# ORIENTED CHROMATIC NUMBER OF CARTESIAN PRODUCTS AND STRONG PRODUCTS OF PATHS 

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#### Abstract

An oriented coloring of an oriented graph $G$ is a homomorphism from $G$ to $H$ such that $H$ is without selfloops and arcs in opposite directions. We shall say that $H$ is a coloring graph. In this paper, we focus on oriented colorings of Cartesian products of two paths, called grids, and strong products of two paths, called strong-grids. We show that there exists a coloring graph with nine vertices that can be used to color every orientation of grids with five columns. We also show that there exists a strong-grid with two columns and its orientation which requires 11 colors for oriented coloring. Moreover, we show that every orientation of every strong-grid with three columns can be colored by 19 colors and that every orientation of every strong-grid with four columns can be colored by 43 colors. The above statements were proved with the help of computer programs.


Keywords: graph, oriented coloring, grid.
2010 Mathematics Subject Classification: 05C15.

## 1. Introduction

Let $G=(V(G), E(G))$ be a simple undirected graph. An orientation of $G$ is a directed graph $\vec{G}=(V(\vec{G}), A(\vec{G}))$ obtained from $G$ by ordering every edge $\{u, v\} \in E(G)$ either from $u$ to $v$ (resulting in an $\operatorname{arc}(u, v) \in A(\vec{G})$ ), or conversely (yielding an $\operatorname{arc}(v, u) \in A(\vec{G})$ ). In this paper, we shall deal with undirected graphs and their orientations. An orientation of a graph is called an oriented graph. An oriented coloring is a coloring $c$ of the vertices of an oriented graph $\vec{G}=(V(\vec{G}), A(\vec{G}))$ such that
(i) no two neighbors have the same color,
(ii) for any two arcs $(u, v)$ and $(y, z) \in A(\vec{G})$, if $c(u)=c(z)$ then $c(v) \neq c(y)$. In other words, if the arc $(y, z)$ goes from color $c(y)$ to $c(z)$, then no other arc can go in the opposite direction, i.e., from $c(z)$ to $c(y)$.

With every oriented coloring $c$ of $\vec{G}$ one can associate a digraph $\vec{H}_{c}$, called the coloring graph of $\vec{G}$, with set of vertices $V\left(\vec{H}_{c}\right)=\{c(x): x \in V(G)\}$ and set of arcs $A\left(\vec{H}_{c}\right)=\{(c(x), c(y)):(x, y) \in A(\vec{G})\}$. Due to conditions (i) and (ii), $\vec{H}_{c}$ is an oriented graph without loops and opposite arcs. An oriented coloring $c$ can then be viewed as a homomorphism (that is an arc-preserving vertex mapping) from $\vec{G}$ to $\vec{H}_{c}$. In this case, $\vec{G}$ is said to be colored by $\vec{H}_{c}$. Similarly, every homomorphism from $\vec{G}$ to an oriented graph $\vec{H}$ can be viewed as a coloring of $\vec{G}$ using the vertices of $\vec{H}$ as colors. The oriented chromatic number $\vec{\chi}(\vec{G})$ of an oriented graph $\vec{G}$ is the smallest number of colors needed for its oriented coloring. The oriented chromatic number $\vec{\chi}(G)$ of an undirected graph $G$ is the maximal chromatic number over all possible orientations of $G$. The oriented chromatic number of a family of graphs is the maximal chromatic number over all possible graphs of the family.

Oriented coloring has been studied in recent years $[2,5,7,8,9,11,13,14$, $15,17,18]$, see [12] for a survey of the main results. Several authors established or bounded the oriented chromatic number for some families of graphs, such as oriented planar graphs [11], outerplanar graphs [14, 15], graphs with bounded degree three [7, 14, 17], $k$-trees [14], Halin graphs [4, 9], graphs with given excess [8] or grids [5, 18].

For a pair of undirected graphs $G$ and $H$, the Cartesian product $G \square H$ of $G$ and $H$ is the graph with vertex set $V(G) \times V(H)$ and where two vertices are adjacent if and only if they are equal in one coordinate and adjacent in the other. The strong product $G \boxtimes H$ of graphs $G$ and $H$ is the graph with vertex set $V(G) \times V(H)$ and where two vertices are adjacent if and only if they are adjacent in one coordinate and adjacent or equal in the other. We use $P_{k}$ to denote the path on $k$ vertices. In this paper we focus on the oriented chromatic number of Cartesian products of paths, called grids, and strong products of paths, called strong-grids.

In [5], Fertin, Raspaud and Roychowdhury have discussed bounds for the oriented chromatic number of $P_{m} \square P_{n}$. They showed that

- $\vec{\chi}\left(P_{m} \square P_{n}\right) \leq 11$, for every $m, n \geq 1$,
- there exists an orientation of $P_{4} \square P_{5}$ that requires 7 colors,
- $\vec{\chi}\left(P_{2} \square P_{2}\right)=4, \vec{\chi}\left(P_{2} \square P_{3}\right)=5$ and $\vec{\chi}\left(P_{2} \square P_{n}\right)=6$, for $n \geq 4$,
- $\vec{\chi}\left(P_{3} \square P_{3}\right)=\vec{\chi}\left(P_{3} \square P_{4}\right)=\vec{\chi}\left(P_{3} \square P_{5}\right)=6$, and $6 \leq \vec{\chi}\left(P_{3} \square P_{n}\right) \leq 7$, for every $n \geq 6$,
- $\vec{\chi}\left(P_{4} \square P_{4}\right)=6$.

They also formulated the two following conjectures:

- every orientation of $P_{m} \square P_{n}$ can be colored by seven colors,
- every orientation of $P_{m} \square P_{n}$ can be colored by $\vec{T}_{7}$.

The coloring graph $\vec{T}_{7}$ is an oriented graph with set of vertices $V\left(\vec{T}_{7}\right)=$ $\{0,1,2, \ldots, 6\}$ and set of $\operatorname{arcs} A\left(\vec{T}_{7}\right)=\left\{(x, x+b(\bmod 7)): x \in V\left(\vec{T}_{7}\right), b=\right.$ 1,2 , or 4\}. Szepietowski and Targan [18] disproved the second conjecture by exhibiting an orientation of $P_{5} \square P_{33}$ that cannot be colored by $\vec{T}_{7}$. By the way, the oriented graph found in [18] can be colored by another coloring graph with 7 vertices. Furthermore, they showed that

- $\vec{\chi}\left(P_{4} \square P_{n}\right)=7$, for every $n \geq 5$,
- $\vec{\chi}\left(P_{3} \square P_{6}\right)=6$,
- $\vec{\chi}\left(P_{3} \square P_{n}\right)=7$, for every $n \geq 7$.

Dybizbański and Nenca [3] disproved the first conjecture by exhibiting an orientation of $P_{7} \square P_{212}$ that requires 8 colors. However, the bounds for $P_{5} \square P_{n}$ were still $7 \leq \vec{\chi}\left(P_{5} \square P_{n}\right) \leq 11$.

Aravind, Narayanan and Subramanian [1] discussed the oriented chromatic number of strong products of paths. They showed that

- $8 \leq \vec{\chi}\left(P_{2} \boxtimes P_{n}\right) \leq 11$, for every $n \geq 5$,
- $10 \leq \vec{\chi}\left(P_{3} \boxtimes P_{n}\right) \leq 67$, for every $n \geq 5$.

Sopena [16] proved that

- $\vec{\chi}\left(P_{k} \boxtimes P_{n}\right) \leq 126$, for every $n, k \geq 3$.

The rest of the paper is organized as follows. In Section 2, we give definitions. In Section 3, we prove the following theorem.

Theorem 1.1. For every $n \geq 5$, $\vec{\chi}\left(P_{5} \square P_{n}\right) \leq 9$.
Moreover, we prove that the family of all orientations of all grids $P_{5} \square P_{n}$ can be colored with one coloring graph of order 9 , which we call $\vec{H}_{9}$.

In Section 4, we show that there exists an orientation of $P_{2} \boxtimes P_{398}$ which cannot be colored by any coloring graph with 10 vertices. This means that the
lower bound for oriented chromatic number of strong-grids with two columns is the same as the upper bound obtained by Aravind, Narayanan and Subramanian in [1]. Moreover, we improve the bounds for the families of all orientations of the following strong-grids: strong-grids with three columns, denoted by $\mathcal{S}_{3} \rightarrow$, and strong-grids with four columns denoted by $\mathcal{S}_{4}$. We show that for every $n$, every orientation of $P_{3} \boxtimes P_{n}$ can be colored by 19 colors and that there exists an $n$ and an orientation of $P_{3} \boxtimes P_{n}$ which requires 11 colors for oriented coloring. It follows that

Theorem 1.2. $11 \leq \vec{\chi}\left(\mathcal{S}_{3}\right) \leq 19$.
Moreover, we prove that for every $n$, every orientation of $P_{4} \boxtimes P_{n}$ can be colored by 43 colors and that there exists an $n$ and an orientation of $P_{4} \boxtimes P_{n}$ that requires 11 colors for oriented coloring. This means that
Theorem 1.3. $11 \leq \vec{\chi}\left(\mathcal{S}_{4}\right) \leq 43$.

## 2. Definitions

Definition 2.1. Let $G_{1}\left(V_{1}, E_{1}\right)$ and $G_{2}\left(V_{2}, E_{2}\right)$ be two undirected graphs. Their Cartesian product, denoted by $G_{1} \square G_{2}$, is the graph where $V\left(G_{1} \square G_{2}\right)=V_{1} \times V_{2}$ and $\left(u_{1}, u_{2}\right),\left(v_{1}, v_{2}\right) \in E\left(G_{1} \square G_{2}\right)$ if either $u_{1}=v_{1}$ and $\left(u_{2}, v_{2}\right) \in E_{2}$, or $u_{2}=v_{2}$ and $\left(u_{1}, v_{1}\right) \in E_{1}$. Their strong product, denoted by $G_{1} \boxtimes G_{2}$, is the graph where $V\left(G_{1} \boxtimes G_{2}\right)=V_{1} \times V_{2}$ and $\left(u_{1}, u_{2}\right),\left(v_{1}, v_{2}\right) \in E\left(G_{1} \boxtimes G_{2}\right)$ if either $u_{1}=v_{1}$ and $\left(u_{2}, v_{2}\right) \in E_{2}$, or $u_{2}=v_{2}$ and $\left(u_{1}, v_{1}\right) \in E_{1}$, or $\left(u_{1}, v_{1}\right) \in E_{1}$ and $\left(u_{2}, v_{2}\right) \in E_{2}$.

The Cartesian product of two paths $P_{m} \square P_{n}$ is called the $m \times n$ grid. The strong product of two paths $P_{m} \boxtimes P_{n}$ is called the $m \times n$ strong-grid.

We shall say that $u \in V$ is a source (respectively a sink) if there is no arc incoming to $u$ (respectively outgoing from $u$ ). A tournament is an orientation of an undirected complete graph.

Let $p$ be a prime number such that $p \equiv 3(\bmod 4), d=\frac{p-1}{2}$, and $c_{1}, c_{2}, \ldots, c_{d}$ be the non-zero quadratic residues of $p$. The directed graph $\vec{T}_{p}$ with set of vertices $V\left(\vec{T}_{p}\right)=\{0,1, \ldots, p-1\}$ and set of $\operatorname{arcs} A\left(\vec{T}_{p}\right)=\left\{\left(x, x+c_{i}(\bmod p)\right)\right.$ : $\left.x \in V\left(\vec{T}_{p}\right), 1 \leq i \leq d\right\}$ is called the Paley tournament of order $p$. It is easy to check that Paley tournaments are arc-transitive, i.e. for any two arcs $(u, v)$, $(x, y)$ in $A\left(\vec{T}_{p}\right)$, there exists an automorphism $f: \vec{T}_{p} \rightarrow \vec{T}_{p}$ satisfying $f(u)=x$, $f(v)=y$, and self-converse, i.e. $T_{p}$ is isomorphic to its converse (a graph obtained by reversing each arc), see [6].

## 3. Cartesian Product of Paths

In this section we prove Theorem 1.1. First, we define the coloring graph $\vec{H}_{9}$ which is obtained from $\vec{T}_{7}$ by adding two vertices, one sink and one source.

More precisely, $V\left(\vec{H}_{9}\right)=\{0,1,2,3,4,5,6,7,8\}$ and $(u, v) \in A\left(\vec{H}_{9}\right)$ if

- $u, v<7$ and $v-u \equiv 1,2$, or $4(\bmod 7)$, or
- $u=7$, or
- $v=8$.

Consider the grid $P_{5} \square P_{n}$ with five columns and $n$ rows. Let us denote by $v_{i}=(1, i), w_{i}=(2, i), x_{i}=(3, i), y_{i}=(4, i), z_{i}=(5, i)$ the five vertices in the $i$-th row, see Figure 1 .


Figure 1. The grid $G(5, n)$.
Suppose that $\vec{G}$ is an orientation of $P_{5} \square P_{n}$. By $S(\vec{G})$ we denote the set of reachable colorings of the last row of $\vec{G}$, namely $S(\vec{G})=\left\{\left(c_{1}, c_{2}, c_{3}, c_{4}, c_{5}\right)\right.$ : there exists a coloring $\gamma: V(\vec{G}) \rightarrow \vec{T}_{7}$, such that $\gamma\left(v_{n}\right)=c_{1}, \gamma\left(w_{n}\right)=c_{2}$, $\left.\gamma\left(x_{n}\right)=c_{3}, \gamma\left(y_{n}\right)=c_{4}, \gamma\left(z_{n}\right)=c_{5}\right\}$.

Let $\Phi(\vec{G})$ be the vector $\left(\delta\left(v_{n}\right), \delta\left(x_{n}\right), \delta\left(z_{n}\right)\right)$, where

$$
\delta(u)=\left\{\begin{array}{lc}
0, & \text { if } u \text { is a sink or a source }, \\
1, & \text { otherwise }
\end{array}\right.
$$

Let $\mathcal{G} \rightarrow$ denote the family of all orientations of all grids with five columns. Let us denote by $\mathcal{G}^{\prime \rightarrow}$ the set of orientation in $\mathcal{G} \rightarrow$ without sinks or sources on first, third and last columns above the last row, i.e., $\mathcal{G}^{\prime \rightarrow}=\bigcup_{n=1}^{\infty}\left\{\vec{G} \in \mathcal{G}^{\rightarrow}\right.$ : $\vec{G}$ is an orientation of $P_{5} \square P_{n}$, where for every $1 \leq i<n$, none of $v_{i}, x_{i}, z_{i}$ is a sink or a source $\}$. First, we show that every $\vec{G} \in \mathcal{G}^{\prime \rightarrow}$ can be colored by the Paley tournament $\vec{T}_{7}$.
Lemma 3.1. For every $\vec{G} \in \mathcal{G}^{\prime} \rightarrow$, there exists a homomorphism $\gamma: \vec{G} \rightarrow \vec{T}_{7}$.
Proof. We use an algorithm to check that the family $\left\{S(\vec{G}): \vec{G} \in \mathcal{G}^{\prime \rightarrow}\right\}$ does not contain the empty set, which means that every $\vec{G} \in \mathcal{G}^{\prime \rightarrow}$ can be colored by $\vec{T}_{7}$. In order to do this, the algorithm looks through all possible pairs $(S, \Phi)$ such
that there exists $\vec{G} \in \mathcal{G}^{\prime} \rightarrow$ for which $S(\vec{G})=S$ and $\Phi(\vec{G})=\Phi$. The algorithm uses a queue to stores such pairs. Since $\vec{T}_{7}$ is arc-transitive and self-converse, we can consider only those orientations of $\vec{G}$ where $\left(v_{n}, w_{n}\right) \in A(\vec{G})$ and only those colorings where $c\left(v_{n}\right)=0$ and $c\left(w_{n}\right)=1$ (see [18] for more details). Note that there are eight possible orientations of the last row with $\left(v_{n}, w_{n}\right) \in A(\vec{G})$ and $3 \times 3 \times 3=27$ possible colorings of the last row satysfying $c\left(v_{n}\right)=0$ and $c\left(w_{n}\right)=1$.

The algorithm starts with the grid $P_{5} \square P_{1}$. For every orientation $\vec{F}$ of $P_{5} \square P_{1}$, it computes the pair $(S(\vec{F}), \Phi(\vec{F}))$ and inserts it to the queue $Q$. Next, the algorithm takes one by one a pair $(S, \Phi)$ from the queue, which contains the information for some $\vec{G} \in \mathcal{G}^{\prime \rightarrow}$ and considers every graph $\vec{R} \in \mathcal{G}^{\prime \rightarrow}$ which can be built from $\vec{G}$ by adding one extra row and which does not contain a sink or a source in the first, the third and the fifth column above the last row. For each such graph $\vec{R}$, the algorithm computes the pair $(S(\vec{R}), \Phi(\vec{R}))$ and, provided it is a new one, inserts it to the queue $Q$. In order to compute the pair $(S(\vec{R}), \Phi(\vec{R}))$, the algorithm does not need to reconstruct the whole graph $\vec{G}$ in its memory and build the graph $\vec{R}$.

After running this algorithm, we have found that the algorithm stops with empty queue $Q$ and no pair of the form $(\emptyset, \Phi)$ is reachable. This means that every $\vec{G} \in \mathcal{G}^{\prime \rightarrow}$ can be colored with $\vec{T}_{7}$.

Proof of Theorem 1.1. Let $\vec{G}$ be an orientation of $P_{5} \square P_{n}, n \geq 5$. We show that there exists a homomorphism $\gamma$ such that $\gamma: \vec{G} \rightarrow \vec{H}_{9}$. First, we construct a new orientation $\vec{G}^{\prime}$ of $P_{5} \square P_{n}$ by reversing some arcs in $\vec{G}$. More precisely, for every row $i$ we have

- if $v_{i}$ is a sink or a source, then we reverse the arc between $v_{i}$ and $w_{i}$,
- if $x_{i}$ is a sink or a source, then we reverse the arc between $w_{i}$ and $x_{i}$,
- if $z_{i}$ is a sink or a source, then we reverse the arc between $y_{i}$ and $z_{i}$.

Note that every vertical arc remains unchanged. Moreover, for each vertex $u$ in the first, the third or the fifth column, we reverse at most one arc incident with $u$. The other end of each reversed arc is in the second or fourth column thus the changes of orientation cannot create a new sink or a source in the first, the third or the fifth column. The graph $\vec{G}^{\prime}$ does not have any sources or sinks in the first, third and fifth column. Hence, by Lemma 3.1, there is a coloring $\gamma^{\prime}: \vec{G}^{\prime} \rightarrow \vec{T}_{7}$. We construct the coloring $\gamma: \vec{G} \rightarrow \vec{H}_{9}$ in the following way.

- If $v$ is in the second or the fourth column, then $\gamma(v):=\gamma^{\prime}(v)$,
- if $v$ is in the first, the third or the fifth column, and is not a sink or a source in $\vec{G}$, then $\gamma(v):=\gamma^{\prime}(v)$,
- if $v$ is in the first, the third or the fifth column, and is a sink or a source in $\vec{G}$, then
* $\gamma(v):=7$ if $v$ is a source in $\vec{G}$,
* $\gamma(v):=8$ if $v$ is a sink in $\vec{G}$.

In order to show that $\gamma$ is an oriented coloring of $\vec{G}$, consider an arc $u v \in \vec{G}$. There are four possible cases.

- If the arc between $u$ and $v$ has been reversed, then one of its ends, say $u$, is a source (or a sink, respectively) in the first, the third, or the fifth column and receives color 7 (or 8 , respectively). The other end $v$ of the arc is in the second or the fourth column, thus it has a color from $\vec{T}_{7}$ and $(7, \gamma(v)) \in A\left(\vec{H}_{9}\right)$ (or $(\gamma(v), 8) \in A\left(\vec{H}_{9}\right)$, respectively).
- If the arc between $u$ and $v$ has not been reversed and the colors of $u$ and $v$ have not been changed, then these colors fit in $\vec{T}_{7}$, and they also fit in $\vec{H}_{9}$.
- If the arc between $u$ and $v$ has not been reversed but the color of one of its ends, say $u$, was changed, then it means that the vertex $u$ is a sink or a source and has a color that matches the color of the vertex $v\left(\gamma(v) \in \vec{T}_{7}\right)$.
- If the arc between $u$ and $v$ has not been reversed but the colors of both its ends were changed, then $u$ and $v$ belong to the same column (the first, the third or the fifth), and one of them is a sink and the other is a source. Their colors are 7 and 8 , and fit in $\vec{H}_{9}$.


## 4. Strong Product of Paths

In this section, we focus on the strong products of paths $S_{k, n}=P_{k} \boxtimes P_{n}$, called strong-grids. Consider the strong-grid $S_{2, n}$ with 2 columns and $n$ rows. Let us denote by $x_{i}=(1, i)$ and $y_{i}=(2, i)$ the two vertices in the $i$-th row.

Let $\vec{S}_{2, n}$ denote the set of orientations of the strong-grid $S_{2, n}$. Since every orientation of a strong-grid $S_{2, n}$ is isomorphic to another orientation with all horizontal edges going in the same direction, we consider only the later ones (see Figure 2), i.e., $\vec{S}_{2, n}^{\prime}=\left\{\vec{S} \in \vec{S}_{2, n}:\left(x_{i}, y_{i}\right) \in A(\vec{S}), 1 \leq i \leq n\right\}$. Suppose that $\vec{S} \in \vec{S}_{2, n}^{\prime}$. By $T(\vec{S}, \vec{H})$ we denote the set of reachable colorings on the last row of $\vec{S}$ by the coloring graph $\vec{H}$, i.e., $T(\vec{S}, \vec{H})=\left\{\left(c_{1}, c_{2}\right)\right.$ : there exists a coloring $\gamma: V(\vec{S}) \rightarrow \vec{H}$, such that $\quad \gamma\left(x_{n}\right)=c_{1}$ and $\left.\gamma\left(y_{n}\right)=c_{2}\right\}$.

We use an algorithm similar to the one used in Lemma 3.1 to prove the following theorem.


Figure 2. The strong-grid $S_{2, n}^{\prime}$.
Theorem 4.1. There exists an integer $n$ such that $\vec{\chi}\left(P_{2} \boxtimes P_{n}\right)=11$.
Proof. Aravind, Narayanan and Subramanian [1] proved that every strong-grid with two columns can be colored by the Paley tournament $\vec{T}_{11}$. We show that there exists an orientation of $P_{2} \boxtimes P_{398}$ which cannot be colored by any coloring graph with ten vertices.

In order to construct a strong-grid that needs eleven colors for any oriented coloring, we use a Extend function. The Extend function for a given oriented strong-grid $\vec{S}_{1}$ and a coloring graph $\vec{H}$ returns a strong-grid $\vec{S}_{2}$, such that:

- $\vec{S}_{2}$ cannot be colored by $\vec{H}$.
- $\vec{S}_{2}$ is constructed by adding new rows to $\vec{S}_{1}$.
- If $\vec{S}_{1}$ cannot be colored by $\vec{H}$, then $\vec{S}_{2}$ is equal to $\vec{S}_{1}$.

The Extend function for a given $\vec{S}_{1}$ and $\vec{H}$ looks through all possible sets $T(\vec{S}, \vec{H})$, where $\vec{S}$ can be built from $\vec{S}_{1}$ by adding some extra rows. Similarly to the algorithm used in Lemma 3.1, the function uses a queue to stores such sets.

The function starts by computing the set $T\left(\vec{S}_{1}, \vec{H}\right)$ and inserts it to the queue $Q$. Next, the algorithm takes one by one a set $C$ from the queue $Q$ and for every orientation of an extra row computes the set of colorings of the next row denoted by $C^{\prime}$ (similarly to Lemma 3.1). Then the set $C^{\prime}$ is inserted into the queue provided it is new one. Moreover, the function puts to an additional memory the triple consisting of the set $C$, the orientation of an extra row and the set $C^{\prime}$. The algorithm stops when an empty set of colorings is reached. After reaching an empty set of colorings, the procedure reconstructs a grid $\vec{S}_{2}$, such
that $T\left(\vec{S}_{2}, \vec{H}\right)=\emptyset$ and $\vec{S}_{2}$ is an extension of $\vec{S}_{1}$. To do this the algorithm uses the information kept in additional memory.

We use the Extend function with all non-isomorphic coloring graphs on ten vertices (there are 9733056 such tournaments). In the first run, we use a single arc. In the next run, we use the grid $\vec{S}$ returned in the previous run, and so on. It is easy to see that if $\vec{S}$ cannot be colored by $\vec{H}$, then any extension $\vec{S}^{\prime}$ of $\vec{S}$ cannot be colored by $\vec{H}$. The result of the last run is a strong-grid $\vec{S}^{\prime \prime}$ that cannot be colored by any of the coloring graphs with ten vertices. The size of $\vec{S}^{\prime \prime}$ may vary depending on the order in which we consider non-isomorphic coloring graphs. Using nauty [10] to generate the list of all non-isomorphic coloring graphs of order 10, we found an orientation of $P_{2} \boxtimes P_{398}$ that admits no oriented coloring with 10 colors.

### 4.1. Proof of Theorem 1.2.

The lower bound for $\vec{\chi}\left(\mathcal{S}_{3}\right)$ follows from Theorem 4.1. We use an algorithm similar to the one used in Section 3 to show that for any $n \geq 1$ any orientation of $P_{3} \boxtimes P_{n}$ can be colored by the Paley tournament $\vec{T}_{19}$. Consider now the grid $S_{3, n}$ with three columns and $n$ rows. Let us denote by $x_{i}=(1, i), y_{i}=(2, i)$, $z_{i}=(3, i)$ the three vertices in the $i$-th row, see Figure 3.


Figure 3. The strong-grid $S_{3, n}$.
Let $\mathcal{S}_{3}$ denote the set of orientations of all strong-grids with three columns. Let $\vec{S}$ be an orientation of $P_{3} \boxtimes P_{n}$, for some $n \geq 1$, and let $T(\vec{S})$ be the set of reachable distinct colorings on the last row of $\vec{S}$ by $\vec{T}_{19}$, i.e., $T(\vec{S})=$ $\left\{\left(c_{1}, c_{2}, c_{3}\right):\right.$ there exists a coloring $\gamma: V(\vec{S}) \rightarrow \vec{T}_{19}$ such that $\gamma\left(x_{n}\right)=c_{1}$, $\left.\gamma\left(y_{n}\right)=c_{2}, \gamma\left(z_{n}\right)=c_{3}\right\}$. Once again, since $\vec{T}_{19}$ is arc-transitive and self-converse, we can consider only those orientations of $\vec{S}$ where $\left(x_{n}, y_{n}\right) \in A(\vec{S})$, and only
those colorings where $c\left(x_{n}\right)=0$ and $c\left(y_{n}\right)=1$. Thus, there are only two orientations of the last row with $\left(x_{n}, y_{n}\right) \in \bar{S}$ and only nine colorings of the last row satisfying $c\left(x_{n}\right)=0$ and $c\left(y_{n}\right)=1$.

The algorithm starts with the strong-grid $P_{3} \boxtimes P_{1}$. For every orientation $\vec{S}$ of $P_{3} \boxtimes P_{1}$ it computes the set $T(\vec{S})$ and inserts it into the queue $Q$. Next, the algorithm takes one by one a set $T$ from the queue. The set $T$ is the set of colorings of the last row of some $\vec{S} \in \mathcal{S}_{3}$. The algorithm considers every graph $\vec{R} \in \mathcal{S}_{3}^{\overrightarrow{ }}$ which can be built from $\vec{S}$ by adding one extra row. For each such a graph $\vec{R}$ the algorithm computes the set $T(\vec{R})$ and, provided it is new one, inserts it into the queue $Q$. Once again the algorithm does not need to reconstruct the whole graph $\vec{R}$ to compute the set $T(\vec{R})$. After running the algorithm, we have found that the algorithm stops with the empty queue $Q$ and the empty set of colorings of the last row is not reached. This means that for every $n$, every orientation of $P_{3} \boxtimes P_{n}$ can be colored by the Paley tournament $\vec{T}_{19}$.

### 4.2. Proof of Theorem 1.3.

Consider now the strong-grid $S_{4,2}=P_{4} \boxtimes P_{2}$. Let us denote by $v=(1,1)$, $w=(2,1), x=(3,1), y=(4,1)$ the vertices of the first row of $S_{4,2}$ and by $v^{\prime}=(1,2), w^{\prime}=(2,2), x^{\prime}=(3,2), y^{\prime}=(4,2)$ the vertices of second row, see Figure 4.


Figure 4. The strong-grid $S_{4,2}$.
The lower bound for $\vec{\chi}\left(\mathcal{S}_{4}\right)$ follows from Theorem 4.1. To prove the upper bound we shall use the following property of the Paley tournament $\vec{T}_{43}$.

Lemma 4.2. For any orientation $\vec{S}$ of $S_{4,2}$, if the first row of $\vec{S}$ can be colored by $\vec{T}_{43}$ with colors $\left(c_{v}, c_{w}, c_{x}, c_{y}\right)$, where $c_{v} \neq c_{x}$ and $c_{w} \neq c_{y}$, then there is a coloring $\gamma: \vec{S} \rightarrow \vec{T}_{43}$, such that $\gamma(v)=c_{v}, \gamma(w)=c_{w}, \gamma(x)=c_{x}, \gamma(y)=c_{y}$, and moreover $\gamma\left(v^{\prime}\right) \neq \gamma\left(x^{\prime}\right)$ and $\gamma\left(w^{\prime}\right) \neq \gamma\left(y^{\prime}\right)$.

Proof. We use a computer algorithm to check the above property. Once again, since $\vec{T}_{43}$ is arc-transitive and self-converse, we can consider only those orientations of $S_{4,2}$ where $(u, v) \in A\left(S_{4,2}\right)$ and only those colorings where $c_{u}=0$ and $c_{v}=1$.

Let $\vec{S}$ be any orientation of $P_{4} \boxtimes P_{n}$. Using Lemma 4.2, we can color $\vec{S}$ row by row.


Figure 5. The oriented grid $\overrightarrow{G_{5}}$.

## 5. Conclusion

In this paper, we have proved that every orientation of $P_{5} \square P_{n}$ can be colored by the coloring graph $\vec{H}_{9}$, see Section 3. However, $\vec{H}_{9}$ cannot be used to color every orientation of every grid $P_{m} \square P_{n}$. To prove that, we use an algorithm similar to the algorithm used in Lemma 3.1 and construct an orientation $\vec{G}_{5}$ of the grid
$P_{5} \square P_{28}$ (see Figure 5) without sink and source vertices on the second, third and fourth column, which cannot be colored by the Paley tournament $\vec{T}_{7}$. One can now easily construct an orientation of $P_{7} \square P_{28}$, by adding two extra columns in $G_{5}$ : before the first one - column 0 , and after the fifth one - column 6 . The arcs incident with all vertices in additional columns are oriented in such a way that there is no sink and source vertices in the first and the fifth column. The resulting orientation cannot be colored by $\vec{H}_{9}$, as we are not able to use color 7 or 8 to color any vertex from column $1-5$. We can only use colors from $\vec{T}_{7}$ but the oriented grid $G_{5}$ cannot be colored by $\vec{T}_{7}$.

On the website https://inf.ug.edu.pl/grids/ we posted the grid $\overrightarrow{G_{5}}$, and the strong-grid $P_{2} \boxtimes P_{398}$ that requires 11 colors. On the same site we posted sample C++ programs that can be used to verify those grids and programs that we have used to prove property of $\vec{T}_{43}$ (Lemma 4.2), $\vec{T}_{19}$ (Theorem 1.2), and $\vec{T}_{7}$ (Lemma 3.1).

Every time when we use Paley tournaments to color certain kind of grids we check that no Paley tournament of smaller order has the expected property. For example $\vec{T}_{43}$ is the smallest Paley tournament with the property described in Lemma 4.2. The fact that this property is not true for $\vec{T}_{31}$ does not mean that any orientation of $P_{4} \boxtimes P_{n}$ cannot be colored by $\vec{T}_{31}$.

## Acknowledgments

The authors would like to thank the anonymous reviewers for helpful comments that contributed to improving the final version of the paper.

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