|V(D)|}{2} + 2$ , a contradiction to (1).  $\square$

Without loss of generality, we assume that  $|X'_X| \geq |X'_Y| + 1$  in the following discussion.

**Claim 2.2.**  $|N^+(x) \cap X''_Y| \geq 3$  and  $|N^-(y) \cap X'_X| \geq 3$  for any  $x \in X'_X$  and  $y \in X''_Y$ .

**Proof.** By symmetry, we only prove that for any  $x \in X'_X$ ,  $|N^+(x) \cap X''_Y| \geq 3$ . Since  $X''_{X0} \neq \emptyset$ , for any  $x \in X'_X$  and  $\bar{x} \in X''_{X0}$ ,  $\left\lceil \frac{|V(D)|}{4} \right\rceil + 1 \leq |N^+(x) \cap N^-(\bar{x})| = |N^+(x) \cap N^-(\bar{x}) \cap X'_Y| + |N^+(x) \cap N^-(\bar{x}) \cap X''_Y| \leq |N^-(\bar{x}) \cap X'_Y| + |N^+(x) \cap X''_Y| = |N^+(x) \cap X''_Y|$ , so  $|N^+(x) \cap X''_Y| \geq \left\lceil \frac{|V(D)|}{4} \right\rceil + 1 \geq 3$ .  $\square$

**Claim 2.3.**  $X'_{Y2} = X''_{X2} = \emptyset$ .

**Proof.** Here, we only prove that  $X'_{Y2} = \emptyset$ . The proof of the statement that  $X''_{X2} = \emptyset$  is similar. Suppose, by a contradiction, there exists  $y \in X'_{Y2}$ . Let

$x_1, x_2 \in N^+(y) \cap X_X''$ . Then

$$\begin{aligned}
\xi^3(D) &\leq \xi(\{x_1, x_2, y\}) \leq |\partial^+(\{y\}) \cup \partial^-(\{x_1, x_2\})| \\
&= |\partial^+(\{y\})| + |\partial^-(\{x_1, x_2\})| - 2 = |[\{y\}, X_X']| + |[\{y\}, X_X'']| \\
&\quad + |[X_Y', \{x_1\}]| + |[X_Y'', \{x_1\}]| + |[X_Y', \{x_2\}]| + |[X_Y'', \{x_2\}]| - 2 \\
&\leq |X_X'| + |[\{y\}, X_X'']| + |[X_Y', \{x_1\}]| + 2|X_Y''| + |[X_Y', \{x_2\}]| - 2 \\
&\leq 3 \max\{|X_X'|, |X_Y''|\} + |[\{y\}, X_X'']| + |[X_Y', \{x_1\}]| \\
&\quad + |[X_Y', \{x_2\}]| - 2 \leq |S| = \lambda^3(D).
\end{aligned}$$

So  $D$  is  $\lambda^3$ -optimal, a contradiction.  $\square$

**Claim 2.4.** For any  $x \in X_X'$ ,  $|N(X) \cap X_Y'| \geq 2$ .

**Proof.** Let  $X_Y' = \{y_1, y_2, \dots, y_p\}$  and let  $S^* = \{s^* : s^* \in N^+(y_i) \cap N^-(y_j) \cap X_X', \text{ where } i, j \in \{1, \dots, p\} \text{ and } i \neq j\}$ . Then  $D[S^* \cup X_Y']$  is strong. Besides, by Claim 2.3, we have that for any  $i, j \in \{1, \dots, p\}$  and  $i \neq j$ ,  $y_i, y_j \in X_{Y_0}' \cup X_{Y_1}'$ . Therefore  $\left\lceil \frac{|V(D)|}{4} \right\rceil + 1 \leq |N^+(y_i) \cap N^-(y_j)| = |N^+(y_i) \cap N^-(y_j) \cap X_X'| + |N^+(y_i) \cap N^-(y_j) \cap X_X''| \leq |N^+(y_i) \cap N^-(y_j) \cap X_X'| + |N^+(y_i) \cap X_X''| \leq |N^+(y_i) \cap N^-(y_j) \cap X_X'| + 1$ . So  $|N^+(y_i) \cap N^-(y_j) \cap X_X'| \geq \left\lceil \frac{|V(D)|}{4} \right\rceil \geq 2$ . Similarly, we can prove that  $|N^+(y_j) \cap N^-(y_i) \cap X_X'| \geq 2$ . On the other hand, since  $|X_Y'| \geq 2$ , we have  $|S^* \cup X_Y'| \geq 3$ . For any  $x \in S^*$ , clearly,  $|N(x) \cap X_Y'| \geq 2$ . Next, we claim that for any  $x \in X_X' \setminus S^*$ ,  $|N^+(x) \cap X_Y''| \leq |[X_Y', \{x\}]|$ .

Suppose there exists  $x^* \in X_X' \setminus S^*$  such that  $|N^+(x^*) \cap X_Y''| > |[X_Y', \{x^*\}]|$ . Since  $D[S^* \cup X_Y']$  is strong and  $|S^* \cup X_Y'| \geq 3$ , we have  $X' \setminus \{x^*\}$  is a 3-restricted edge cut. Therefore  $|\partial^+(X' \setminus \{x^*\})| = |S| - |N^+(x^*) \cap X_Y''| + |[X_Y', \{x^*\}]| < |S|$ , a contradiction to the minimality of  $S$ . Thus  $|[X_Y', \{x\}]| \geq |N^+(x) \cap X_Y''|$ . By Claim 2.2, we have that  $|[X_Y', \{x\}]| \geq 3$ . The proof of Claim 2.4 is complete.  $\square$

Let  $x_1 \in X_X'$  such that  $|N^+(x_1) \cap X_Y''| \leq |N^+(u) \cap X_Y''|$  for any  $u \in X_X'$ , and let  $y_1, y_2 \in N(x_1) \cap X_Y'$ . Then

$$\begin{aligned}
\xi^3(D) &\leq |\partial^+(\{x_1, y_1, y_2\})| = |[\{x_1, y_1, y_2\}, X' \setminus \{x_1, y_1, y_2\}]| + |[\{x_1, y_1, y_2\}, X'']| \\
&\leq 2(|X_X'| - 1) + |X_Y'| - 2 + |[\{x_1\}, X_Y'']| + |[\{y_1\}, X_X'']| + |[\{y_2\}, X_X'']| \\
&\leq 3|X_X'| - 5 + |[\{x_1\}, X_Y'']| + |[\{y_1\}, X_X'']| \\
&\quad + |[\{y_2\}, X_X'']| (|X_X'| \geq |X_Y'| + 1) \leq |[\{x_1\}, X_Y'']| \times |X_X'| \\
&\quad + |[\{y_1\}, X_X'']| + |[\{y_2\}, X_X'']| \leq |S| = \lambda^3(D).
\end{aligned}$$

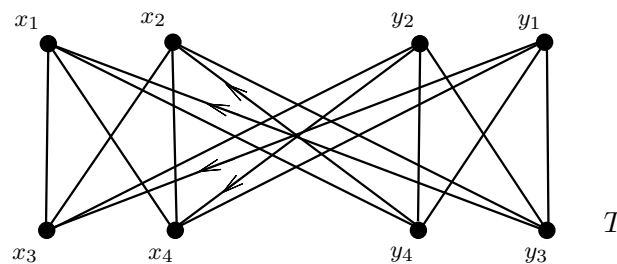
So  $D$  is  $\lambda_3$ -optimal, a contradiction.

The proof is complete.  $\blacksquare$

From Theorem 2, we have following corollaries.

**Corollary 13.** *Let  $D = (X, Y, A(D))$  be a strong bipartite digraph with  $\delta(D) \geq 3$ . If for any  $u, v \in V(D)$  in the same partite,  $d^+(u) + d^-(v) \geq |V(D)| - 1$ , then  $D$  is  $\lambda^3$ -optimal.*

**Corollary 14.** *Let  $D = (X, Y, A(D))$  be a strong bipartite digraph with  $|V(D)| \geq 6$ . If  $\delta(D) \geq \left\lfloor \frac{|V(D)|}{2} \right\rfloor$ , then  $D$  is  $\lambda^3$ -optimal.*



(Unordered edges represent two arcs with the same end-vertices and opposite directions.)

Figure 1. The example from Remark 15.

**Remark 15.** To show that the condition “ $|N^+(u) \cap N^-(v)| \geq \left\lceil \frac{|V(D)|}{4} \right\rceil + 1$  for any  $u, v \in V(D)$  in the same partite” in Theorem 2 is sharp, we consider the digraph  $T$  shown in Figure 1. Clearly,  $|V(D)| \geq 6$  and  $D$  is strong. There exists  $x_1, y_1$  in the same partite such that  $|N^+(x_1) \cap N^-(y_1)| = 2 < 3 = \left\lceil \frac{|V(T)|}{4} \right\rceil + 1$ . Clearly,  $\partial^+(\{x_1, x_2, x_3, x_4\})$  is a 3-restricted edge cut and  $\xi^3(T) = |\partial^+(\{x_1, x_2, x_3\})| = 5$ . Therefore,  $\lambda^3(T) \leq |\partial^+(\{x_1, x_2, x_3, x_4\})| = 4 < 5 = \xi^3(T)$  and  $T$  is not  $\lambda^3$ -optimal.

Besides, since  $d^+(x_3) + d^-(y_4) = 6 < 7 = |V(T)| - 1$  and  $\delta(T) = 3 < 4 = \left\lfloor \frac{|V(D)|}{2} \right\rfloor$ , this example also shows that the conditions “ $d^+(u) + d^-(v) \geq |V(D)| - 1$  for any  $u, v \in V(D)$  in the same partite” in Corollary 13 and “ $\delta(D) \geq \left\lfloor \frac{|V(D)|}{2} \right\rfloor$ ” in Corollary 14 are sharp.

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