$T ext{-}\mathrm{COLORINGS}, \, \mathrm{DIVISIBILITY} \,\, \mathrm{AND} \,\, \mathrm{THE} \,\, \mathrm{CIRCULAR}$ CHROMATIC NUMBER

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Abstract

Let T be a T-set, i.e., a finite set of nonnegative integers satisfying $0 \in T$, and G be a graph. In the paper we study relations between the T-edge spans $\exp_T(G)$ and $\exp_{d \cap T}(G)$, where d is a positive integer and

$$d \odot T = \{0 \le t \le d \pmod{T+1} \colon d \mid t \Rightarrow t/d \in T\}.$$

We show that $\exp_{d\odot T}(G)=d\exp_T(G)-r$, where $r,\,0\leq r\leq d-1$, is an integer that depends on T and G. Next we focus on the case $T=\{0\}$ and show that

$$\exp_{d\odot\{0\}}(G) = \lceil d(\chi_c(G) - 1) \rceil,$$

where $\chi_c(G)$ is the circular chromatic number of G. This result allows us to formulate several interesting conclusions that include a new formula for the circular chromatic number

$$\chi_c(G) = 1 + \inf \{ \exp_{d \odot \{0\}}(G)/d : d \ge 1 \}$$

and a proof that the formula for the T-edge span of powers of cycles, stated as conjecture in [Y. Zhao, W. He and R. Cao, The edge span of T-coloring on graph C_n^d , Appl. Math. Lett. 19 (2006) 647–651], is true.

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1. Introduction

In the paper we study relations between two different generalizations of ordinary vertex colorings: T-colorings and (k,d)-colorings. Let G be a graph with n-vertex set V and edge set E. Given integers $1 \le d \le k$, by a (k,d)-coloring of G we mean any function $c \colon V \to [0,k-1]$ ($[a,b] := \{a,a+1,\ldots,b\}$ for any integers $a \le b$) such that

$$d \le |c(u) - c(v)| \le k - d$$

whenever $uv \in E$. This notion may be viewed as a generalization of a k-coloring since (k,d)-colorings of G are k-colorings of G and (k,1)-colorings are the same as k-colorings that use colors from the interval [0,k-1]. The *circular chromatic number*, introduced by Vince [12] as a generalization of the chromatic number, is defined by the formula

$$\chi_c(G) = \inf \{ k/d : G \text{ has a } (k, d) \text{-coloring} \}.$$

The circular chromatic number was studied by many authors, see [14, 15] for a survey of results. It was shown for example [12] that the distance between the circular and ordinary chromatic number does not exceed 1, i.e.

$$\chi(G) - 1 < \chi_c(G) < \chi(G)$$
.

In the same paper Vince proved two useful facts: (1) G has a (k,d)-coloring if and only if $\chi_c(G) \leq k/d$; (2) $\chi_c(G)$ is a rational number which has a form k/d, where $k \leq n$. We will use these observations to show that there is a relation between $\chi_c(G)$ and $\exp_T(G)$ the T-edge span defined below. Given a T-set T, i.e., a finite set that consists of nonnegative integers and satisfies $0 \in T$, by a T-coloring of G we mean any function $c: V \to \mathbb{Z}$ such that

$$|c(u) - c(v)| \notin T$$

whenever $uv \in E$. T-colorings were introduced as a model for the frequency assignment problem in [5]. This notion also may be viewed as a generalization of ordinary vertex colorings since T-colorings are vertex colorings and vertex

colorings are $\{0\}$ -colorings. The T-edge span, introduced by Cozzens and Roberts [1], is defined as

$$esp_T(G) = min\{esp(c): c \text{ is a } T\text{-coloring of } G\},\$$

where $\exp(c) = \max\{|c(u) - c(v)| : uv \in E\}$ is the edge span of c (if G is an empty graph then $\exp(c) = 0$). If we replace $\exp(c)$ by $\operatorname{sp}(c)$ (the span of c, i.e., $\max\{|c(u) - c(v)| : u, v \in V\}$) we will receive the T-span of G. Both parameters were studied by many authors, there are results concerning computational complexity of the problem of computing $\operatorname{sp}_T(G)$ [2, 3], the behaviour of the greedy algorithm [7] and formulas describing $\operatorname{sp}_T(G)$ and $\operatorname{esp}_T(G)$ for some T-sets T and some graphs G [8, 9, 13].

The remainder of the paper is organized as follows. In Section 2 we study relations between $\exp_T(G)$ and $\exp_{d\odot T}(G)$, where d is a positive integer and $d\odot T=\{0\leq t\leq d (\max T+1)\colon d\,|\, t\Rightarrow t/d\in T\}$. We show that $\exp_{d\odot T}(G)=d\exp_T(G)-r$, where $r,\,0\leq r\leq d-1$, is an integer that depends on T and G. In Section 3 we study the distance between the T-span and T-edge span and show that it cannot exceed $\max T$. We also give examples that prove that this bound is tight. Section 4 contains our main results. We show that if T is an interval, i.e., T=[0,d-1] (or equivalently $T=d\odot\{0\}$), then (k,d)-colorings $(k\geq d)$ are nonnegative T-colorings with span bounded by k-1 and edge span bounded by k-d. We use this relation to show that

$$\exp_{d \odot \{0\}}(G) = \lceil d(\chi_c(G) - 1) \rceil.$$

We also discuss whether it is possible to extend this relation to all T-sets. Using the above formula we show that

$$\chi_c(G) = 1 + \inf \{ \exp_{d \odot \{0\}}(G) / d \colon d \ge 1 \}$$

and discuss how these formulas allow us to move known results from the world of the T-edge span to the world of the circular chromatic number and vice versa. The last section is devoted to the powers of cycles investigated in [13]. The authors conjectured and partially proved that

$$\exp_{d \odot \{0\}}(C_n^p) = pd + \lceil rd/q \rceil,$$

where $q \geq 2$ and r are the quotient and the remainder of the division of n by p+1, respectively. We show that it is true in general.

2. T-Edge Span and $d \odot T$ -Edge Span

The operation \odot was introduced in [6], where it was shown that $\operatorname{sp}_{d\odot T}(G) = d\operatorname{sp}_T(G)$. Below we prove a similar formula for the T-edge span, but before we proceed we need to recall the following result.

Lemma 1 (Lemma 2.2(i) of [6]). If a and b are real numbers, then $\lfloor |a-b| \rfloor \leq \lfloor \lfloor a \rfloor - \lfloor b \rfloor \rfloor \leq \lceil |a-b| \rceil$.

Lemma 2. Let G be a graph, T be a T-set and d be a positive integer.

- (1) If c is a T-coloring of G, then dc is a $d \odot T$ -coloring of G.
- (2) If c is a $d \odot T$ -coloring of G, then $\lfloor c/d \rfloor$ is a T-coloring of G.

Proof. Let uv be an edge of G (if G is empty, then our claim is obvious).

- $(1) \text{ If } |c(u)-c(v)| \geq \max T + 1, \text{ then } |dc(u)-dc(v)| \geq d \left(\max T + 1\right) = \max d \odot T + 1. \text{ If } |c(u)-c(v)| < \max T + 1 \text{ and } |dc(u)-dc(v)| \in d \odot T, \text{ then the definition of } d \odot T \text{ gives } |c(u)-c(v)| \in T, \text{ a contradiction. Hence } |dc(u)-dc(v)| \notin d \odot T \text{ in both cases.}$
- (2) If $|c(u)-c(v)| \ge \max d\odot T+1 = d (\max T+1)$, then $|\lfloor c(u)/d \rfloor \lfloor c(v)/d \rfloor| \ge \lfloor |c(u)-c(v)|/d \rfloor \ge \max T+1$ by Lemma 1. If $|c(u)-c(v)| < \max d\odot T+1$, then the definition of $d\odot T$ gives d||c(u)-c(v)| and, by Lemma 1, $|\lfloor c(u)/d \rfloor \lfloor c(v)/d \rfloor| = |c(u)-c(v)|/d \notin T$. Hence $|\lfloor c(u)/d \rfloor \lfloor c(v)/d \rfloor| \notin T$ in both cases.

Theorem 3. Let G be a graph, T be a T-set and d be a positive integer. There is an integer $0 \le r \le d-1$ such that $\exp_{d \cap T}(G) = d \exp_T(G) - r$.

Proof. Let c be a T-coloring of G such that $\exp(c) = \exp_T(G)$. By Lemma 2, dc is a $d \odot T$ -coloring of G. Hence

(1)
$$\exp_{d \cap T}(G) \le \exp(dc) = d \exp(c) = d \exp_{T}(G).$$

Let c' be a $d \odot T$ -coloring of G such that $\exp(c') = \exp_{d \odot T}(G)$. By Lemma 2, $\lfloor c'/d \rfloor$ is a T-coloring of G. Let uv be an edge of G such that $\exp(\lfloor c'/d \rfloor) = \lfloor \lfloor c'(u)/d \rfloor - \lfloor c'(v)/d \rfloor \rfloor$ (if G is empty our claim is obvious). Then

$$d \exp_{T}(G) - d \leq d \exp(\lfloor c'/d \rfloor) - d = d|\lfloor c'(u)/d \rfloor - \lfloor c'(v)/d \rfloor| - d$$

$$\leq d \lceil |c'(u) - c'(v)|/d \rceil - d \leq d \lceil \exp(c')/d \rceil - d$$

$$= d \lceil \exp_{d \odot T}(G)/d \rceil - d < \exp_{d \odot T}(G).$$

To complete the proof it suffices to combine (1) with (2).

The open problem is a formula for r. Later we will show how to compute r provided that $T = \{0\}$ and that r can be any integer from [0, d-1].

Corollary 4. Let G be a graph, T be a T-set and d be a positive integer. Then $\exp_T(G) = \left[\exp_{d \cap T}(G)/d\right]$.

3. The Distance Between the T-Span and T-Edge Span

It is known [1] that $\exp_T(G) \leq \sup_T(G)$. We are going to show that $\sup_T(G) \leq \exp_T(G) + \max T$ and give examples in which the difference $\sup_T(G) - \exp_T(G)$ equals $\max T$.

Lemma 5. Let G be a graph and T be a T-set. If $c': V \to \mathbb{Z}$ is a T-coloring of G and $c: V \to \mathbb{Z}$ is the remainder of the division of c' by $\exp(c') + \max T + 1$, i.e., $c(v) = c'(v) \mod (\exp(c') + \max T + 1)$ for $v \in V$, then

- (1) c is a T-coloring of G.
- (2) $\operatorname{sp}(c) \le \operatorname{esp}(c') + \max T$.
- (3) $\exp(c) \le \exp(c') + \max T + 1 \min(\mathbb{N} \setminus T).$

Proof. Observe that (2) follows immediately from the definition of c. To prove (1) and (3), take an edge uv of G (if G is empty, our claim is obvious). Let q be the quotient of the division of c' by $\exp(c') + \max T + 1$. Without loss of generality we may assume that $q(u) \geq q(v)$. It is easy to see that $q(u) \leq q(v) + 1$ since otherwise

$$\exp(c') \ge |c'(u) - c'(v)|
= |(\exp(c') + \max T + 1)(q(u) - q(v)) + c(u) - c(v)|
\ge (\exp(c') + \max T + 1)|q(u) - q(v)| - |c(u) - c(v)|
\ge 2(\exp(c') + \max T + 1) - \exp(c') - \max T
= \exp(c') + \max T + 2 > \exp(c').$$

Hence there are two cases to consider.

(a) q(u) = q(v) + 1. Then $|c'(u) - c'(v)| = |(\exp(c') + \max T + 1)(q(u) - q(v)) + (c(u) - c(v))| = |\exp(c') + \max T + 1 + (c(u) - c(v))|$. Since $\exp(c') + \max T + 1 > \exp(c') \ge |c'(u) - c'(v)|$ and $|c(u) - c(v)| \le \exp(c') + \max T$, we have $|c(u) - c(v)| = \exp(c') + \max T + 1 - |c'(u) - c'(v)|$. This gives $|c(u) - c(v)| \ge \max T + 1$ and $|c(u) - c(v)| \le \exp(c') + \max T + 1 - \min(\mathbb{N} \setminus T)$ since $|c'(u) - c'(v)| \notin T$ implies $|c'(u) - c'(v)| \ge \min(\mathbb{N} \setminus T)$.

(b) q(u) = q(v). Then |c'(u) - c'(v)| = |c(u) - c(v)|, which gives $|c(u) - c(v)| \notin T$ and $|c(u) - c(v)| \le \exp(c') \le \exp(c') + \max T + 1 - \min(\mathbb{N} \setminus T)$.

Corollary 6. Let G be a graph and T be a T-set. Then

- (1) There is a T-coloring c of G such that $\operatorname{sp}(c) \leq \operatorname{esp}_T(G) + \max T$ and $\operatorname{esp}(c) \leq \operatorname{esp}_T(G) + \max T + 1 \min(\mathbb{N} \setminus T)$.
- (2) If T is an interval, then there is a T-coloring c of G such that $\exp(c) = \exp_T(G)$ and $\operatorname{sp}(c) \leq \exp_T(G) + \max T$.
- (3) $\exp_T(G) \le \sup_T(G) \le \exp_T(G) + \max T$.

Proof. (1) Let c' be a T-coloring of G satisfying $\exp(c') = \exp_T(G)$ and c be the remainder of the division of c' by $\exp_T(G) + \max T + 1$. The claim follows from Lemma 5.

- (2) Follows from (1) since $\min(\mathbb{N} \setminus T) = \max T + 1$ if T is an interval.
- (3) Follows from (1) and the definition of the T-span.

The above inequalities are tight. It is known [1] that $\exp_T(G) = \operatorname{sp}_T(G)$ for all weakly perfect graphs and all T-sets T. It is also easy to see that if T is an interval, then $\operatorname{sp}_T(C_{2n+1}) = 2 \max T + 2 (\operatorname{sp}_T(G) = (\max T + 1)(\chi(G) - 1))$ if T is an interval, see [1]) and $\operatorname{esp}_T(C_{2n+1}) = \lceil (\max T + 1)(1 + 1/n) \rceil$ (see Theorem 8) which gives $\operatorname{sp}_T(C_{2n+1}) = \operatorname{esp}_T(C_{2n+1}) + \max T$ provided that $n \geq \max T + 1$.

4. The Relation Between (k, d)-Colorings and T-Colorings

Now we are ready to prove that there is a relation between (k, d)-colorings and T-colorings provided that T is an interval.

Lemma 7. Let G be a graph and d be a positive integer. If T = [0, d-1], then for every function $c: V \to \mathbb{Z}$ and every integer $k \ge d$ the following conditions are equivalent:

- (1) c is a T-coloring of G such that $\operatorname{sp}(c) \leq k-1$ and $\operatorname{esp}(c) \leq k-d$.
- (2) $c \min c(V)$ is a (k, d)-coloring of G.

Proof. Let uv be an edge of G (our claim is obvious if G is empty) and $c' = c - \min c(V)$.

- (\Rightarrow) c is a T-coloring of G and T is an interval, so $|c'(u) c'(v)| = |c(u) c(v)| \ge d$. Moreover, $|c'(u) c'(v)| = |c(u) c(v)| \le \exp(c) \le k d$ and $c'(V) \subseteq [0, \operatorname{sp}(c)] \subseteq [0, k 1]$.
- (\Leftarrow) c' is a (k,d)-coloring of G, so $|c(u)-c(v)|=|c'(u)-c'(v)|\geq d$ and $|c(u)-c(v)|=|c'(u)-c'(v)|\leq k-d$. This proves that c is a T-coloring and gives $\exp(c)\leq k-d$. To complete the proof it suffices to observe that $c'(V)\subseteq [0,k-1]$ implies $\operatorname{sp}(c)=\operatorname{sp}(c')\leq k-1$.

Theorem 8. Let G be a graph and d be a positive integer. If T = [0, d-1], then

$$\operatorname{esp}_T(G) = \lceil d(\chi_c(G) - 1) \rceil.$$

Proof. Without loss of generality we assume that G is not empty. Then $k = \lceil d\chi_c(G) \rceil - 1 \ge d$. If $\exp_T(G) \le k - d$, then, by Corollary 6, there is a T-coloring c of G such that $\exp(c) = \exp_T(G) \le k - d$ and $\operatorname{sp}(c) \le \exp_T(G) + d - 1 \le k - 1$.

Lemma 7 implies now that $c - \min c(V)$ is a (k, d)-coloring, which finally gives $d\chi_c(G) \leq k$, a contradiction. Hence

$$esp_T(G) \ge k - d + 1.$$

On the other hand, $(k+1)/d \ge \chi_c(G)$ so there exists a (k+1,d)-coloring c of G. Without loss of generality we assume that $\min c(V) = 0$. By Lemma 7, c has to be a T-coloring of G with $\exp(c) \le k - d + 1$. This gives

$$\exp_T(G) \le k - d + 1.$$

Combining these inequalities together, we get $\exp_T(G) = k - d + 1 = \lceil d\chi_c(G) \rceil - d = \lceil d(\chi_c(G) - 1) \rceil$.

Since T is an interval, we know that $|T| = \max T + 1$ and the above formula may be expressed as

$$\operatorname{esp}_T(G) = \lceil |T|(\chi_c(G) - 1) \rceil.$$

This resembles Tesman's inequality $\operatorname{sp}_T(G) \leq |T|(\chi(G) - 1)$ which holds for all T-sets T and all graphs G [11], so it is interesting to ask the following question.

Does
$$\exp_T(G) \leq \lceil |T|(\chi_c(G) - 1) \rceil$$
 for all T-sets T and all graphs G?

Unfortunately, the answer is negative even for odd cycles. To show this, let us consider integers $1 \le k \le n-1$ and set $T = \{0, 2, ..., 2k\}$ and $G = C_{2n+1}$. Then $\lceil |T|(\chi_c(G)-1) \rceil = \lceil (k+1)(1+1/n) \rceil = k+2$ and $\exp_T(C_{2n+1}) \ge 2k+2$ since otherwise the differences of colors assigned to adjacent vertices of G in any T-coloring of G with minimal edge span would be odd and their sum would not be 0, a contradiction.

Theorem 8 shows also that the value of integer r of Theorem 3 can be arbitrary. Indeed, if we take $0 \le r \le d-1$ and a planar graph G such that $\chi_c(G) = 3 - r/d$ (which exists by [10]), then $\chi(G) = 3$ and $\exp_{d \odot \{0\}}(G) = \lceil d(\chi_c(G) - 1) \rceil = \lceil d(2 - r/d) \rceil = 2d - r = d(\chi(G) - 1) - r = d \exp_{\{0\}}(G) - r$. The open question is if this is true for all T-sets T.

Theorem 9. Let G be a graph. Then

$$\chi_c(G) = 1 + \inf \big\{ \operatorname{esp}_{d \odot \{0\}}(G) / d \colon d \ge 1 \big\}.$$

Moreover, if $\chi_c(G) = k/d$ $(1 \le d \le k)$, then $\chi_c(G) = 1 + \exp_{d \odot \{0\}}(G)/d$.

Proof. $\chi_c(G) - 1 \le \exp_{d \odot \{0\}}(G)/d$ by Theorem 8. To complete the proof it suffices to observe that if $\chi_c(G) = k/d$, then the same theorem gives $\chi_c(G) - 1 = \exp_{d \odot \{0\}}(G)/d$.

Theorems 8 and 9 have two important consequences. Firstly, if we know a formula for $\chi_c(G)$, then we can easily obtain a formula for $\exp_T(G)$ for all T-sets T that are intervals. For example, Fan [4] proved that $\chi_c(G) = \chi(G)$ if the complement of G is non-Hamiltonian, which gives

Corollary 10. If G is a graph whose complement is non-Hamiltonian, then

$$\exp_{d \odot \{0\}}(G) = d(\chi(G) - 1) = \sup_{d \odot \{0\}}(G)$$

for every $d \geq 1$.

Secondly, if the problem of computing $\chi_c(G)$ for graphs G from a certain class \mathcal{G} is polynomially solvable, then we can compute $\exp_T(G)$ for $G \in \mathcal{G}$ and any interval T in a polynomial time, too.

5. Powers of Cycles

Let $p \ge 1$ and $n \ge 2p + 2$ be integers. Let q and r are the quotient and the remainder of the division of n by p + 1, respectively.

Zhao et al. in [13] proved the following theorem.

Theorem 11. If q = pl + t for $l \ge 0$, $0 \le t \le p - 1$ such that $p \ge td$, then

$$\operatorname{esp}_{d \odot \{0\}}(C_n^p) = pd + \lceil rd/q \rceil.$$

Moreover, they conjectured that this equality holds for any $n \geq 2p + 2$, not only when $p \geq td$. We will show that it is true. Recall that it is known that if G is a n-vertex graph, then $\chi_c(G) \geq n/\alpha(G)$, where $\alpha(G)$ is the independence number of G.

Theorem 12. $\chi_c(C_n^p) = n/q$.

Proof. Let $v_0, v_1, \ldots, v_{n-1}$ be a cyclic ordering of vertices of C_n^p . We claim that a function given by

$$c(v_i) = (iq) \bmod n$$

is a (n,q)-coloring of C_n^p . Indeed, the definition of c gives $0 \le c \le n-1$ and, if $v_i v_j$ (i > j) is an edge of C_n^p , then either $1 \le i - j \le p$ and $|c(v_i) - c(v_j)| = (i - j)q$ or $1 \le n + j - i \le p$ and $|c(v_i) - c(v_j)| = (n - i + j)q$. In both cases it is easy to verify that $q \le |c(v_i) - c(v_j)| \le qp \le n - q$.

To complete the proof it suffices to observe that $\alpha(C_n^p) \leq q$ and use inequality $\chi_c(G) \geq n/\alpha(G)$.

Theorem 13. $\exp_{d\odot\{0\}}(C_n^p) = pd + \lceil rd/q \rceil$.

Proof. Follows immediately from Theorems 8 and 12.

6. Conclusion

We proved the general relation between the circular chromatic number and T-edge span for $T = d \odot \{0\}$. Moreover, we applied it to solve an open conjecture concerning the T-edge span for powers of cycles C_n^p .

Possible further fields of research include for example finding the necessary conditions for $\exp_T(G) \leq \lceil |T|(\chi_c(G)-1) \rceil$, or analyzing dependence between $\exp_T(G)$ and $\chi_c(G)$ on the structure of a set T.

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