Discussiones Mathematicae Graph Theory 38 (2018) 323–349 doi:10.7151/dmgt.2041

# $\mathcal{P}$ -APEX GRAPHS

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Dedicated to the memory of Professor Horst Sachs (1927 – 2017)

#### Abstract

Let  $\mathcal{P}$  be an arbitrary class of graphs that is closed under taking induced subgraphs and let  $\mathcal{C}(\mathcal{P})$  be the family of forbidden subgraphs for  $\mathcal{P}$ . We investigate the class  $\mathcal{P}(k)$  consisting of all the graphs G for which the removal of no more than k vertices results in graphs that belong to  $\mathcal{P}$ . This approach provides an analogy to apex graphs and apex-outerplanar graphs studied previously. We give a sharp upper bound on the number of vertices of graphs in  $\mathcal{C}(\mathcal{P}(1))$  and we give a construction of graphs in  $\mathcal{C}(\mathcal{P}(k))$  of relatively large order for  $k \geq 2$ . This construction implies a lower bound on the maximum order of graphs in  $\mathcal{C}(\mathcal{P}(k))$ . Especially, we investigate  $\mathcal{C}(\mathcal{W}_r(1))$ , where  $\mathcal{W}_r$  denotes the class of  $P_r$ -free graphs. We determine some forbidden subgraphs for the class  $\mathcal{W}_r(1)$  with the minimum and maximum number of vertices. Moreover, we give sufficient conditions for graphs belonging to  $\mathcal{C}(\mathcal{P}(k))$ , where  $\mathcal{P}$  is an additive class, and a characterisation of all forests in  $\mathcal{C}(\mathcal{P}(k))$ . Particularly we deal with  $\mathcal{C}(\mathcal{P}(1))$ , where  $\mathcal{P}$  is a class closed under substitution and obtain a characterisation of all graphs in the corresponding  $\mathcal{C}(\mathcal{P}(1))$ . In order to obtain desired results we exploit some hypergraph tools and this technique gives a new result in the hypergraph theory.

**Keywords:** induced hereditary classes of graphs, forbidden subgraphs, hypergraphs, transversal number.

2010 Mathematics Subject Classification: 05C75, 05C15.

#### 1. INTRODUCTION

We only consider finite and simple graphs and follow [1] for graph-theoretical terminology and notation not defined here. A graph G is an *apex graph* if it contains a vertex w such that G - w is planar. Although apex graphs seem to be close to planar graphs, some of their properties are far from corresponding properties of planar graphs (for example, see [18]).

A result of Robertson and Seymour (see [19]) says that every proper minorclosed class of graphs  $\mathcal{P}$  can be characterized by a finite family of *forbidden minors* (minor-minimal graphs not in  $\mathcal{P}$ ). Evidently, the class of apex graphs is minor-closed but the long-standing problem of finding the complete family of forbidden minors for this class is still open.

However, Dziobak in [9] introduced an apex-outerplanar graph that is a conceptual analogue to an apex graph. Namely, a graph G is *apex-outerplanar* if there exists  $w \in V(G)$  such that G - w is outerplanar. Moreover, Dziobak provided the complete list of 57 forbidden minors for this class.

Another attempt to extend the concept of an apex graph is presented in [20] where an l-apex graph is defined. A graph G is an l-apex graph if it can be made planar by removing at most l vertices.

This paper concerns classes of graphs that generalize the aforementioned. Formally, by a *class of graphs* we mean an arbitrary family of non-isomorphic graphs. The empty class of graphs and the class of all graphs are called *trivial*. A class of graphs  $\mathcal{P}$  is *induced hereditary* if it is closed with respect to taking induced subgraphs. Such a class  $\mathcal{P}$  can be uniquely characterized by the family of *forbidden subgraphs*  $\mathcal{C}(\mathcal{P})$  that is defined as a set

 $\{G: G \notin \mathcal{P} \text{ and } H \in \mathcal{P} \text{ for each proper induced subgraph } H \text{ of } G\}.$ 

By  $\mathbf{L}_{\leq}$  we denote the class of all non-trivial induced hereditary classes of graphs. Each class  $\mathcal{P} \in \mathbf{L}_{\leq}$  has a non-empty family of forbidden subgraphs, consisting of graphs with at least two vertices. Moreover,  $\mathcal{C}(\mathcal{P})$  contains only connected graphs when  $\mathcal{P}$  is *additive*, i.e., closed under taking the union of disjoint graphs. By  $\mathbf{L}_{\leq}^{a}$  we denote the family of all non-trivial induced hereditary and additive classes of graphs.

Let  $\mathcal{P} \in \mathbf{L}_{\leq}$  and let k be a non-negative integer. A graph G is a  $\mathcal{P}(k)$ -apex graph if there is  $W \subseteq V(G)$ ,  $|W| \leq k$  (W is allowed to be the empty set), such that G - W belongs to  $\mathcal{P}$ . We denote the set of all  $\mathcal{P}(k)$ -apex graphs by  $\mathcal{P}(k)$  for short.

We can see immediately that if k is a non-negative integer and  $\mathcal{P} \in \mathbf{L}_{\leq}$ , then  $\mathcal{P}(k) \in \mathbf{L}_{\leq}$  too. On the other hand, the additivity of  $\mathcal{P} \in \mathbf{L}_{\leq}$  implies the additivity of  $\mathcal{P}(k)$  if and only if k = 0. Indeed,  $\mathcal{P}(0) = \mathcal{P}$ . Moreover, if  $\mathcal{P} \in \mathbf{L}_{\leq}^{a}$ , then  $\mathcal{C}(\mathcal{P}) \neq \emptyset$  and assuming that  $F \in \mathcal{C}(\mathcal{P})$  we can easily see that the union of

## $\mathcal{P} ext{-}\operatorname{Apex}\ \operatorname{Graphs}$

k + 1 disjoint copies of F is in  $\mathcal{C}(\mathcal{P}(k))$ . Thus, for  $k \ge 1$ , it yields the existence of at least one disconnected graph that is forbidden for  $\mathcal{P}(k)$ . Hence, for  $k \ge 1$ , the class  $\mathcal{P}(k)$  is not additive.

Lewis and Yannakakis in [17] have shown that for any non-trivial induced hereditary class  $\mathcal{P}$  containing infinitely many graphs and for a given positive integer k, the decision problem: "does G belong to  $\mathcal{P}(k)$ ?" is NP-complete.

In this paper, we investigate the classes  $\mathcal{P}(k)$ , in particular we focus on forbidden subgraphs for the classes  $\mathcal{P}(k)$  (i.e., we study graphs in  $\mathcal{C}(\mathcal{P}(k))$ ). Additionally, we use hypergraphs as an effective tool in the research on  $\mathcal{P}(k)$ .

Let  $\mathcal{H}$  be a hypergraph with vertex set  $V(\mathcal{H})$  and edge set  $\mathcal{E}(\mathcal{H})$  and let  $W \subseteq V(\mathcal{H})$ . The hypergraph  $\mathcal{H}[W]$  induced in  $\mathcal{H}$  by W has vertex set W and edge set  $\{E \in \mathcal{E}(\mathcal{H}) : E \subseteq W\}$ . To simplify the notation we write  $\mathcal{H} - W$  instead of  $\mathcal{H}[V(\mathcal{H}) \setminus W]$  and, moreover,  $\mathcal{H} - v$  instead of  $\mathcal{H} - \{v\}$  when v is a vertex of  $\mathcal{H}$ . Analogously, we write  $\mathcal{H} - E$  to denote the hypergraph obtained from  $\mathcal{H}$  by the deletion of the edge E from  $\mathcal{E}(\mathcal{H})$ .

By  $\mathcal{H}_1 \cup \mathcal{H}_2$  we mean the union of disjoint hypergraphs  $\mathcal{H}_1$  and  $\mathcal{H}_2$ , i.e., the hypergraph with vertex set  $V(\mathcal{H}_1) \cup V(\mathcal{H}_2)$  and edge set  $\mathcal{E}(\mathcal{H}_1) \cup \mathcal{E}(\mathcal{H}_2)$ . Moreover, notations  $2\mathcal{H}_1$ ,  $\mathcal{H}_1 \cup \mathcal{H}_1$ , and their generalization are used interchangeably. The symbol  $\mathcal{H}_1 \leq \mathcal{H}_2$  denotes that the hypergraph  $\mathcal{H}_1$  is isomorphic to a subhypergraph of  $\mathcal{H}_2$  induced by some of its vertex subset. Let r be a non-negative integer. A hypergraph  $\mathcal{H}$  is r-uniform if each edge in  $\mathcal{E}(\mathcal{H})$  has exactly r vertices. A set  $T \subseteq V(\mathcal{H})$  is called a *transversal* of the hypergraph  $\mathcal{H}$  if  $T \cap E \neq \emptyset$  for each  $E \in \mathcal{E}(\mathcal{H})$ . By  $\tau(\mathcal{H})$  we denote the cardinality of the minimum transversal of  $\mathcal{H}$ , i.e.,

$$\tau(\mathcal{H}) = \min\{|T|: T \text{ is a transversal of } \mathcal{H}\}.$$

A hypergraph  $\mathcal{H}$  is  $\tau$ -vertex critical if for any  $v \in V(\mathcal{H})$  the inequality  $\tau(\mathcal{H}-v) \leq \tau(\mathcal{H}) - 1$  holds. If a  $\tau$ -vertex critical hypergraph  $\mathcal{H}$  satisfies  $\tau(\mathcal{H}) = l$  for some positive integer l, then we call it  $\tau$ -vertex l-critical.

Recall that each graph is a hypergraph, which allows us to use these notations also for graphs. The symbols  $K_n$ ,  $P_n$ ,  $C_n$  are used only for graphs and denote the complete graph, the path and the cycle with n vertices, respectively.

This paper is organized as follows. We start with  $\tau$ -vertex *l*-critical hypergraphs in Section 2. We prove an upper bound on the order of a  $\tau$ -vertex 2-critical hypergraph and describe the construction of  $\tau$ -vertex *l*-critical hypergraphs with large number of vertices. Next, in Section 3, we prove some results on relations between  $\tau$ -vertex (k+1)-critical hypergraphs and graphs in  $\mathcal{C}(\mathcal{P}(k))$  for  $\mathcal{P} \in \mathbf{L}_{\leq}$ . In Section 4, for  $\mathcal{P} \in \mathbf{L}_{\leq}^a$  we show some sufficient conditions that have to be satisfied by a graph to be in  $\mathcal{C}(\mathcal{P}(k))$  and we characterize all forests in  $\mathcal{C}(\mathcal{P}(k))$ . Section 5 deals with the class  $\mathcal{P}$  of graphs that does not contain  $P_r$  as an induced subgraph. We determine some forbidden subgraphs for  $\mathcal{P}(1)$  with minimum and maximum order in this case. In Section 6 we characterize all graphs in  $\mathcal{C}(\mathcal{P}(1))$ , where  $\mathcal{P}$  is a class of graphs that is induced hereditary and closed under substitution (for the definition see Section 6).

## 2. $\tau$ -Vertex Critical Hypergraphs

A hypergraph  $\mathcal{H}$  is  $\tau$ -edge *l*-critical if  $\tau(\mathcal{H}) = l$  and the deletion of an edge decreases the transversal number of the resulting hypergraph. It is clear that the class of  $\tau$ -edge *l*-critical hypergraphs without isolated vertices forms a subclass of the class of  $\tau$ -vertex *l*-critical hypergraphs. On the other hand, it is easy to prove that the maximum order of hypergraphs in both classes is the same. In this section we prove that an *r*-uniform  $\tau$ -vertex 2-critical hypergraph has at most  $\left\lfloor \frac{(r+2)^2}{4} \right\rfloor$  vertices. Our proof is different than Tuza's proof in [21] concerning a corresponding theorem for *r*-uniform  $\tau$ -edge 2-critical hypergraphs.

Next, for  $l \geq 3$  we give the construction of an *r*-uniform  $\tau$ -vertex *l*-critical hypergraph with a large order. Gyárfás *et al.* [15] proved that each *r*-uniform  $\tau$ -vertex *l*-critical hypergraph has order bounded from above by  $\binom{l+r-2}{r-2}l+l^{r-1}$ . This bound is probably far from the exact value of the maximum number of vertices in a hypergraph that is *r*-uniform  $\tau$ -vertex *l*-critical. Our construction gives a large lower bound on the maximum order of a hypergraph that is *r*-uniform  $\tau$ -vertex *l*-critical.

**Theorem 1.** Let r be an integer,  $r \geq 2$ , and let  $\mathcal{H}$  be a  $\tau$ -vertex 2-critical hypergraph. If for each  $E \in \mathcal{E}(\mathcal{H})$  we have  $|E| \leq r$ , then

$$|V(\mathcal{H})| \le \left\lfloor \frac{(r+2)^2}{4} \right\rfloor.$$

Moreover, the bound is sharp.

**Proof.** Denote by  $\mathcal{H}'$  a hypergraph obtained from  $\mathcal{H}$  by the optional deletion of some edges in such a way that  $\tau(\mathcal{H}) = \tau(\mathcal{H}') = 2$  and  $\tau(\mathcal{H}' - E') \leq 1$  for each edge E' of  $\mathcal{H}'$ . Let  $\mathcal{E}' = \mathcal{E}(\mathcal{H}')$  and assume  $\mathcal{E}' = \{E'_1, \ldots, E'_m\}$ . Observe that each vertex of  $\mathcal{H}'$  is contained in at least one of the edges in  $\mathcal{E}(\mathcal{H}')$ . Otherwise, if there is  $x \in V(\mathcal{H}')$  such that x belongs to no edge in  $\mathcal{E}(\mathcal{H}')$ , then  $\tau(\mathcal{H} - x) = 2$  giving a contradiction to the  $\tau$ -vertex criticality of  $\mathcal{H}$ .

Let a bipartite graph B be the incidence graph of the hypergraph  $\mathcal{H}'$ . Thus  $B = (V(\mathcal{H}), \mathcal{E}'; E(B))$ , where  $vE' \in E(B)$  if and only if  $v \in E'$ . The previous consideration says that  $d_B(v) \geq 1$  for all  $v \in V(\mathcal{H})$  and  $d_B(E'_i) \leq r$  for all  $i \in \{1, \ldots, m\}$ . The last condition implies  $|E(B)| \leq mr$ .

**Claim 2.** For every  $E'_i$  there is a vertex, say  $v_i \in V(\mathcal{H}) \subseteq V(B)$ , such that  $v_i \notin E'_i$  but  $v_i \in E'_j \in \mathcal{E}'$  for all  $j \neq i$ .

**Proof.** Delete a vertex  $E'_i$  from the graph B. The graph  $B - E'_i$  is an incidence graph of the hypergraph  $\mathcal{H}' - E'_i$ , so  $\tau(\mathcal{H}' - E'_i) = 1$ , i.e., there is a vertex, say x, which is adjacent in B to every  $E'_j$ ,  $j \neq i$ . Obviously the vertex x is not adjacent to  $E'_i$ , otherwise in the hypergraph  $\mathcal{H}'$  there would be a 1-element transversal  $\{x\}$ , which is impossible. Thus x can play the role of  $v_i$  from the statement.  $\Box$ 

By Claim 2, in the graph *B* there is a set of *m* vertices  $\{v_1, \ldots, v_m\}$  with  $d_B(v_i) = m - 1$ , for  $i \in \{1, \ldots, m\}$ . Since  $d_B(v) \ge 1$  for each  $v \in V(\mathcal{H})$  we have  $m(m-1) + (n-m) \le |E(B)| \le mr$ , where  $n = |V(\mathcal{H})|$ . It leads to the inequality  $n \le -m^2 + (r+2)m$ . Thus for fixed *r*, the maximum *n* is  $\lfloor \frac{(r+2)^2}{4} \rfloor$  and it is achieved at  $m = \lfloor \frac{r}{2} \rfloor + 1$  or at  $m = \lfloor \frac{r}{2} \rfloor + 1$ .

Finally, we prove that the bound is sharp. All the previous arguments imply that the structure of the  $\tau$ -vertex 2-critical hypergraph with maximum number of vertices must be defined in the following way. For  $m = \lfloor \frac{r}{2} \rfloor + 1$  or  $\lfloor \frac{r}{2} \rfloor + 1$  let  $U = \{1, \ldots, m\}$  and let  $A_i = \{a_1^i, \ldots, a_{r+1-m}^i\}$  with  $i \in U$ . The *r*-uniform hypergraph  $\mathcal{H}$  such that  $V(\mathcal{H}) = U \cup \bigcup_{i=1}^m A_i$  and  $E(\mathcal{H}) = \{E_1, \ldots, E_m\}$  where  $E_i = (U \setminus \{i\}) \cup A_i$  for  $i \in \{1, \ldots, m\}$ , confirms the sharpness of the inequality given in the assertion.

The construction from the proof of Theorem 1 can be generalized in an easy way resulting in the following r-uniform  $\tau$ -vertex l-critical hypergraph with a large number of vertices.

**Construction 1.** Let k, r, x be integers,  $k \ge 1$ ,  $r \ge 3$  and  $r \ge x \ge 1$  and let  $U = \{1, \ldots, k, k+1, \ldots, k+x\}$ . Next let  $m = \binom{k+x}{x}$  and let  $\{U_1, \ldots, U_m\}$  be the family of all x-element subsets of U. Additionally, let  $A_i = \{a_1^i, \ldots, a_{r-x}^i\}$  with  $i \in \{1, \ldots, m\}$  be m pairwise disjoint sets each of which is also disjoint with U.

We define an *r*-uniform hypergraph  $\mathcal{H}^* = \mathcal{H}^*(k, r, x)$  in the following way:  $E(\mathcal{H}^*) = \{E_1, \ldots, E_m\}$ , where  $E_i = U_i \cup A_i$ ,  $i \in \{1, \ldots, m\}$ ;  $V(\mathcal{H}^*) = \bigcup_{i=1}^m E_i = U \cup A$ , where  $A = \bigcup_{i=1}^m A_i$ .

**Theorem 3.** If k, r, x are integers such that  $k \ge 1$ ,  $r \ge 3$  and  $r \ge x \ge 1$ , then  $\mathcal{H}^*(k, r, x)$  is  $\tau$ -vertex (k + 1)-critical.

**Proof.** Let  $\mathcal{H}^*(k, r, x) = \mathcal{H}^*$ . We use the notations connected with  $\mathcal{H}^*$  given in Construction 1. Observe that an arbitrary (k + 1)-element subset of U is a transversal of  $\mathcal{H}^*$ . Thus  $\tau(\mathcal{H}^*) \leq k + 1$ . Suppose, for a contradiction, that Tis a transversal of  $\mathcal{H}^*$  and  $|T| \leq k$ . If  $T \subseteq U$ , then  $U \setminus T$  contains at least one x-element subset  $U_i$  and consequently  $E_i$  is an edge of  $\mathcal{H}^* - T$ . Hence T is not a transversal of  $\mathcal{H}^*$ , a contradiction. Thus  $T \setminus U = S \neq \emptyset$ . Denote  $t = |T \cap U|$ and s = |S|. There are at least  $\binom{k+x-t}{x}$  edges of  $\mathcal{H}^*$  each of which has nonempty intersection with S. It follows  $\binom{k+x-t}{x} \leq s$ . Recall that  $s+t \leq k$ . It means  $\binom{k+x-t}{x} \leq k-t$ , which is impossible for any x satisfying  $r \geq x \geq 1$ .

To observe the  $\tau$ -vertex criticality of  $\mathcal{H}^*$  it is enough to show that for each  $v \in V(\mathcal{H}^*)$  the condition  $\tau(\mathcal{H}^* - v) \leq k$  holds. If  $v \in U$ , then the removal of any k vertices of U, all different from v, results in a hypergraph without edges. If  $v \in A_i$  for some  $i \in \{1, \ldots, m\}$ , then the k-element transversal  $U \setminus U_i$  realizes the inequality  $\tau(\mathcal{H}^* - v) \leq k$ .

In the next lemma we find the maximum order of  $\mathcal{H}^*(k, r, x)$ . This result gives a lover bound on the maximum number of vertices in an *r*-uniform  $\tau$ -vertex (k + 1)-critical hypergraph.

Given k, r we introduce  $n(x) = \binom{k+x}{x}(r-x) + k + x = \binom{k+x}{k}(r-x) + k + x$ . Lemma 4. If k, r are integers such that  $k \ge 1, r \ge 3$ , then

$$\max_{1 \le x \le r} |V(\mathcal{H}^*(k, r, x))| = \max_{1 \le x \le r} n(x) = n\left(\left\lceil \frac{k(r-1)}{k+1} \right\rceil\right).$$

**Proof.** By Construction 1 we have  $\max_{1 \le x \le r} |V(\mathcal{H}^*(k, r, x))| = \max_{1 \le x \le r} n(x)$ . Consider the difference function  $D(x) = n(x) - n(x+1) = -1 + \binom{k+x}{k} [(r-x) - \frac{k+x+1}{x+1}((r-x)-1)] = -1 + \binom{k+x}{k} \frac{(r-x)(-k)+k+x+1}{x+1} = -1 + \frac{(k+x)!}{k!(x+1)!} [(x+1)(k+1) - kr] = -1 + \frac{1}{x+1} \prod_{i=1}^{k} (1 + \frac{x}{i}) [(x+1)(k+1) - kr].$ 

Since x, k and r are positive integers,  $D(x) \ge 0$  if and only if  $(x+1)(k+1) - kr \ge 1$  and therefore the maximum n(x) is reached at the smallest x such that  $D(x) \ge 0$ , i.e., at  $x = \left\lceil \frac{k(r-1)}{k+1} \right\rceil$ .

#### 3. GRAPH APPROACH

In this section we formulate some results on relations between  $\tau$ -vertex (k + 1)critical hypergraphs and forbidden subgraphs for  $\mathcal{P}(k)$ . They are preceded by the helpful lemmas.

**Lemma 5.** Let k be a non-negative integer and  $\mathcal{P} \in \mathbf{L}_{\leq}$ . If  $F \in \mathcal{C}(\mathcal{P}(k))$ , then  $F \in \mathcal{P}(k+1) \setminus \mathcal{P}(k)$ .

**Proof.** By the definition of  $\mathcal{C}(\mathcal{P}(k))$  it follows that  $F \notin \mathcal{P}(k)$ . Moreover, for an arbitrary  $v \in V(F)$  we have  $F - v \in \mathcal{P}(k)$ . It means that there exists a set W, contained in V(F - v), such that  $|W| \leq k$  and  $(F - v) - W \in \mathcal{P}$ . Because  $|W \cup \{v\}| \leq k + 1$  it leads to  $F \in \mathcal{P}(k + 1)$ .

Let  $\mathcal{P} \in \mathbf{L}_{\leq}$  and G be a graph. By  $\mathcal{H}_{\mathcal{P}}(G)$  we denote a hypergraph whose vertex set is V(G) and whose edge set is  $\{W \subseteq V(G) : G[W] \in \mathcal{C}(\mathcal{P})\}$ . Note the following facts.

**Remark 1.** Let k be a non-negative integer,  $\mathcal{P} \in \mathbf{L}_{\leq}$  and G be a graph.

- (i)  $G \in \mathcal{P}(k)$  if and only if  $\tau(\mathcal{H}_{\mathcal{P}}(G)) \leq k$ .
- (ii)  $G \in \mathcal{P}(k+1) \setminus \mathcal{P}(k)$  if and only if  $\tau(\mathcal{H}_{\mathcal{P}}(G)) = k+1$ .

**Lemma 6.** Let k be a non-negative integer and  $\mathcal{P} \in \mathbf{L}_{\leq}$ . A graph G is a forbidden subgraph for  $\mathcal{P}(k)$  if and only if  $\mathcal{H}_{\mathcal{P}}(G)$  is  $\tau$ -vertex (k + 1)-critical.

**Proof.** Suppose that  $G \in \mathcal{C}(\mathcal{P}(k))$ . By Lemma 5 and Remark 1,  $\tau(\mathcal{H}_{\mathcal{P}}(G)) = k + 1$ . Moreover, for each  $v \in V(G)$  we have  $G - v \in \mathcal{P}(k)$ , which again by Remark 1 implies  $\tau(\mathcal{H}_{\mathcal{P}}(G-v)) \leq k$ . Since  $\mathcal{H}_{\mathcal{P}}(G-v) = \mathcal{H}_{\mathcal{P}}(G) - v$  we conclude that  $\mathcal{H}_{\mathcal{P}}(G)$  is  $\tau$ -vertex (k + 1)-critical.

Now assume that  $\mathcal{H}_{\mathcal{P}}(G)$  is  $\tau$ -vertex (k + 1)-critical. Remark 1 and the equality  $\mathcal{H}_{\mathcal{P}}(G - v) = \mathcal{H}_{\mathcal{P}}(G) - v$  yield  $G \in \mathcal{P}(k + 1) \setminus \mathcal{P}(k)$  and  $G - v \in \mathcal{P}(k)$  for each  $v \in V(G)$ . Hence  $G \in \mathcal{C}(\mathcal{P}(k))$ .

Lemma 6 and Theorem 3 make it easy to formulate one more observation.

**Corollary 1.** Let k, r, x be integers such that  $k \ge 1, r \ge 3, r \ge x \ge 1$  and let  $\mathcal{P} \in \mathbf{L}_{\le}$ . If G is a graph such that  $\mathcal{H}_{\mathcal{P}}(G)$  is isomorphic to  $\mathcal{H}^*(k, r, x)$  defined in Construction 1, then G is a forbidden subgraph for  $\mathcal{P}(k)$ .

A graph G is a host-graph of a hypergraph  $\mathcal{H}$  if  $V(G) = V(\mathcal{H})$  and for each edge e of G there is an edge E of  $\mathcal{H}$  satisfying  $e \subseteq E$ . For an arbitrary family  $\mathcal{F}$  of graphs, a graph G is an  $\mathcal{F}$ -host-graph of a hypergraph  $\mathcal{H}$  when it is a host-graph of  $\mathcal{H}$  such that  $G[E] \in \mathcal{F}$  for each edge E of  $\mathcal{H}$ .

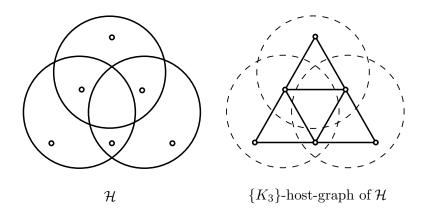


Figure 1. The example of a host-graph of a hypergraph.

Observe that for a given family of graphs  $\mathcal{F}$  and a hypergraph  $\mathcal{H}$  an  $\mathcal{F}$ -hostgraph of a hypergraph  $\mathcal{H}$  does not necessarily exist. However, we can easily find a family  $\mathcal{F}$  and a hypergraph  $\mathcal{H}$  having an  $\mathcal{F}$ -host-graph. As an example, for a fixed positive integer r, take  $\mathcal{F} = \{K_r\}$  and any r-uniform hypergraph  $\mathcal{H}$  (see Figure 1).

Furthermore, if G is a  $\mathcal{C}(\mathcal{P})$ -host-graph of a hypergraph  $\mathcal{H}$  then  $\mathcal{H}_{\mathcal{P}}(G)$  is not necessarily isomorphic to  $\mathcal{H}$  (see Figure 1 again). We use  $\mathcal{C}(\mathcal{P})$ -host-graphs to describe forbidden subgraphs for  $\mathcal{P}(k)$  with large number of vertices. In Section 2, we have constructed the family of hypergraphs  $\mathcal{H}^*(k, r, x)$  that are r-uniform  $\tau$ -vertex (k+1)-critical and have large number of vertices. So, a  $\mathcal{C}(\mathcal{P})$ -host-graph of a hypergraph  $\mathcal{H}^*(k, r, x)$  could be potentially a forbidden subgraph for  $\mathcal{P}(k)$ . First we give some examples of families  $\mathcal{F}$  of graphs for which an  $\mathcal{F}$ -host-graph of  $\mathcal{H}^*$  from Construction 1 exists.

Let G be a graph. The symbols  $\omega(G)$  and  $\alpha(G)$  denote the order of the maximum clique and the cardinality of the maximum independent set of G, respectively.

**Lemma 7.** Let  $\mathcal{F}$  be a family of graphs. Next let k, r, x be integers,  $k \ge 1, r \ge 3$ ,  $r > x \ge 1$  and  $\mathcal{H}^* = \mathcal{H}^*(k, r, x)$  be a hypergraph from Construction 1.

- (i) If there is  $F \in \mathcal{F}$  such that |V(F)| = r and  $\omega(F) \ge x$ , then there exists an  $\mathcal{F}$ -host-graph of the hypergraph  $\mathcal{H}^*$ .
- (ii) If there is  $F \in \mathcal{F}$  such that |V(F)| = r and  $\alpha(F) \ge x$ , then there exists an  $\mathcal{F}$ -host-graph of the hypergraph  $\mathcal{H}^*$ .
- (iii) If there is  $F \in \mathcal{F}$  such that |V(F)| = r and moreover  $r \ge x + k$ , then there exists an  $\mathcal{F}$ -host-graph of the hypergraph  $\mathcal{H}^*$ .

**Proof.** Using the notations from Construction 1 we show how to obtain an  $\mathcal{F}$ -host-graph G of the hypergraph  $\mathcal{H}^*$ . First we prove statements (i) and (ii). In the hypergraph  $\mathcal{H}^*$  we add all the edges between vertices in U to obtain  $K_{x+k}$  for (i) and we leave U independent for (ii). Then we choose  $F \in \mathcal{F}$  such that |V(F)| = r and  $\omega(F) \geq x$  (for (i)) or  $\alpha(F) \geq x$  (for (ii)). Now in each set  $A_i$  from Construction 1 we enter a part of F such that each  $E_i$  induces F in G. Observe that the assumption  $\omega(F) \geq x$  or  $\alpha(F) \geq x$  guarantees that all steps of this procedure can be done. To construct an  $\mathcal{F}$ -host-graph G for (iii) we choose an arbitrary vertex subset W of F of the cardinality k + x. Such a subset always exists since  $r \geq k + x$ . Next, we join some of the vertices in U by edges in such a way that the resulting graph is isomorphic to F[W]. Then, similarly to above, in each set  $A_i$  from Construction 1 we enter a part of the graph F such that each  $E_i$  induces F in the graph G.

Consider  $\mathcal{P} \in \mathbf{L}_{\leq}$  and a hypergraph  $\mathcal{H}^* = \mathcal{H}^*(k, r, x)$ . As we mentioned before if G is a  $\mathcal{C}(\mathcal{P})$ -host-graph of a hypergraph  $\mathcal{H}$ , then  $\mathcal{H}_{\mathcal{P}}(G)$  may be nonisomorphic to  $\mathcal{H}$ . Hence we do not know whether a  $\mathcal{C}(\mathcal{P})$ -host-graph of  $\mathcal{H}^*$  is a forbidden subgraph for  $\mathcal{P}(k)$  or not. In the next theorem, we solve this problem positively for some cases, regardless of whether the hypergraphs  $\mathcal{H}_{\mathcal{P}}(G)$  and  $\mathcal{H}^*$ are isomorphic.

## $\mathcal{P}$ -Apex Graphs

A set S is a vertex-cut-set in a connected graph G if G - S has at least two connected components. For a positive integer x, a connected graph G is x-vertex connected if it does not contain any vertex-cut-set of the cardinality less than x. As usual, for a given graph G and  $v \in V(G)$ , we denote by  $N_G(v)$  the set of neighbours of v in G.

**Theorem 8.** Let k, r, x be integers,  $k \ge 1$ ,  $r \ge 3$ ,  $r > x \ge 1$ , and let  $\mathcal{H}^* = \mathcal{H}^*(k, r, x)$  be the hypergraph from Construction 1. If  $\mathcal{P} \in \mathbf{L}_{\le}$  is a class of graphs such that  $\mathcal{C}(\mathcal{P})$  consists only of x-vertex connected graphs of order at least r, then each  $\mathcal{C}(\mathcal{P})$ -host-graph of the hypergraph  $\mathcal{H}^*$  is a forbidden subgraph for  $\mathcal{P}(k)$ .

**Proof.** In the proof we refer to the notations from Construction 1. Let G be an arbitrary  $\mathcal{C}(\mathcal{P})$ -host-graph of the hypergraph  $\mathcal{H}^*$ . Applying Lemma 6, the aim is to show that  $\mathcal{H}_{\mathcal{P}}(G)$  is  $\tau$ -vertex (k+1)-critical.

First we prove that any (k + 1)-element subset W of U is a transversal of  $\mathcal{H}_{\mathcal{P}}(G)$ , i.e., for any (k+1)-element subset W of U, the graph G-W does not contain any induced subgraph F satisfying  $F \in \mathcal{C}(\mathcal{P})$ . Suppose that this is not the case and let F be a subgraph of G - W such that  $F \in \mathcal{C}(\mathcal{P})$ . Denote by  $U'_1, \ldots, U'_m$  the subsets of V(G-W) that correspond to  $U_1, \ldots, U_m$  in G. Thus,  $|U'_i| \leq x - 1$  for each  $i \in \{1, \ldots, m\}$ . Furthermore, since r > x, it follows that V(F) is not contained in U - W and consequently F must contain at least one vertex of some  $A_i$  with  $i \in \{1, \ldots, m\}$ . Because of the symmetry, we may assume that  $A' = A_1 \cap V(F) \neq \emptyset$ . Since  $|A' \cup U'_1| < r$ , there is a vertex of F that does not belong to  $A' \cup U'_1$ . Hence, we can divide vertices of F into three parts  $V_1 = V(F) \cap A', V_2 = V(F) \cap U'_1$  and  $V_3 = V(F) \setminus (V_1 \cup V_2)$ . By our earlier observation  $V_3 \neq \emptyset$ . Since  $N_G(A_1) \subseteq U_1$ , it follows that  $N_F(V_1) \subseteq V_2$ . Thus,  $V_2$  is a vertex-cut-set of F. Furthermore,  $|V_2| \leq |U_1'| \leq x - 1$ , which contradicts that F is x-vertex connected and proves  $\tau(\mathcal{H}_{\mathcal{P}}(G)) \leq k+1$ . Recall that, by the construction of G, each edge of  $\mathcal{H}^*$  is an edge of  $\mathcal{H}_{\mathcal{P}}(G)$ . It means, by Theorem 3, that  $\tau(\mathcal{H}_{\mathcal{P}}(G)) \ge k+1$  and consequently  $\tau(\mathcal{H}_{\mathcal{P}}(G)) = k+1$ .

Now, we prove the  $\tau$ -vertex criticality of  $\mathcal{H}_{\mathcal{P}}(G)$ . By Remark 1 and the fact that  $\mathcal{H}_{\mathcal{P}}(G-v) = \mathcal{H}_{\mathcal{P}}(G) - v$ , we have to argue that for any  $i \in \{1, \ldots, m\}$  and for any  $v \in A_i$  we obtain  $G-v \in \mathcal{P}(k)$ . Let  $W' = U - U_i$ . Observe that |W'| = kand  $U_j \cap W' \neq \emptyset$  for  $j \neq i$ . We show that  $(G-v) - W' \in \mathcal{P}$  or equivalently that (G-v) - W' does not contain an induced subgraph isomorphic to any  $F \in \mathcal{C}(\mathcal{P})$ . Let  $U''_1, \ldots, U''_m$  be subsets of V(G - W') that correspond to  $U_1, \ldots, U_m$  in G. Thus,  $|U''_j| \leq x - 1$  for each  $j \neq i$  and  $|U''_i| = x$ . Suppose that there is  $F \in \mathcal{C}(\mathcal{P})$ such that  $F \leq (G - v) - W'$ . It is clear that there is  $j \neq i$  such that F contains at least one vertex of  $A_j$ . Therefore, similarly as above, we can divide V(F) into three parts  $V_1 = V(F) \cap A_j, V_2 = V(F) \cap U'_j$  and  $V_3 = V(F) \setminus (V_1 \cup V_2)$  with  $V_3 \neq \emptyset$ . Since  $N_F(V_1) \subseteq V_2$ , the set  $V_2$  is a vertex cut-set of F, contrary to the x-vertex connectivity of F. Theorem 8 gives us a very fruitful tool to construct forbidden subgraphs for  $\mathcal{P}(k)$ .

**Corollary 2.** Let k, x be positive integers and let  $\mathcal{P} \in \mathbf{L}_{\leq}$  be a class of graphs such that each graph in  $\mathcal{C}(\mathcal{P})$  is x-vertex connected of order at least r. If r is the order of some  $F \in \mathcal{C}(\mathcal{P})$  and  $r \geq 3$ , and  $r \geq k + x$ , then there exists G that is a forbidden subgraph for  $\mathcal{P}(k)$  and  $|V(G)| = k + x + \binom{k+x}{x}(r-x)$ .

**Theorem 9.** Let  $\mathcal{P} \in \mathbf{L}_{\leq}$ . If  $r = \max\{|F| : F \in \mathcal{C}(\mathcal{P})\}$  and  $G \in \mathcal{C}(\mathcal{P}(1))$ , then  $|V(G)| \leq \left\lfloor \frac{(r+2)^2}{4} \right\rfloor$ . Moreover, this bound is achieved for infinitely many classes  $\mathcal{P} \in \mathbf{L}_{\leq}$ .

**Proof.** By Lemma 6 and Theorem 1 we only need to show the last sentence of the statement. However, if we put k = 1 and  $x = \left\lceil \frac{r-1}{2} \right\rceil$  in Corollary 2, then for  $r \ge 3$  we obtain a forbidden subgraph for  $\mathcal{P}(k)$  with  $\left\lfloor \frac{(r+2)^2}{4} \right\rfloor$  vertices and hence the theorem follows.

The next remark is an immediate consequence of Theorem 9 and the fact that  $(\mathcal{P}(k))(1) = \mathcal{P}(k+1)$ .

**Remark 2.** Let k be a non-negative integer and  $\mathcal{P} \in \mathbf{L}_{\leq}$ . If  $\mathcal{C}(\mathcal{P})$  is finite, then the family  $\mathcal{C}(\mathcal{P}(k))$  is also finite.

#### 4. The Structure of Forbidden Subgraphs

At the beginning of this section we describe connected forbidden subgraphs for  $\mathcal{P}(k)$  in terms of connected forbidden subgraphs for  $\mathcal{P}(l)$ , where l < k. To do it we use the following hypergraph tool.

**Remark 3.** If  $\mathcal{H}_1 \cup \mathcal{H}_2$  is the union of disjoint hypergraphs  $\mathcal{H}_1$  and  $\mathcal{H}_2$ , then

$$\tau(\mathcal{H}_1 \cup \mathcal{H}_2) = \tau(\mathcal{H}_1) + \tau(\mathcal{H}_2).$$

Note that the definition of the  $\tau$ -vertex criticality of a hypergraph and Remark 3 imply the following observation.

**Remark 4.** Let s be an integer,  $s \ge 2$ . The union  $\mathcal{H}_1 \cup \cdots \cup \mathcal{H}_s$  of disjoint hypergraphs  $\mathcal{H}_1, \ldots, \mathcal{H}_s$  is  $\tau$ -vertex critical if and only if for each  $i \in \{1, \ldots, s\}$  the hypergraph  $\mathcal{H}_i$  is  $\tau$ -vertex critical.

The next result is the consequence of Remark 4.

**Theorem 10.** Let k, s be integers,  $k \ge 0$ ,  $s \ge 1$  and  $\mathcal{P} \in \mathbf{L}^a_{\le}$ . The union  $F_1 \cup \cdots \cup F_s$  of disjoint connected graphs  $F_1, \ldots, F_s$  is a forbidden subgraph for  $\mathcal{P}(k)$  if and only if there exist non-negative integers  $k_1, \ldots, k_s$  such that  $\sum_{i=1}^s k_i = k+1-s$  and for each  $i \in \{1, \ldots, s\}$  the graph  $F_i$  is a forbidden subgraph for  $\mathcal{P}(k_i)$ .

**Proof.** From Lemma 6 we have  $F_1 \cup \cdots \cup F_s \in \mathcal{C}(\mathcal{P}(k))$  if and only if  $\mathcal{H}_{\mathcal{P}}(F_1 \cup \cdots \cup F_s)$  is  $\tau$ -vertex (k+1)-critical. Since  $\mathcal{H}_{\mathcal{P}}(F_1 \cup \cdots \cup F_s) = \mathcal{H}_{\mathcal{P}}(F_1) \cup \cdots \cup \mathcal{H}_{\mathcal{P}}(F_s)$  and because of Remarks 3, 4 we know that it is equivalent to the conditions  $\tau(\mathcal{H}_{\mathcal{P}}(F_1)) + \cdots + \tau(\mathcal{H}_{\mathcal{P}}(F_s)) = k + 1$  and for each  $i \in \{1, \ldots, s\}$  the hypergraph  $\mathcal{H}_{\mathcal{P}}(F_i)$  is  $\tau$ -vertex critical. It means that there exist non-negative integers  $k_1, \ldots, k_s$  such that for each  $i \in \{1, \ldots, s\}$  the hypergraph  $\mathcal{H}_{\mathcal{P}}(F_i)$  is  $\tau$ -vertex  $(k_i + 1)$ -critical and moreover  $\sum_{i=1}^{s} (k_i + 1) = k + 1$ . From Lemma 6 these conditions are equivalent to the statement  $F_i \in \mathcal{C}(\mathcal{P}(k_i))$  for each  $i \in \{1, \ldots, s\}$  and  $\sum_{i=1}^{s} k_i = k + 1 - s$ .

**Corollary 3.** Let k be a non-negative integer and  $\mathcal{P} \in \mathbf{L}^a_{\leq}$ . If F is the union of disjoint connected graphs  $F_1, \ldots, F_s$  and  $F \in \mathcal{C}(\mathcal{P}(k))$ , then  $s \leq k + 1$ .

**Corollary 4.** Let k be a non-negative integer and  $\mathcal{P} \in \mathbf{L}^{a}_{\leq}$  and let  $|\mathcal{C}(\mathcal{P})| = p$ . The number of forbidden subgraphs for  $\mathcal{P}(k)$  that have exactly k + 1 connected components is equal to  $\binom{k+p}{k+1}$ .

**Proof.** From Theorem 10 we know that forbidden subgraphs for  $\mathcal{P}(k)$  with exactly k + 1 connected components have the form  $F_1 \cup \cdots \cup F_{k+1}$ , where for each  $i \in \{1, \ldots, k+1\}$  the condition  $F_i \in \mathcal{C}(\mathcal{P})$  holds. Let  $\mathcal{C}(\mathcal{P}) = \{H_1, \ldots, H_p\}$ . Thus, if  $m_i$  denotes  $|\{l : F_l = H_i\}|$ , then we actually are interested in the number of sequences  $(m_1, \ldots, m_p)$  whose elements are non-negative integers and for which the equality  $m_1 + \cdots + m_p = k + 1$  holds, which leads to the assertion.

The remaining part of this section is devoted to other constructions of forbidden subgraphs for  $\mathcal{P}(k)$  in terms of forbidden subgraphs for  $\mathcal{P}$ . In this consideration the structure of  $\mathcal{H}_{\mathcal{P}}(G)$  is unknown. It means that our results are based only on the analysis of graph structures.

**Construction 2.** Let s be a positive integer,  $G_1, \ldots, G_s$  be graphs and T be a forest with the vertex set  $\{x_1, \ldots, x_s\}$ . By  $T(G_1, \ldots, G_s)$  we denote the family of all graphs obtained from disjoint  $G_1, \ldots, G_s$  by the addition of exactly |E(T)| new edges, such that a new edge joins an arbitrary vertex of  $G_i$  with an arbitrary vertex of  $G_j$  when  $x_i x_j$  is an edge of T. Next we use a symbol  $(G_1, \ldots, G_s)$  to denote the family of all graphs  $T(G_1, \ldots, G_s)$  taken over all s-vertex forests T and all possible orderings of their vertices.

**Theorem 11.** If k is a non-negative integer and  $\mathcal{P} \in \mathbf{L}^a_{\leq}$  and  $G_1, \ldots, G_{k+1} \in \mathcal{C}(\mathcal{P})$ , then each graph G in  $(G_1, \ldots, G_{k+1})$  is a forbidden subgraph for  $\mathcal{P}(k)$ .

**Proof.** Suppose that  $G \in (G_1, \ldots, G_{k+1})$ . It follows that there exists a forest T with k + 1 vertices  $x_1, \ldots, x_{k+1}$  such that  $G \in T(G_1, \ldots, G_{k+1})$ . Observe that  $G \notin \mathcal{P}(k)$  since it contains k + 1 disjoint induced subgraphs that are forbidden subgraphs for  $\mathcal{P}$ .

Next, let  $v \in V(G)$ . We show that there exist k vertices  $u_2, \ldots, u_{k+1}$  in  $V(G) \setminus \{v\}$  such that the graph resulting from G by the removal of  $v, u_2, \ldots, u_{k+1}$  is in  $\mathcal{P}$ .

The construction of G implies the existence of the unique index i such that  $v \in V(G_i)$ . Let  $x_{j_1}, \ldots, x_{j_{k+1}}$  be a new ordering of vertices of T such that  $x_{j_1} = x_i$  and for  $l \geq 2$  each vertex  $x_{j_l}$  has at most one neighbour in  $\{x_{j_1}, \ldots, x_{j_{l-1}}\}$ . Such an ordering can be done by brute-force search algorithm. Suppose, without loss of generality, that  $x_{j_l} = x_l$  for each  $l \in \{1, \ldots, k+1\}$ . Consequently,  $G_{j_l} = G_l$  for each  $l \in \{1, \ldots, k+1\}$  and especially  $G_i = G_1$ .

Now we describe how to choose vertices  $u_2, \ldots, u_{k+1}$ . For each  $j \in \{2, \ldots, k+1\}$  there is at most one edge  $x_l x_j$  with l < j. Thus when such an edge exist we take as  $u_j$  the vertex of  $G_j$  that is the end of the unique edge joining  $G_j$  with  $G_l$  (see the construction of G), otherwise  $u_j$  is an arbitrary vertex of  $G_j$ . Observe that  $G - \{v, u_2, \ldots, u_{k+1}\}$  is the union of k+1 disjoint graphs  $G_1 - v$  and  $G_j - u_j$  for  $j \in \{2, \ldots, k+1\}$ . The assertion follows by the additivity of  $\mathcal{P}$  and properties of all  $G_j$ .

**Theorem 12.** Let k be a non-negative integer and  $\mathcal{P} \in \mathbf{L}_{\leq}^{a}$ . A forest G is a forbidden subgraph for  $\mathcal{P}(k)$  if and only if  $G \in (G_1, \ldots, G_{k+1})$ , where  $G_1, \ldots, G_{k+1}$ are trees that are forbidden subgraphs for  $\mathcal{P}$ .

**Proof.** By Theorem 11, it is enough to prove that if G is simultaneously a forest and a forbidden subgraph for  $\mathcal{P}(k)$ , then there are graphs  $G_1, \ldots, G_{k+1}$  belonging to  $\mathcal{C}(\mathcal{P})$  and there exists a (k+1)-vertex forest T such that  $G \in T(G_1, \ldots, G_{k+1})$ . To do it we use the induction on k.

By the additivity of  $\mathcal{P}$ , each forest that is a forbidden subgraph for  $\mathcal{P}(0) = \mathcal{P}$ is a tree. The conclusion follows from the fact that there is only one 1-vertex forest  $T = K_1$  and each graph G can be represented as  $K_1(G)$ , which means as T(G).

Assume that the implication is true for parameters less than k and  $k \ge 1$ . First suppose that G has at least two connected components  $H_1, \ldots, H_s$ . Obviously, each of them is a tree. By Theorem 10,  $H_i \in \mathcal{C}(\mathcal{P}(k_i))$ , where  $\sum_{i=1}^{s} k_i = k + 1 - s$ . Because all  $k_i$  are non-negative integers and  $s \ge 2$  we obtain  $0 \le k_i \le k - 1$  for each  $i \in \{1, \ldots, s\}$ . By the induction hypothesis,  $H_i \in T_i(G_1^i, \ldots, G_{k_i+1}^i)$ , which implies

 $G \in T(G_1^1, \ldots, G_{k_1+1}^i, \ldots, G_1^s, \ldots, G_{k_s+1}^s),$ 

where T is the union of disjoint  $T_1, \ldots, T_s$  and  $G_j^l \in \mathcal{C}(\mathcal{P})$  for each  $l \in \{1, \ldots, s\}$ and  $j \in \{1, \ldots, k_l + 1\}$ . Since each  $T_i$  has exactly  $k_i + 1$  vertices, the forest T has  $\sum_{i=1}^{s} (k_i + 1)$  vertices, which means T has k + 1 vertices. Thus G has a required form.

Now suppose that G is connected, which means G is a tree.

**Claim 13.** There is  $x \in V(G)$  such that G - x has at least one connected component in  $\mathcal{P}$  and if  $H_1, \ldots, H_p$  are all connected components of G - x belonging to  $\mathcal{P}$ , then the graph induced in G by  $V(H_1) \cup \cdots \cup V(H_p) \cup \{x\}$  is not in  $\mathcal{P}$ .

**Proof.** We describe the procedure which finds the required x in a finite number of steps.

Let  $v_0$  be an arbitrary vertex of G that is not a leaf (such a vertex always exists because  $k \ge 1$ , which implies  $|V(G)| \ge 3$ ). Next let  $G_1$  be an arbitrary connected component of  $G - v_0$  such that  $G_1 \notin \mathcal{P}$  (since G is in  $\mathcal{C}(\mathcal{P}(k))$  and  $k \ge 1$  such a connected component exists).

Let  $v_1$  be the unique neighbour of  $v_0$  in  $G_1$ . If  $G_1 - v_1 \in \mathcal{P}$ , then  $x = v_1$ . Otherwise, let  $G_2$  be an arbitrary connected component of  $G_1 - v_1$  such that  $G_2 \notin \mathcal{P}$  and let  $v_2$  be the unique neighbour of  $v_1$  in  $G_2$ . If  $G_2 - v_2 \in \mathcal{P}$ , then  $x = v_2$ . Otherwise, since G is finite, we find the finite sequence of vertices  $v_0, \ldots, v_q$  and the sequence of graphs  $G = G_0, G_1, \ldots, G_q$  such that  $G_i - v_i \notin \mathcal{P}$  for  $i \in \{0, \ldots, q-1\}, G_q \notin \mathcal{P}$  and  $G_q - v_q \in \mathcal{P}$ . Moreover for  $i \in \{1, \ldots, q\}$  the graph  $G_i$  is a connected component of  $G_{i-1} - v_{i-1}$  and  $v_i$  is the unique neighbour of  $v_{i-1}$  in  $G_i$ .

Observe that  $v_q$  can play the role of x. Indeed, the procedure implies that the connected components of  $G_q - v_q$  are simultaneously the connected components of  $G - v_q$ .

Let x be a vertex that satisfies the assumptions of Claim 13. Recall that G is a tree, which means that G - x is a forest. Since G is a forbidden subgraph for  $\mathcal{P}(k)$  we obtain  $G - x \notin \mathcal{P}(k-1)$ . It follows that G - x contains an induced subgraph  $G' \in \mathcal{C}(\mathcal{P}(k-1))$  that is a forest. By the induction hypothesis V(G') can be partitioned into k sets  $V_1, \ldots, V_k$  such that for each  $i \in \{1, \ldots, k\}$  the graph  $G'_i$  induced by  $V_i$  in G - x is forbidden for  $\mathcal{P}$ . Because  $\mathcal{P}$  is additive, all of the graphs  $G'_i$  are connected and as subgraphs of G - x they are trees. Additionally,  $(V(G'_1) \cup \cdots \cup V(G'_k) \cup \{x\}) \cap V(H_i) = \emptyset$  for  $i \in \{1, \ldots, p\}$  (keep in mind that  $H_1, \ldots, H_p \in \mathcal{P}$ , see Claim 13).

Recall that, by Claim 13,  $V(H_1) \cup \cdots \cup V(H_p) \cup \{x\}$  contains at least one subset that induces a graph, say  $G'_{k+1}$ , forbidden for  $\mathcal{P}$ . Hence  $G'_1, \ldots, G'_{k+1}$ are disjoint induced subgraphs of G, each of which is in  $\mathcal{C}(\mathcal{P})$ . Suppose, for a contradiction, that there is a vertex  $u \in V(G) \setminus \bigcup_{i=1}^{k+1} V(G'_i)$ . Since  $G \in \mathcal{C}(\mathcal{P}(k))$ we can find at most k different vertices of G-u such that the removal of all of them from G-u results in a graph in  $\mathcal{P}$ . Because G contains disjoint induced subgraphs  $G'_1, \ldots, G'_{k+1}$  that are forbidden for  $\mathcal{P}$ , it is impossible, giving a contradiction. It means  $V(G) = \bigcup_{i=1}^{k+1} V(G'_i)$  and, since G is a tree, there is a tree T with k+1 vertices such that  $G \in T(G'_1, \ldots, G'_{k+1})$ .

Below we present one more construction of graphs that are forbidden for  $\mathcal{P}(k)$ .

**Construction 3.** Let  $G_1, \ldots, G_s$  be rooted graphs, which means that for each  $i \in \{1, \ldots, s\}$  the graph  $G_i$  has a marked vertex  $v_i$ , called its root. Next let H be a graph with  $V(H) = \{x_1, \ldots, x_s\}$ . We take disjoint  $H, G_1, \ldots, G_s$  and identify vertices  $v_i$  with  $x_i$  for all  $i \in \{1, \ldots, s\}$ . By  $H|G_1, \ldots, G_s|$  we denote the family of all graphs of this type taken over all possible choices of roots  $v_1, \ldots, v_s$ . More precisely, for each graph G in  $H|G_1, \ldots, G_s|$  we have  $V(G) = \bigcup_{i=1}^s V(G_i)$  and  $E(G) = \bigcup_{i=1}^s E(G_i) \cup \{v_i v_j : x_i x_j \in E(H)\}$  with a choice of roots  $v_1, \ldots, v_s$ . Now we use a symbol  $|G_1, \ldots, G_s|$  to denote the union of sets  $H|G_1, \ldots, G_s|$  taken over all s-vertex graphs H.

**Theorem 14.** If k is a non-negative integer and  $\mathcal{P} \in \mathbf{L}^a_{\leq}$  and  $G_1, \ldots, G_{k+1} \in \mathcal{C}(\mathcal{P})$ , then each graph G in  $|G_1, \ldots, G_{k+1}|$  is a forbidden subgraph for  $\mathcal{P}(k)$ .

**Proof.** By the assumption  $G \in |G_1, \ldots, G_{k+1}|$ , we have that  $G \in H|G_1, \ldots, G_{k+1}|$  for some (k+1)-vertex graph H. Let  $x_i = v_i$  be a common vertex of H and  $G_i$ , described in Construction 3.

Because G contains disjoint induced subgraphs  $G_1, \ldots, G_{k+1}$  it follows that  $G \notin \mathcal{P}(k)$ . If  $v \in V(G)$ , then  $v \in V(G_j)$  for exactly one index  $j \in \{1, \ldots, k+1\}$ . The graph obtained from G - v by the removal of the vertex set S, where  $S = \{x_l : l \neq j\}$ , has at least k + 1 connected components each of which is in  $\mathcal{P}$ . The additivity of  $\mathcal{P}$  implies  $G - v \in \mathcal{P}(k)$ .

### 5. $P_r$ -Free Graphs

In this section we focus our attention on the class  $\mathcal{W}_r$  of graphs not containing  $P_r$ as an induced subgraph. We determine the minimum and maximum number of vertices of a graph in  $\mathcal{C}(\mathcal{W}_r(1))$ . First we consider  $\mathcal{C}(\mathcal{W}_3(1))$ . Because of Theorem 9, each graph that is forbidden for  $\mathcal{W}_3(1)$  has at most six vertices. Searching all non-isomorphic graphs of this type we can derive that  $\mathcal{C}(\mathcal{W}_3(1))$  has 14 elements:  $C_4, C_5, C_6, P_6, 2P_3, F_1, \ldots, F_9$ , where the graphs  $F_i$  for  $i \in \{1, \ldots, 9\}$  are depicted in Figure 2. Similar arguments we apply to the classes  $\mathcal{O}$  of edgeless graphs and  $\mathcal{K}$  of complete graphs. In this case, the facts  $\mathcal{C}(\mathcal{O}) = \{K_2\}$  and  $\mathcal{C}(\mathcal{K}) = \{\overline{K_2}\}$ yield  $\mathcal{C}(\mathcal{O}(1)) = \{K_3, P_4, C_4, 2K_2\}$  and  $\mathcal{C}(\mathcal{K}(1)) = \{\overline{K_3}, P_4, C_4, 2K_2\}$ .

Of course the brute searching method is not too effective if forbidden subgraphs have big orders. Thus for  $r \ge 4$  we start with determining forbidden subgraphs for  $W_r(1)$  with the minimum number of vertices. If  $G \in C(W_r(1))$ ,

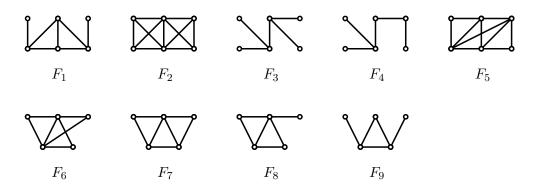


Figure 2. All the graphs in  $\mathcal{C}(\mathcal{W}_3(1)) \setminus \{C_4, C_5, C_6, P_6, 2P_3\}$ .

then G must contain an induced subgraph  $P_r$  after deletion of any vertex. Thus r + 1 is the lower bound on the number of vertices of a graph in  $\mathcal{C}(\mathcal{W}_r(1))$ . We conclude the following fact.

**Proposition 1.** If r is an integer,  $r \ge 3$ , then  $C_{r+1}$  is a forbidden subgraph for  $W_r(1)$  with the minimum number of vertices.

By Theorem 9 we have that the upper bound on the number of vertices of a graph in  $\mathcal{C}(\mathcal{W}_r(1))$  is  $\lfloor \frac{(r+2)^2}{4} \rfloor$ . However, for r = 4 we find no graph that realizes this bound. For any  $r \geq 5$  there exists a graph in  $\mathcal{C}(\mathcal{W}_r(1))$  of order  $\lfloor \frac{(r+2)^2}{4} \rfloor$ . To prove this fact we use the class of graphs that contains all the complements of graphs in  $\mathcal{W}_r$ .

For a given class of graphs  $\mathcal{P} \in \mathbf{L}_{\leq}$  let us define  $\overline{\mathcal{P}} = \{\overline{G} : G \in \mathcal{P}\}$ . It is a known fact that if  $\mathcal{P} \in \mathbf{L}_{\leq}$ , then  $\overline{\mathcal{P}}$  is also in  $\mathbf{L}_{\leq}$ . Moreover, there is a coincidence between forbidden subgraphs for  $\mathcal{P}$  and  $\overline{\mathcal{P}}$  given by the equality  $\mathcal{C}(\overline{\mathcal{P}}) = \{\overline{F} : F \in \mathcal{C}(\mathcal{P})\}$  [2]. Let  $\mathcal{P}_1, \mathcal{P}_2$  be classes of graphs. By  $\mathcal{P}_1 \circ \mathcal{P}_2$  we denote the class of all graphs G whose vertex set can be partitioned into two parts  $V_1, V_2$ (possible empty) such that, for all  $i \in \{1, 2\}$ , if  $V_i$  is non-empty, then  $G[V_i] \in \mathcal{P}_i$ . In that case  $\mathcal{P}_1 \circ \mathcal{P}_2$  is called a *product* of  $\mathcal{P}_1$  and  $\mathcal{P}_2$ . In [4] it is proved that  $F \in \mathcal{C}(\mathcal{P}_1 \circ \mathcal{P}_2)$  if and only if  $\overline{F} \in \mathcal{C}(\overline{\mathcal{P}}_1 \circ \overline{\mathcal{P}}_2)$ . It is easy to observe that for each class of graphs  $\mathcal{P}$  and a positive integer k, the class  $\mathcal{P}(k)$  is identical with  $\mathcal{P} \circ \mathcal{Q}$ , where Q consists of all the graphs of order at most k. Moreover, for such  $\mathcal{Q}$  we have  $\overline{\mathcal{Q}} = \mathcal{Q}$ . Hence, taking into account the previous consideration, we have the following observation.

**Proposition 2.** If  $\mathcal{P} \in \mathbf{L}_{\leq}$ , then

- (i)  $G \in \mathcal{P}(k)$  if and only if  $\overline{G} \in \overline{\mathcal{P}}(k)$ , and
- (ii)  $F \in \mathcal{C}(\mathcal{P}(k))$  if and only if  $\overline{F} \in \mathcal{C}(\overline{\mathcal{P}}(k))$ , and

337

(iii)  $\overline{G} \in \overline{\mathcal{P}}(k)$  if and only if  $\overline{G} \in \overline{\mathcal{P}}(k)$ .

Let us consider  $\overline{\mathcal{W}_r}$ . Thus,  $\mathcal{C}(\overline{\mathcal{W}_r}) = \{\overline{P_r}\}$  and, by Proposition 2, it follows that  $G \in \mathcal{C}(\overline{\mathcal{W}_r}(1))$  if and only if  $\overline{G} \in \mathcal{C}(\mathcal{W}_r(1))$ . As a consequence, the complement of a forbidden subgraph for  $\overline{\mathcal{W}_r}(1)$  with the maximum number of vertices is a forbidden subgraph for  $\mathcal{W}_r(1)$  with the maximum number of vertices. Since the vertex connectivity of  $\overline{P_r}$  is relatively big we will be able to apply Theorem 8. First we give the supporting observation.

**Lemma 15.** If r is an integer,  $r \ge 5$ , then  $\overline{P_r}$  is  $\left\lceil \frac{r-1}{2} \right\rceil$ -connected.

**Proof.** Let  $G = \overline{P_r}$ . Observe that the vertices of G can be divided into two sets  $W_1, W_2$  such that subgraphs induced by  $W_i$  for  $i \in \{1, 2\}$  are complete graphs and  $|W_1| = \lceil \frac{r}{2} \rceil$ ,  $|W_2| = \lfloor \frac{r}{2} \rfloor = \lceil \frac{r-1}{2} \rceil$ . Suppose that there is a vertex-cut-set S of G such that  $|S| < \lfloor \frac{r}{2} \rfloor$ . Thus G - S has two disjoint subgraphs  $G_1$  and  $G_2$  such that there is no edge joining a vertex of  $G_1$  with a vertex of  $G_2$ . Furthermore, observe that  $V(G_1) = W_1 \setminus S$  and  $V(G_2) = W_2 \setminus S$  and moreover,  $V(G_1) \neq \emptyset$  and  $V(G_2) \neq \emptyset$ . Let us denote  $W'_1 = W_1 \setminus S$  and  $W'_2 = W_2 \setminus S$ . So, by our assumptions, there is no edge joining a vertex of  $W'_1$  with a vertex of  $W'_2$  in G. This implies that in  $\overline{G}$  each vertex of  $W'_1$  is adjacent to each vertex of  $W'_2$ . If  $|W'_1| \geq 2$  and  $|W'_2| \geq 2$ , then  $\overline{G}$  contains  $C_4$ , which contradicts that  $\overline{G} = P_r$ . If one of the sets  $W'_1, W'_2$  contains exactly one vertex, then since  $|S| < \lfloor \frac{r}{2} \rfloor$ , there are at least three vertices in the second set. Thus  $\overline{G}$  has a vertex of degree three, which again gives a contradiction with the assumption that  $\overline{G}$  is a path.

By Lemma 7 we have the additional fact.

**Lemma 16.** Let r be an integer,  $r \ge 5$ . There exists a  $\{\overline{P_r}\}$ -host-graph of a hypergraph  $\mathcal{H}^*\left(1, r, \left\lceil \frac{r-1}{2} \right\rceil\right)$  given in Construction 1.

Finally, by Theorem 8, Lemma 16 and Proposition 2, we obtain the conclusion.

**Theorem 17.** Let r be an integer,  $r \geq 5$ . The complement of a  $\{\overline{P_r}\}$ -hostgraph of the hypergraph  $\mathcal{H}^*\left(1, r, \left\lceil \frac{r-1}{2} \right\rceil\right)$ , given in Construction 1, is a forbidden subgraph for  $\mathcal{W}_r(1)$  with the maximum number of vertices.

In Figure 3 we present the complement of a forbidden subgraph for  $W_5(1)$ . Theorem 17 says that this graph has the maximum number of vertices among all the graphs in  $\mathcal{C}(W_5(1))$ . Moreover, by Proposition 2, the graph in Figure 3 is in  $\mathcal{C}(\overline{W_5(1)})$  and also in  $\mathcal{C}(\overline{W_5}(1))$  and realizes the maximum order among all the graphs in both these families.

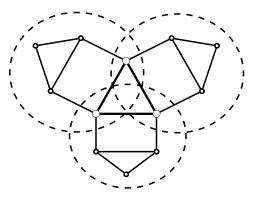


Figure 3. The complement of the graph in  $\mathcal{C}(\mathcal{W}_5(1))$  with the maximum order.

#### 6. Classes of Graphs That Are Closed Under Substitution

Let  $H, G_1, \ldots, G_n$  be graphs and  $v_1, \ldots, v_n$  be an arbitrary ordering of the set V(H). By  $H[G_1, \ldots, G_n]$  we denote the graph resulting from H by the simultaneous substitution of each vertex  $v_i$  with the graph  $G_i$ . Here the substitution of the vertex v with the graph G in the graph H means the removal of v and joining all the vertices of G with all the neighbours of v in H. A class  $\mathcal{P}$  of graphs is closed under substitution if for any graphs  $H, G_1, \ldots, G_n \in \mathcal{P}$  and every ordering of V(H), the graph  $H[G_1, \ldots, G_n]$ , called a substitution graph, is also in  $\mathcal{P}$ . By  $\mathbf{L}^*_{\leq}$  we denote the class of all non-trivial induced hereditary classes of graphs that are closed under substitution. The smallest of such ones (in the sense of the number of elements) is  $\{K_1\}$ , among most notable we should list the classes  $\mathcal{W}_r$ , where r = 2 or  $r \geq 4$ . Observe that  $P_4$ -free graphs are just cographs. In this section we characterize all forbidden subgraphs for  $\mathcal{P}(1)$  where  $\mathcal{P} \in \mathbf{L}^*_{<}$ .

A set  $W \subseteq V(G)$  is a module in a graph G if for each two vertices  $x, y \in W$ ,  $N_G(x) \setminus W = N_G(y) \setminus W$ . The trivial modules in G are V(G),  $\emptyset$  and singletons. A graph having only trivial modules is called *prime*. By **PRIME** we denote the class of all prime graphs that have at least two vertices.

In 1997 Giakoumakis [14] proved that for each class of graphs  $\mathcal{P} \in \mathbf{L}_{\leq}$  its closure under substitution  $\mathcal{P}^*$  consisting of all the graphs in  $\mathcal{P}$  and all their substitution graphs can be characterized by  $\mathcal{C}(\mathcal{P}^*)$  that consists of all minimal prime extensions of all the graphs in  $\mathcal{C}(\mathcal{P})$ . It has to be said that G' is a *minimal prime extension of* G if it is a prime induced supergraph of G and it does not contain as a proper induced subgraph any other prime induced supergraph of G.

Since for each class  $\mathcal{P} \in \mathbf{L}_{\leq}^*$  we have  $\mathcal{P} = \mathcal{P}^*$  (by the definition of  $\mathbf{L}_{\leq}^*$ ), the Giakoumakis consideration leads to the following conclusion.

**Remark 5.** If  $\mathcal{P} \in \mathbf{L}_{\leq}$ , then  $\mathcal{P} \in \mathbf{L}_{\leq}^*$  if and only if  $\mathcal{C}(\mathcal{P}) \subseteq \mathbf{PRIME}$ .

In [4] the following two theorems concerning  $C(\mathcal{P}_1 \circ \mathcal{P}_2)$  when both  $\mathcal{P}_1$ ,  $\mathcal{P}_2$  are in  $\mathbf{L}^*_{\leq}$  have been proven.

**Theorem 18** [4]. Let  $\mathcal{P}_1, \mathcal{P}_2 \in \mathbf{L}^*_{\leq}$  and let  $H \in \mathbf{PRIME}$  with  $V(H) = \{v_1, \ldots, v_n\}$ . If  $G = H[G_1, \ldots, G_n]$  and  $G \in \mathcal{C}(\mathcal{P}_1 \circ \mathcal{P}_2)$ , then  $H \notin \mathcal{P}_1$  or  $H \notin \mathcal{P}_2$  and there exists a partition (A, B, C, D) of  $\{1, \ldots, n\}$  (empty parts are allowed), such that

(i)  $G_i = K_1$  for  $i \in A$ , and (ii)  $G_i \in \mathcal{C}(\mathcal{P}_2) \cap \mathcal{P}_1$  for  $i \in B$ , and (iii)  $G_i \in \mathcal{C}(\mathcal{P}_1) \cap \mathcal{P}_2$  for  $i \in C$ , and (iv)  $G_i \in \mathcal{C}(\mathcal{P}_1 \cup \mathcal{P}_2)$  for  $i \in D$ .

A graph G, different from  $K_1$ , is strongly decomposable if in its description  $G = H[G_1, \ldots, G_n]$  with  $H \in \mathbf{PRIME}$ , all the graphs  $G_i$  satisfy  $|V(G_i)| \geq 2$ . In the next theorem we will restrict our attention to graphs that are strongly decomposable and are forbidden subgraphs for a product of classes of graphs.

**Theorem 19** [4]. Let  $\mathcal{P} \in \mathbf{L}_{\leq}^* \setminus \{\mathcal{O}, \mathcal{K}, \{K_1\}\}$ . A graph G is a forbidden subgraph for  $\mathcal{P}_1 \circ \mathcal{P}_2$  and it is strongly decomposable if and only if there exists a representation  $H[G_1, \ldots, G_n]$  of G, with  $H \in \mathbf{PRIME}$ ,  $V(H) = \{v_1, \ldots, v_n\}$ , such that either for j = 1 and l = 2 or for j = 2 and l = 1 the following three conditions hold:

- (i)  $H \in \mathcal{C}(\mathcal{P}_j)$ , and
- (ii) for each  $i \in \{1, \ldots, n\}$ ,  $G_i \in \mathcal{C}(\mathcal{P}_l)$ , and
- (iii) for  $M = \{i \in \{1, ..., n\} : G_i \notin \mathcal{P}_j\}$  and for each  $s \in \{1, ..., n\} \setminus M$ the subgraph of H induced by  $\{v_i : i \in M \cup \{s\}\}$  is in  $\mathcal{P}_l$ ; moreover, if  $M = \{1, ..., n\}$ , then  $H \in \mathcal{P}_l$ .

Observe that **PRIME** includes only two graphs,  $K_2$ ,  $\overline{K_2}$ , with two vertices, no graph on three vertices and only one graph,  $P_4$ , with four vertices. Next  $\mathcal{C}(\mathcal{O}) = \{K_2\}, \mathcal{C}(\mathcal{K}) = \{\overline{K_2}\}, \mathcal{C}(\{K_1\}) = \{K_2, \overline{K_2}\}$ . Thus if  $\mathcal{P} \in \mathbf{L}_{\leq}^* \setminus \{\mathcal{O}, \mathcal{K}, \{K_1\}\}$ , then the family  $\mathcal{C}(\mathcal{P})$  has to contain at least one graph in **PRIME**  $\setminus \{K_2, \overline{K_2}\}$ . Since each graph on at least 4 vertices contains as an induced subgraph  $K_2$  or  $\overline{K_2}$  and graphs in  $\mathcal{C}(\mathcal{P})$  are not comparable with respect to induced subgraph relation, we conclude that  $\mathcal{C}(\mathcal{P}) \cap \{K_2, \overline{K_2}\} = \emptyset$ . Hence we have the following fact.

**Remark 6.** If  $\mathcal{P} \in \mathbf{L}^*_{\leq} \setminus \{\mathcal{O}, \mathcal{K}, \{K_1\}\}$ , then  $\{K_2, \overline{K_2}\} \subseteq \mathcal{P}$ .

Recall that  $\mathcal{P}(1) = \mathcal{P} \circ \{K_1\}$  and  $\{K_1\} \in \mathbf{L}^*_{\leq}$ . Hence, from Theorem 19, we obtain the following immediate consequence.

**Corollary 5.** If  $\mathcal{P} \in \mathbf{L}_{\leq}^* \setminus \{\mathcal{O}, \mathcal{K}, \{K_1\}\}$ , then G is a forbidden subgraph for  $\mathcal{P}(1)$  that is strongly decomposable if and only if  $G = K_2[H_1, H_2]$  or  $G = \overline{K_2}[H_1, H_2] = H_1 \cup H_2$  or  $G = H_1[G_1, \ldots, G_n]$ , where  $H_1, H_2 \in \mathcal{C}(\mathcal{P})$  and  $G_1, \ldots, G_n \in \{K_2, \overline{K_2}\}$ .

**Proof.** We apply Theorem 19 together with the notations. If  $\mathcal{P} = \mathcal{P}_j$  and  $\{K_1\} = \mathcal{P}_l$ , then, by Remark 6,  $M = \emptyset$  and the graph induced in H by  $\{v_i : i \in M \cup \{s\}\}$  is  $K_1$ . Consequently we obtain that  $H_1[G_1, \ldots, G_n]$  is forbidden for  $\mathcal{P} \circ \{K_1\} = \mathcal{P}(1)$ . If  $\mathcal{P} = \mathcal{P}_l$  and  $\{K_1\} = \mathcal{P}_j$ , then H is one of the graphs  $K_2, \overline{K_2}$ . By Remark 6 we have  $M = \{1, 2\}$  and we obtain that  $K_2[H_1, H_2]$  and  $H_1 \cup H_2$  are graphs in  $\mathcal{C}(\mathcal{P}(1))$ . Theorem 19 guarantees no other strongly decomposable graphs in  $\mathcal{C}(\mathcal{P}(1))$ .

In [5] the author explained that an arbitrary graph can be obtained from a prime graph by the iterative substitution of some of its vertices by prime graphs. This procedure corresponds to the well-known construction (which has been discovered many times and is based on the Gallai Theorem [13]) called a *tree decomposition of a graph*. For a given graph G, all prime graphs applied in this tree-like iterative procedure and all their prime induced subgraphs create the unique family denoted by  $Z^*(G)$ . In the next investigation we use the following fact from this field.

**Lemma 20** [5]. Let G, G' be graphs. If  $G' \in \mathbf{PRIME}$ , then  $G' \leq G$  if and only if  $G' \in Z^*(G)$ .

Consequently we have the following observation.

**Lemma 21.** If  $\mathcal{P} \in \mathbf{L}^*_{\leq}$  and G is a graph, then  $G \in \mathcal{P}$  if and only if  $Z^*(G) \subseteq \mathcal{P}$ .

**Proof.** If  $G \in \mathcal{P}$ , then all induced subgraphs of G are in  $\mathcal{P}$ , which means  $Z^*(G) \subseteq \mathcal{P}$ .

Suppose that  $Z^*(G) \subseteq \mathcal{P}$  and, for a contradiction,  $G \notin \mathcal{P}$ . Hence there is an induced subgraph of G, say F, such that  $F \in \mathcal{C}(\mathcal{P})$  (obviously  $F \notin \mathcal{P}$ ). Remark 5 implies that F is prime, which by Lemma 20 leads to  $F \in Z^*(G)$ , and gives a contradiction.

We use Lemma 21 in proofs of forthcoming results.

**Lemma 22.** Let  $\mathcal{P} \in \mathbf{L}_{\leq}^*$  and  $H_1, H_2 \in \mathcal{C}(\mathcal{P})$ . If  $v_1, \ldots, v_n$  is an arbitrary ordering of the set  $V(H_1)$ , then  $H_1[H_2, K_1, \ldots, K_1]$  is a forbidden subgraph for  $\mathcal{P}(1)$ .

**Proof.** Let  $G = H_1[H_2, K_1, \ldots, K_1]$  and let  $V(G) = \{u_1, \ldots, u_l, v_2, \ldots, v_n\}$ , where  $v_1$  is substituted with vertices  $u_1, \ldots, u_l$  of  $H_2$ . Hence for each  $i \in \{1, \ldots, l\}$ the vertices  $u_i, v_2, \ldots, v_n$  induce  $H_1$  in G. First we observe that  $G - v \notin \mathcal{P}$  for any vertex  $v \in V(G)$ . Indeed, if  $v = v_i$  for some  $i \in \{2, \ldots, n\}$ , then  $H_2$  is an induced subgraph of G - v. If  $v = u_i$  for some  $i \in \{1, \ldots, l\}$ , then  $H_1$  is an induced subgraph of G - v.

Now we argue that for each  $v \in V(G)$  there is  $x \in V(G) \setminus \{v\}$  such that  $G - \{v, x\} \in \mathcal{P}$ . If  $v \in \{v_2, \ldots, v_n\}$ , then we choose as x one of the vertices  $u_1, \ldots, u_l$ . If  $v \in \{u_1, \ldots, u_l\}$ , then we choose as x one of the vertices  $v_2, \ldots, v_n$ . In both cases  $Z^*(G - \{v, x\})$  contains only proper prime induced subgraphs of  $H_1$  and  $H_2$ , which means  $Z^*(G - \{v, x\}) \subseteq \mathcal{P}$  and, by Lemma 21, implies  $G - \{v, x\} \in \mathcal{P}$ .

**Lemma 23.** Let  $\mathcal{P} \in \mathbf{L}_{\leq}^*$ ,  $H_1, H_2 \in \mathcal{C}(\mathcal{P})$  and  $X \in \mathbf{PRIME}$ . If  $v_1, \ldots, v_n$  is an ordering of the set V(X) such that  $X[\{v_2, \ldots, v_n\}] = H_1$  and  $X - v_i \in \mathcal{P}$  for each  $i \in \{2, \ldots, n\}$ , then  $X[H_2, K_1, \ldots, K_1]$  is a forbidden subgraph for  $\mathcal{P}(1)$ .

**Proof.** Let  $G = X[H_2, K_1, \ldots, K_1]$  and let  $V(G) = \{u_1, \ldots, u_l, v_2, \ldots, v_n\}$ , where  $v_1$  is substituted with vertices  $u_1, \ldots, u_l$  of  $H_2$ . Thus G contains two disjoint subgraphs  $H_1, H_2$  induced by vertices  $v_2, \ldots, v_n$  and  $u_1, \ldots, u_l$ , respectively. Hence  $G \notin \mathcal{P}(1)$ .

Now we argue that each pair of vertices  $u_i, v_j$ , with  $i \in \{1, \ldots, l\}$  and  $j \in \{2, \ldots, n\}$  satisfies the condition  $G - \{u_i, v_j\} \in \mathcal{P}$ . Indeed,  $Z^*(G - \{u_i, v_j\})$  contains only prime graphs that are induced subgraphs of  $H_2 - u_i$  and  $X - v_j$ . Both these graphs are in  $\mathcal{P}$ , which implies  $Z^*(G - \{u_i, v_j\}) \subseteq \mathcal{P}$ . Lemma 21 yields  $G - \{u_i, v_j\} \in \mathcal{P}$ , as we desired.

Now we are ready to prove that  $G-v \in \mathcal{P}(1)$  for each  $v \in V(G)$ , which means that for each vertex  $v \in V(G)$  there is  $x \in V(G) \setminus \{v\}$  such that  $G-\{x,v\} \in \mathcal{P}$ . If  $v = u_i$  for some  $i \in \{1, \ldots, l\}$ , then we put  $x = v_j$  for an arbitrary  $j \in \{2, \ldots, n\}$ , and if  $v = v_j$  for some  $j \in \{2, \ldots, n\}$ , then we put  $x = u_i$  for an arbitrary  $i \in \{1, \ldots, l\}$ . The earlier consideration confirms that  $G - \{x, v\} \in \mathcal{P}$  in both cases.

**Theorem 24.** Let  $\mathcal{P} \in \mathbf{L}^*_{\leq} \setminus \{\mathcal{O}, \mathcal{K}, \{K_1\}\}$ . A graph G is a forbidden subgraph for  $\mathcal{P}(1)$  if and only if G has one of the following forms:

- (i)  $G = G_1[H_1, H_2]$ , or
- (ii)  $G = H_1[G_1, \ldots, G_{|V(H_1)|}], \text{ or }$
- (iii)  $G = H_1[H_2, K_1, \dots, K_1]$ , or
- (iv)  $G = X[H_2, K_1, \dots, K_1], or$
- (v)  $G = Y[G_1, \ldots, G_s, K_1, \ldots, K_1],$

where  $H_1, H_2 \in \mathcal{C}(\mathcal{P})$  and  $G_i \in \{K_2, \overline{K_2}\}$  for all permissible *i*; further  $X, Y \in$ **PRIME** and, assuming that  $V(X) = \{v_1, \ldots, v_{n_1}\}$  and  $V(Y) = \{u_1, \ldots, u_{n_2}\}$ , the following conditions are fulfilled:

•  $X[\{v_2,\ldots,v_{n_1}\}] \in \mathcal{C}(\mathcal{P}), and$ 

342

- for each  $i \in \{2, \ldots, n_1\}$ ,  $X v_i \in \mathcal{P}$ , and
- $n_2 \ge s+2$ , and
- for each  $i \in \{1, \ldots, s\}$ ,  $Y u_i \in \mathcal{P}$ , and
- for each  $i \in \{s+1, \ldots, n_2\}$ ,  $Y-u_i \notin \mathcal{P}$  and there exists  $j \in \{s+1, \ldots, n_2\} \setminus \{i\}$ satisfying  $Y - \{u_i, u_j\} \in \mathcal{P}$ .

**Proof.** Lemmas 22, 23 and Corollary 5 show that graphs having forms (i), (ii), (iii) or (iv) are forbidden subgraphs for  $\mathcal{P}(1)$ . Recall that a graph G belongs to  $\mathcal{C}(\mathcal{P}(1))$  if the graph resulting by the removal of any vertex of G does not belong to  $\mathcal{P}$  and for each vertex  $v \in V(G)$  there exists another vertex  $x \in V(G)$  such that  $G - \{v, x\} \in \mathcal{P}$ . Observe that if a graph has the form (v), then it satisfies these conditions. Namely, if v is one of the vertices of  $G_i$  with  $i \in \{1, \ldots, s\}$ , then we choose another vertex of  $G_i$  as x. If v is one of the vertices  $u_i$  with  $i \geq s + 1$ , then the role of x is played by  $u_j$  given by the assumptions of the theorem. In both cases the conclusion follows by the construction of G.

Corollary 5 characterizes all strongly decomposable graphs in  $\mathcal{C}(\mathcal{P}(1))$ . It means that to finish the proof it is enough to show that if G is not strongly decomposable and forbidden for  $\mathcal{P}(1)$ , then G has either the form (iii) or (iv) or (v). The mentioned earlier observation that graphs in  $\mathcal{C}(\mathcal{P}(1))$  are pairwise incomparable with respect to the induced subgraph relation allows us to to simplify analysis. Namely, it is enough to show that such G contains as an induced subgraph a graph of one of the forms (i), (ii), (iii), (iv), (v). As a consequence, we observe that G has to be of the corresponding form.

Assume that G is not strongly decomposable. By Theorem 18, Remark 5 and the iterative construction of graphs via prime graphs, we can assume that G has a form  $W[U_1, \ldots, U_l, K_1, \ldots, K_1]$ , where  $W, U_1, \ldots, U_l \in \mathbf{PRIME}$  and  $V(W) = \{w_1, \ldots, w_l, w_{l+1}, \ldots, w_n\}$  with  $n \ge l+1$  (we adopt the convention that l = 0 is equivalent to  $G = W[K_1, \ldots, K_1] = W$ ). Moreover, graphs  $U_1, \ldots, U_l$ are forbidden subgraphs for  $\mathcal{P}$  or are elements of the set  $\{K_2, \overline{K_2}\}$ .

Suppose that two of the graphs  $U_1, \ldots, U_l$ , say  $U_i, U_j$ , are forbidden subgraphs for  $\mathcal{P}$ . Hence  $K_2[U_i, U_j]$  or  $\overline{K_2}[U_i, U_j]$  is an induced subgraph of G depending on whether or not  $w_i, w_j$  are adjacent in W. In both cases it leads to the conclusion that G contains an induced subgraph of the form (i).

In the next part of the proof we assume that at most one among graphs  $U_1, \ldots, U_l$  is in  $\mathcal{C}(\mathcal{P})$  and, without loss of generality, only  $U_1$  can be such a graph. Following this assumption  $W \notin \mathcal{P}$ . If not, then  $Z^*(G - v) \subseteq \mathcal{P}$ , where v is an arbitrary vertex of  $U_1$  and next, by Remark 6,  $G - v \in \mathcal{P}$  giving  $G \in \mathcal{P}(1)$ , which is impossible. Thus  $W \notin \mathcal{P}$ .

Now we consider the case  $U_1 \in \mathcal{C}(\mathcal{P})$ . It means that if  $l \geq 2$ , then  $U_2, \ldots, U_l \in \{K_2, \overline{K_2}\}$ . If there is  $W' \leq W$  such that  $W' \in \mathcal{C}(\mathcal{P})$  with  $w_1 \in V(W')$ , then G contains an induced subgraph of the form (iii). Otherwise, since  $W \notin \mathcal{P}$ 

there is  $W' \leq W$  such that  $W' \in \mathcal{C}(\mathcal{P})$  but  $w_1 \notin V(W')$  and moreover, for  $W'' = W[\{w_1\} \cup V(W')]$  we have  $W'' - x \in \mathcal{P}$  for each  $x \in V(W')$ . Observe that  $W''[U_1, K_1, \ldots, K_1] \leq G$  and  $W''[U_1, K_1, \ldots, K_1]$  is of the form (iv), which completes the proof in this case.

Suppose that  $U_1 \notin \mathcal{C}(\mathcal{P})$ . Hence  $G = W[U_1, \ldots, U_l, K_1, \ldots, K_1]$ , where  $U_1$ ,  $\dots, U_l \in \{K_2, \overline{K_2}\}$ . Assume that  $V(G) = \{w_1^1, w_1^2, \dots, w_l^1, w_l^2, w_{l+1}, \dots, w_n\},\$ where for  $i \in \{1, \ldots, l\}$   $w_i$  is substituted with vertices  $w_i^1, w_i^2$  of either  $K_2$  or  $\overline{K_2}$ . Next we show that  $W - w_i \notin \mathcal{P}$  for  $i \in \{l+1, \ldots, n\}$ . For a contradiction, let  $W - w_i \in \mathcal{P}$  for some *i* from the range. Hence, because  $K_2, \overline{K_2} \in \mathcal{P}$ , by Remark 6, we have  $Z^*(G - w_i) \subseteq \mathcal{P}$ . It implies, by Lemma 21, that  $G \in \mathcal{P}(1)$ and gives a contradiction. Therefore  $W - w_i \notin \mathcal{P}$  for  $i \in \{l+1, \ldots, n\}$ . By the definition of  $\mathcal{C}(\mathcal{P}(1))$  we know that there exists a vertex  $v \in V(G) \setminus \{w_i\}$ such that  $G - \{w_i, v\} \in \mathcal{P}$ . We ask whether or not v could be  $w_t^j$  for some  $t \in \{1, \ldots, l\}$  and  $j \in \{1, 2\}$ . Without loss of generality, let  $v = w_t^2$  for some t from the range. Thus  $G[\{w_1^1, \dots, w_l^1, w_{l+1}, \dots, w_{i-1}, w_{i+1}, \dots, w_n\}] = W - w_i$ . We observed previously that  $W - w_i \notin \mathcal{P}$ , which means that  $G - \{w_i, w_t^2\} \notin \mathcal{P}$  and excludes this possibility. Thus v must be  $w_j$  for some  $j \in \{l+1, \ldots, n\} \setminus \{i\}$  and moreover, it implies  $n \ge l+2$ . Finally, we show that if  $l \ge 1$ , then  $W - w_i \in \mathcal{P}$ for each  $i \in \{1, \ldots, l\}$ . If not, then  $W - w_i \notin \mathcal{P}$  for some  $i \in \{1, \ldots, l\}$ . It implies  $G - \{w_i^1, w_i^2\} \notin \mathcal{P}$ . By the definition of graphs in  $\mathcal{C}(\mathcal{P}(1))$  we know that there exists  $v \in V(G) \setminus \{w_i^2\}$  such that  $G - \{v, w_i^2\} \in \mathcal{P}$ . Obviously  $v \neq w_i^1$ . Moreover,  $W - w_t \leq G - \{w_t, w_i^2\}$  for each  $t \in \{l+1, \ldots, n\}$  and  $W \leq G - \{w_t^j, w_i^2\}$  for each  $t \in \{1, \ldots, l\} \setminus \{i\}$  and  $j \in \{1, 2\}$ . Because  $W - w_t \notin \mathcal{P}$  for  $t \in \{l + 1, \ldots, n\}$  and  $W \notin \mathcal{P}$ , we obtain a contradiction. Hence we conclude that  $W - w_i \in \mathcal{P}$  for each  $i \in \{1, \ldots, l\}$ . Thus, adopting l = s and  $n = n_2$ , G satisfies all the conditions that define the form (v) in this case.

In Figures 4, 5(d), 5(e), and 6 we present all possible graphs in  $\mathcal{C}(\mathcal{W}_4(1))$  that have forms pointed out in Theorem 24(i), 24(iii) and Theorem 24(iv). Some examples of graphs in  $\mathcal{C}(\mathcal{W}_4(1))$  having the construction given by Theorem 24(ii) are shown in Figures 5(a), 5(b), 5(c). Figure 7 illustrates Theorem 24(v). It refers to cases s = 0, s = 1, s = 2, represented by Y being  $C_5$ ,  $\overline{P_5}$ ,  $P_6$ , respectively. It should be mentioned here that the graph in Figure 3 has the form given by Theorem 24(v) with s = 0.

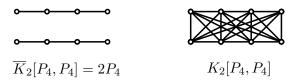


Figure 4. All the graphs in  $\mathcal{C}(\mathcal{W}_4(1))$  of the form given in Theorem 24(i).

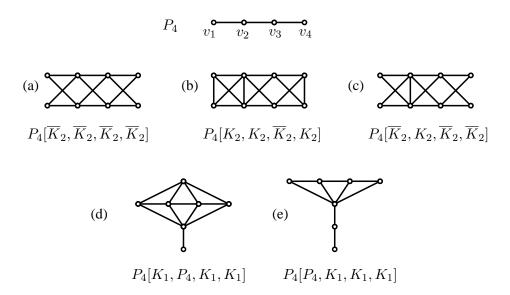


Figure 5. Some examples of graphs in  $C(W_4(1))$  of the form given in Theorem 24(ii) ((a), (b), (c)) and all the graphs in  $C(W_4(1))$  of the form given in Theorem 24(iii)) ((d), (e)).

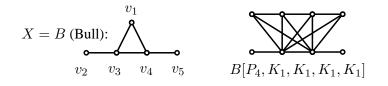


Figure 6. The unique graph in  $\mathcal{C}(\mathcal{W}_4(1))$  of the form given in Theorem 24(iv).

## 7. Concluding Remarks

In this final section we would like to present relations between the concept of a  $\mathcal{P}(k)$ -apex graph and a concept of an (H, k)-stable graph. According to [12, 16], let H be a fixed graph, a graph G is (H, k)-stable whenever the deletion of any set of k edges of G results in a graph that still contains a subgraph isomorphic to H.

An (H, k)-stable graph G is *minimal* if for every  $A \subseteq E(G)$ , |A| = k, there is  $e \in E(G) \setminus A$  such that (G - A) - e does not contain a subgraph isomorphic to H. Let us denote by Stab(H, k) the set of all minimal (H, k)-stable graphs.

**Proposition 3.** Let k be an integer and H be a graph such that  $|V(H)| \ge 4$ . Next let  $\mathcal{Q}$  be the class of all graphs that do not contain L(H) (the line graph of H) as an induced subgraph. If  $G \in Stab(H, k)$ , then  $L(G) \in C(\mathcal{Q}(k))$ .

**Proof.** On the contrary, suppose that  $L(G) \notin C(\mathcal{Q}(k))$ . Consider now two cases.

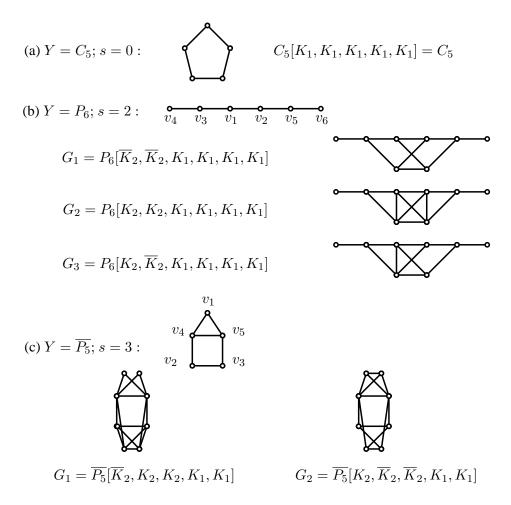


Figure 7. Some examples of graphs in  $\mathcal{C}(\mathcal{W}_4(1))$  of the form given in Theorem 24(v).

Case 1.  $L(G) \in \mathcal{Q}(k)$ . It follows that there is a set  $B \subseteq V(L(G)), |B| \leq k$ such that  $L(G) - B \in \mathcal{Q}$ . The graph L(G) - B is a line graph of some graph G'. Thus  $L(G) - B = L(G') \not\geq L(H)$ . From Whitney's Theorem [22] and assumptions it follows that  $G' \not\geq H$ . The graph G' is obtained by removing at most k edges from the graph G which correspond in a unique way to the vertices of the set B. This contradicts our assumption that  $G \in Stab(H, k)$ .

Case 2.  $L(G) \geq F \in \mathcal{C}(\mathcal{Q}(k))$ . If L(G) = F, then the conclusion is obvious. Suppose that  $L(G) \neq F$ . Thus F is a line graph of some graph G' which is a proper spanning subgraph of G. Let  $e \in E(G) \setminus E(G')$ . From the assumption  $G \in Stab(H, k)$  it follows that for the edge e there is a set  $B' \subseteq E(G) \setminus \{e\}$ , |B'| = k such that (G - e) - B' has no subgraph H. Obviously,  $|B' \cap E(G')| \leq k$ . Since  $G' \subseteq G - e$ , then G' - B' has no subgraph H. This fact implies that there is a set  $A' \subseteq V(F)$ , |A'| = k such that  $F - A' \in Q$ . This contradicts our assumption that  $F \in \mathcal{C}(\mathcal{Q}(k))$  and the proof is complete.

In [16] the minimum size of  $(P_4, k)$ -stable graphs was determined. In Section 5 of this paper we deal with the minimum and maximum order of graphs in  $\mathcal{C}(\mathcal{W}_r(k))$ . Since  $L(P_{r+1}) = P_r$  we have the following observation.

**Corollary 6.** Let k, r be integers,  $r \geq 3$ . If  $G \in Stab(P_{r+1}, k)$ , then  $L(G) \in C(\mathcal{W}_r(k))$ .

Let us define a vertex version of the *H*-stability. Let *H* be a graph and *k* be a positive integer. A graph *G* of order at least *k* is said to be (H,k)-vertex stable if for any set *S* of *k* vertices the subgraph G - S contains an induced subgraph isomorphic to *H*. An (H,k)-vertex stable graph *G* is minimal if for every  $W \subseteq V(G)$ , |W| = k, there is  $v \in V(G) \setminus W$  such that (G - W) - v does not contain *H*. Let us denote by  $Stab_V(H,k)$  the set of all minimal (H,k)-vertex stable graphs. Observe the following fact.

**Proposition 4.** If k is an integer and H is a connected graph, then  $Stab_V(H, k) = C(\mathcal{P}(k))$ , where  $\mathcal{P}$  is the class of all graphs that do not contain H as an induced subgraph.

**Proof.** If  $G \in \mathcal{C}(\mathcal{P}(k))$ , then  $G - v \in \mathcal{P}(k)$  and  $G - v \notin \mathcal{P}(k-1)$  for every  $v \in V(G)$ . In the case when  $G - v \in \mathcal{P}(k-1)$  for an vertex v, then there is a set  $A \subseteq V(G)$ , |A| = k - 1 such that  $(G - v) - A \in \mathcal{P}$ . This contradicts our assumption that  $G \in \mathcal{C}(\mathcal{P}(k))$ . It implies that for every set  $A \subseteq V(G)$ , |A| = k we have  $G - A \geq H$ , i.e.,  $G \in Stab_V(H, k)$ . Thus,  $\mathcal{C}(\mathcal{P}(k)) \subseteq Stab_V(H, k)$ .

Now let  $G \in Stab_V(H, k)$ . Then for every  $A \subseteq V(G)$ , |A| = k, there is  $v \in V(G) \setminus A$  such that (G - A) - v does not contain H as an induced subgraph. It follows that for every  $v \in V(G)$  there is a set  $A \subseteq V(G)$ , |A| = k such that  $(G - v) - A \in \mathcal{P}$ , i.e.,  $G \in \mathcal{C}(\mathcal{P}(k))$ . Hence  $Stab_V(H, k) \subseteq \mathcal{C}(\mathcal{P}(k))$ .

Yet another version of an (H, k)-stable graph was studied in a series of papers [3,6-8,10,11] where the (H, k)-vertex stability was considered taking into account, instead of induced subgraphs, subgraphs of G isomorphic to H. In case of  $H = K_q$ , both concepts coincide.

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Received 17 January 2018 Accepted 5 February 2018