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# COMPLETELY INDEPENDENT SPANNING TREES IN *k*-TH POWER OF GRAPHS

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## Abstract

Let  $T_1, T_2, \ldots, T_k$  be spanning trees of a graph G. For any two vertices u, v of G, if the paths from u to v in these k trees are pairwise openly disjoint, then we say that  $T_1, T_2, \ldots, T_k$  are completely independent. Araki showed that the square of a 2-connected graph G on n vertices with  $n \ge 4$  has two completely independent spanning trees. In this paper, we prove that the k-th power of a k-connected graph G on n vertices with  $n \ge 2k$  has k completely independent spanning trees. In fact, we prove a stronger result: if G is a connected graph on n vertices with  $\delta(G) \ge k$  and  $n \ge 2k$ , then the k-th power  $G^k$  of G has k completely independent spanning trees.

**Keywords:** completely independent spanning tree, power of graphs, spanning trees.

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### 1. INTRODUCTION

The graphs considered in this paper are finite, undirected, and simple (no loops or multiple edges). The vertex set and the edge set of G are denoted by V(G) and E(G), respectively. For a vertex  $v \in V(G)$ , the neighbour set  $N_G(v)$  is the set of vertices adjacent to v,  $\deg_G(v) = |N_G(v)|$  is the degree of v. For a subgraph H of G,  $N_H(v)$  is the set of the neighbour of v which are in H, and  $\deg_H(v) = |N_H(v)|$ is the degree of v in H. The set of (close) neighbour of a edge e in G is denoted by  $N_G(e)$  ( $N_G[e]$ ). When no confusion can occur, we shall write N(v), N(e), N[e], instead of  $N_G(v)$ ,  $N_G(e)$ ,  $N_G[e]$ , respectively. We denote by  $\delta(G)$  the minimum degree of the vertices of G. For a subset  $U \subseteq V(G)$ , the subgraph induced by U is denoted by G[U], which is the graph on U whose edges are precisely the edges of G with both ends in U. We use G - U to denote the subgraph induced by  $V(G) \setminus U$ , that is, the graph obtained from G by deleting all the vertices of Utogether with all the edges with at least one end in U. If  $U = \{u\}$ , then we shall use G - u instead of  $G - \{u\}$ . For a subset U, V of V(G), we denote  $E_G(U, V)$ the set of edges of G with one end in U and the other end in V.

For  $u \in V(G)$  and  $U \subseteq V(G)$ , the distance between u and U, denoted by  $\operatorname{dist}_G(u, U)$ , is the length of a shortest path from u to a vertex in U. When U consists of a single vertex, we write  $\operatorname{dist}_G(u, v)$  instead of  $\operatorname{dist}_G(u, \{v\})$ . For a positive integer k, the k-th power  $G^k$  of a graph G is the graph  $G^k$  whose vertex set is V(G), two distinct vertices being adjacent in  $G^k$  if and only if their distance in G is at most k. If k = 1,  $G^1 = G$ . In particular, the graph  $G^2$  is referred to as the square of G, the graph  $G^3$  as the cube of G. We say G is k-connected if |V(G)| > k and G - X is connected for every set  $X \subset V(G)$  with |X| < k. For simplicity, we denote  $[k] = \{1, \ldots, k\}$ .

A tree T of G is a spanning tree of G if V(T) = V(G). A leaf is a vertex of degree 1. An internal vertex is a vertex of degree at least 2. A rooted tree T is a tree with a specified vertex x, called the root of T. A x-tree T refer to a rooted tree with root x. The level of a vertex v of the x-tree T is the length of the path from the root x to v, the depth of the x-tree T is the maximum level of a vertex in the tree, denoted by D(T). A graph is called homeomorphically irreducible if it contains no vertices of degree 2. A homeomorphically irreducible tree is called a HIT, and a homeomorphically irreducible spanning tree of a graph is called a HIST of the graph. A caterpillar is a tree in which the internal vertices induce a path.

Let x, y be two vertices of G. An (x, y)-path is a path with the two ends xand y. Two (x, y)-paths  $P_1$ ,  $P_2$  are openly disjoint if they have no common edge and no common vertex except for the two ends x and y. Let  $T_1, T_2, \ldots, T_k$  be spanning trees in a graph H. For any two vertices u, v of H, if the paths from u to v in these k trees are pairwise openly disjoint, then we say that  $T_1, T_2, \ldots, T_k$  are completely independent spanning trees(CISTs) in G. The concept of completely independent spanning trees was proposed by Hasunuma [4]. In [4], Hasunuma gave a characterization for CISTs and proved that the underlying graph of a kconnected line digraph always contains k CISTs. It is well known [7, 9] that every 2k-edge-connected graph has k edge-disjoint spanning trees. Motivated by this, Hasunuma [5] conjectured that every 2k-connected graph has k CISTs. However, Péterfalvi [8] disproved the conjecture by constructing a k-connected graph, for each  $k \geq 2$ , which does not have two CISTs. Recently, Araki [1] provided a new characterization of the existence of k CISTs and showed the following results.

**Theorem 1** [1]. Let G be a graph with  $n \ge 7$  vertices. If  $\delta(G) \ge n/2$ , then G has two completely independent spanning trees.

In [6], Hong *et al.* give a generalization of Theorem 1.1.

**Theorem 2** [1]. If G is a 2-connected graph G on n vertices with  $n \ge 4$ , then the square  $G^2$  has two completely independent spanning trees.

It is interesting to note that the above Dirac's conditions and Fleischner's conditions is sufficient for a graph to be Hamiltonian. So, Araki [1] asked that whether other sufficient conditions for a graph to be Hamiltonian also imply the existence of two CISTs. In [3], Fan *et al.* confirmed that the well-known Ore's condition also implies the existence two CISTs.

In this paper, we generalize Theorem 2. In fact, we prove a stronger result which is Theorem 7.

First, we give the preliminaries of our results as follows.

Let  $(V_1, V_2, \ldots, V_k)$  be a partition of the vertex set V(G) and, for  $i \neq j$ ,  $B(V_i, V_j, G)$  be a bipartite graph with the edge set  $\{uv \mid uv \in E(G), u \in V_i \text{ and } v \in V_j\}$ . If the graph G is clear from the context, we may use  $B(V_1, V_2)$  instead of  $B(V_1, V_2, G)$ . A partition  $(V_1, V_2, \ldots, V_k)$  is called a CIST-partition of G if it satisfies the following two conditions:

(1) for i = 1, 2, ..., k, the induced subgraph  $G[V_i]$  is connected, and

(2) for any  $i \neq j$ , the bipartite graph  $B(V_i, V_j)$  has no tree components, that is, every connected component H of  $B(V_i, V_j)$  satisfies  $|E(H)| \geq |V(H)|$ .

The following result obtained by Araki [1] plays a key role in our proof.

**Lemma 3** [1]. A connected graph G has k completely independent spanning trees if and only if there is a CIST-partition  $(V_1, \ldots, V_k)$  of V(G).

**Lemma 4** [2]. Let G be a graph with every edge in at least two triangles. Then G contains a HIST.

Now we give the definition of a good vertex x of H to use in the proof of our result.

Given a graph H and a partition  $(V_1, \ldots, V_k)$  of its vertex set, let

$$d_{x_i} = dst_H(x, V_i), x \in V(H), i \in [k].$$

For every  $x \in V(H)$ , there exists a corresponding sequence  $(d_{t_1}, d_{t_2}, \ldots, d_{t_k})$ such that  $d_{t_1} \ge d_{t_2} \ge \cdots \ge d_{t_k}$  and  $t_1, t_2, \ldots, t_k$  is a permutation of  $x_1, x_2, \ldots, x_k$ . We say that a vertex x is good with respect to H if  $d_{t_j} \le k - j$   $(j \in [k])$ .

**Lemma 5.** Let G be a connected graph and  $H \subseteq G$ . Suppose that there are q components  $H_1, H_2, \ldots, H_q$  in G - H and S is a subset of V(H) with the following proporty: for every component  $H_s$  ( $s \in [q]$ ) of G - H, there exist a vertex  $u \in V(H_s)$  and a vertex  $v \in S$  such that  $uv \in E(G)$ . If  $H^k$  has a CIST-partition  $(V_1, \ldots, V_k)$  and every vertex of S is good with respect to H, then  $G^k$  has k completely independent spanning trees.

**Proof.** Let  $(V_1, \ldots, V_k)$  be a CIST-partition of the vertex set  $V(H^k)$ , we try to find a CIST-partition of the vertex set  $V(G^k)$  by extending the partition  $(V_1, \ldots, V_k)$ .

Let  $H_1, H_2, \ldots, H_q$  be q components of G - H. For every component  $H_s$  $(s \in [q])$ , we choose a spanning tree  $T_s$  and a vertex  $u \in V(T_s)$  such that  $uv \in E(G)$ , where  $v \in S$ . We may assume that  $T'_s = T_s \cup \{vu\}$  and  $T'_s$  is a v-tree. Let  $d_{x_i} = \text{dist}_H(v, V_i), i \in [k]$ . For the vertex v, there exists a sequence  $(d_{t_1}, d_{t_2}, \ldots, d_{t_k})$  such that  $d_{t_1} \ge d_{t_2} \ge \cdots \ge d_{t_k}$  and  $t_1, t_2, \ldots, t_k$  is a permutation of  $x_1, x_2, \ldots, x_k$ . Let  $\alpha$  be a one-to-one correspondence from [k] to [k] such that  $\text{dist}_H(v, V_{\alpha(j)}) = d_{t_j}$ .

Let  $L_j$  be the vertex set of all vertices in *j*-th level of  $T'_s$  for  $j \in \{1, \ldots, D(T'_s)\}$ . For every  $w \in L_j$ , we assign it to  $V_{\alpha(j \pmod{k})}$ . For other components, we repeat the above operation, and it follows that we obtain a new partition  $(V'_1, V'_2, \ldots, V'_k)$ of the vertex set  $V(G^k)$ . It remains to show that  $(V'_1, V'_2, \ldots, V'_k)$  is a CISTpartition of  $V(G^k)$ .

If  $1 \leq j \leq k$ , since every vertex of S is good with respect to H and  $L_j \subset V'_{\alpha(j)}$  for any  $w \in L_j$ , we have

$$\operatorname{dist}_G(w, V_{\alpha(j)}) \leq \operatorname{dist}_G(v, V_{\alpha(j)}) + \operatorname{dist}_G(v, w) \leq k - j + j = k.$$

Thus,

$$E_{G^k}(\{w\}, V_{\alpha(j)}) \neq \emptyset.$$

If  $k+1 \leq j \leq D(T'_s)$ , for any  $w_1 \in L_j, w_2 \in L_{j-k}$ , then

$$L_j, L_{j-k} \subset V_{\alpha(j \pmod{k})}, \operatorname{dist}_G(w_1, w_2) \leq k.$$

Thus,

$$w_1w_2 \in E(G^k).$$

It is easy to see that the induced graph  $G^k[V'_i]$  is connected for  $i \in [k]$ .

Note that  $\deg_{B(V'_i,V'_j)}(w) \geq 1$  for every vertex  $w \in V'_i \setminus V_i, j' \neq i'$ . Since  $B(V_i, V_j)$  has no tree components and the vertex w is adjacent to  $V_j$  in  $G^k$  by alternative path between  $V'_i$  and  $V'_j$ , we get that  $B(V'_i, V'_j)$  has no tree component.

Hence,  $(V'_1, V'_2, \ldots, V'_k)$  is a CIST-partition of  $V(G^k)$ . By Lemma 3,  $G^k$  has k completely independent spanning trees.

**Lemma 6.** For a homeomorphically irreducible tree (HIT) T with  $|V(T)| \ge 2k$ , the k-th power  $T^k$  of T has k completely independent spanning trees.

**Proof.** We first consider the longest path  $P = x_0 x_1 \cdots x_{|P|}$  of T.

Case 1. |P| < k. If |P| < k, then  $T^k$  is a complete graph. Also,  $|V(T)| \ge 2k$ and any partition  $(V_1, V_2, \ldots, V_k)$  of V(T) with  $|V_i| \ge 2$  is a CIST-partition of  $V(T^k)$ . Hence, by Lemma 3,  $T^k$  has k completely independent spanning trees.

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Case 2.  $|P| \ge k$ . Denote  $P_0 = x_0 x_1 \cdots x_k$ . Since T is a homeomorphically irreducible tree (HIT), we choose a caterpillar  $T_0$  such that its internal vertices are  $V(P_0 - \{x_0, x_k\})$  and its leaf vertices are  $N(x_1) \cup \cdots \cup N(x_{k-1}) \setminus V(P_0 - \{x_0, x_k\})$ . We regard  $T_0$  as a rooted tree which is rooted at  $x_0$  in the following proof.

Let  $L_i$  be the set of all vertices with the same level of  $x_i$  in  $T_0$ , where  $i \in \{0, 1, \ldots, k\}$ . Note that

$$L_0 = \{x_0\}, \ L_1 = \{x_1\}, \ |L_i| \ge 2, \ i \in \{2, \dots, k\}.$$

We have  $|T_0| \ge 2k$  and the distance between x and y is at most k for every pair  $x, y \in V(T_0)$ . Thus,  $T^k[T_0]$  is a complete graph and any partition  $(V_1, V_2, \ldots, V_k)$  of the vertex set  $V(T^k[T_0])$  with  $|V_i| \ge 2$  is a CIST-partition. Specially, we choose a partition of  $V(T^k[T_0])$  as  $(L_0 \cup L_1, L_2, \ldots, L_k)$ .

Since P is a longest path in T and  $T_0$  is a caterpillar,  $E_T(L_0 \cup L_1 \cup L_2 \cup V(P_0 - x_k), T - T_0)$  is empty. Let  $S = V(T_0 - (L_0 \cup L_1 \cup L_2 \cup V(P_0 - x_k)))$ , it follows that  $xx_{i-1} \in E(T)$  for any  $x \in S \cap L_i$   $(i \geq 3)$  and

$$dist_T(x, L_i) = 0,$$
  

$$dist_T(x, L_j) = i - j, \ 1 \le j < i,$$
  

$$dist_T(x, L_j) = j - i + 2, \ i + 1 \le j \le k.$$

In addition, there exists a corresponding sequence  $(d_{t_1}, d_{t_2}, \ldots, d_{t_k})$  such that  $d_{t_1} \geq d_{t_2} \geq \cdots \geq d_{t_k}$  and  $t_1, t_2, \ldots, t_k$  is a permutation of  $\operatorname{dist}_T(x, L_0 \cup L_1)$ ,  $\operatorname{dist}_T(x, L_2), \ldots, \operatorname{dist}_T(x, L_k)$ . So,  $d_{t_j} \leq k - j$   $(j \in [k])$ . Hence, every vertex of S is good with respect to  $T_0$ . By Lemma 5,  $T^k$  has k completely independent spanning trees.

**Theorem 7.** If G is a connected graph on n vertices with  $n \ge 2k$  and  $\delta(G) \ge k$ , then the k-th power  $G^k$  of G has k completely independent spanning trees.

**Proof.** If k = 1, then the theorem holds trivially. Therefore, we may assume that  $k \geq 2$ . Now, suppose that k = 2. Since  $\delta(G) \geq 2$ , we have that G has a cycle  $C_m$ .

If  $|V(C_m)| \ge 2k = 4$ , then let  $C_m = x_1 x_2 \cdots x_m$  and let a partition of the vertex set  $V(C_m)$  be as follows:

$$V_i = \{x_j \mid j \equiv i \pmod{2}, 1 \le j \le m\}, i \in [2].$$

Since  $m \geq 4$ , we have  $|V_i| \geq 2$ . It is easy to see that the induced graph  $G_m^2[V_i]$  is connected for  $i \in [2]$ . Actually,  $G_m^2[V_i]$  is either a path or a cycle. If  $x_j \in V_i$ , then  $|V_{3-i} \cap \{v | \operatorname{dist}_{C_m}(x_j, v) \leq 2\}| \geq 2$ . Thus,  $\operatorname{deg}_{B(V_i, V_{3-i})}(x_j) \geq 2$  and  $B(V_i, V_{3-i})$  has no tree component. Hence,  $(V_1, V_2)$  is a CIST-partition of  $C_m^2$ .

Let  $H_i$   $(i \in [q])$  be the connected components of graphs  $G - C_m$ . Since G is connected, we have that there exists an edge which connects  $H_i$  to  $C_m$  for every

 $i \in [q]$ . Also, we have  $\operatorname{dist}_G(x, V_i) = 0$  for every vertex  $x \in V_i$   $(i \in [2])$  and  $\operatorname{dist}_G(x, V_{3-i}) = 1$ . Hence, for the set  $S = V(C_m)$ , every vertex of S is good with respect to  $C_m$ . By Lemma 5,  $G^2$  has 2 completely independent spanning trees.

Now we assume that |C| < 4 for any cycle C of G.

We choose a cycle  $C = x_1 x_2 x_3$ . Note that  $n \ge 4$ , therefore there exists a vertex  $y \in V(G-C)$  such that  $yx_i \in E(G)$ . Without loss of generality, we assume that  $yx_1 \in E(G)$ . Let  $H = C \cup \{yx_1\}$  and let a partition of the vertex set  $V(H^2)$  be as follows:

$$V_1 = \{x_1, x_2\}, V_2 = \{x_3, y\}.$$

Since  $n \ge 4$ , we have  $|V_i| \ge 2$ . It is easy to see that  $(V_1, V_2)$  is a CISTpartition of the vertex set  $V(H^2)$ . As we have stated in the previous case, for the set S = V(H), every vertex of S is good with respect to H. By Lemma 5,  $G^2$  has 2 completely independent spanning trees.

Thus, we only consider the case  $k \geq 3$ .

Case 1. There exists an edge xy such that  $|N(x) - N[y]| \ge k - 1$  and  $|N(y) - N[x]| \ge k - 1$ .

$$xy = x_1y_1, H = G[N[x_1] \cup N[y_1]] - N(x_1) \cap N(y_1),$$
$$N(x_1) = \{y_1, x_2, x_3, \dots, x_{\deg_G(x_1)}\},$$
$$N(y_1) = \{x_1, y_2, y_3, \dots, y_{\deg_G(y_1)}\}.$$

Since  $\delta(G) \geq k \geq 3$  and dist $(u, v) \leq 3$  for any two vertices u, v, combine with  $|V(H)| \geq 2k$ ,  $H^k$  is a complete graph. Thus, any partition  $(V_1, V_2, \ldots, V_k)$ of the vertex set  $V(H^k)$  with  $|V_i| \geq 2$  is a CIST-partition. Specially, we choose a partition of the vertex set  $V(H^k)$  as follows:

$$V_{1} = \{x_{1}, y_{1}\},$$
  

$$V_{i} = \{x_{j} \mid j \equiv i \pmod{(k-1)}, \ 2 \leq j \leq \deg_{G}(x_{1})\}$$
  

$$\cup \{y_{j} \mid j \equiv i \pmod{(k-1)}, \ 2 \leq j \leq \deg_{G}(y_{1})\}, \ i \in \{2, \dots, k\}.$$

For every  $x \in V_i$   $(i \ge 2)$ , we have

$$dist_G(x, V_i) = 0$$
,  $dist_G(x, V_1) = 1$ ,  $dist_G(x, V_j) = 2$ ,  $j \neq i$ .

Hence, for the set S = V(H), every vertex of S is good with respect to H. By Lemma 5,  $G^k$  has k completely independent spanning trees.

Case 2. There exists an edge xy such that  $|N(x) - N[y]| \ge k - 2$  and  $|N(x) - N[y]| \ge k - 2$  and which does not satisfy Case 1. Let

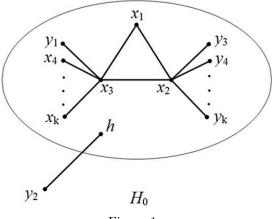
$$H = G[N[xy]].$$

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Since  $\delta(G) \ge k$ ,  $|V(H)| \ge 2k - 1$ .

Case 2.1. If d(x) > k or d(y) > k, then xy satisfies the Case 1 which is a contradiction.

Case 2.2. If d(x) = k, d(y) = k, then |V(H)| = 2k - 1. Since  $n \ge 2k$ , there exists a vertex  $w \in V(G - H)$  such that w is adjacent to a vertex h of  $H \setminus \{x, y\}$ . In other words,  $h \in N(x) \setminus \{y\}$  or  $N(y) \setminus \{x\}$ . We suppose that  $H_0 = G[V(H) \cup w]$ .





(1) We first consider the case  $k \ge 4$ . If  $h \in N(x) \setminus \{y\}$  (or  $h \in N(y) \setminus \{x\}$ ), then we label  $H_0$  as in Figure 1. Let

$$x = x_3, \ y = x_2 \text{ (or } x = x_2, \ y = x_3),$$
$$N(x) \cap N(y) = \{x_1\}, \ w = y_2,$$
$$N(x_3) = \{x_2, x_1, y_1, x_4, x_5, \dots, x_k\},$$
$$N(x_2) = \{x_3, x_1, y_3, y_4, y_5, \dots, y_k\}.$$

Since  $\delta(G) \ge k \ge 4$ ,  $H_0^k$  is a complete graph. Thus, any partition  $(V_1, V_2, \ldots, V_k)$  of the vertex set  $V(H_0^k)$  with  $|V_i| \ge 2$  is a CIST-partition. Specially, we choose a partition of the vertex set  $V(H_0^k)$  as follows:

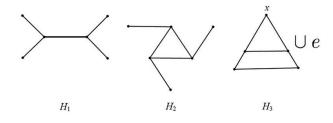
$$V_i = \{x_i, y_i\}, \ i \in [k].$$

If 
$$h = x_{\ell}$$
 ( $\ell = 1, 4, 5, ..., k$ ) or  $h = y_{\ell}$  ( $\ell = 1$ ), then for  $j \in [k]$  we have  
 $\operatorname{dist}_{G}(x_{1}, V_{1}) = 0$ ,  $\operatorname{dist}_{G}(x_{1}, V_{2}) = 1$ ,  $\operatorname{dist}_{G}(x_{1}, V_{j}) \leq 2$  ( $j \neq 1, 2$ ),  
 $\operatorname{dist}_{G}(y_{1}, V_{1}) = 0$ ,  $\operatorname{dist}_{G}(y_{1}, V_{3}) = 1$ ,  $\operatorname{dist}_{G}(y_{1}, V_{j}) \leq 2$  ( $j \neq 1, 3$ ),

$$\begin{aligned} \operatorname{dist}_{G}(y_{2}, V_{2}) &= 0, \ \operatorname{dist}_{G}(y_{2}, V_{\ell}) = 1, \ \operatorname{dist}_{G}(y_{2}, V_{3}) \leq 2, \\ \operatorname{dist}_{G}(y_{2}, V_{j}) &\leq 3 \ (j \neq 2, 3, \ell), \\ \operatorname{dist}_{G}(y_{i}, V_{i}) &= 0, \ \operatorname{dist}_{G}(y_{i}, V_{2}) = 1, \ \operatorname{dist}_{G}(y_{i}, V_{j}) \leq 2, 3 \leq i \leq k \ (j \neq 2, i), \\ \operatorname{dist}_{G}(x_{i}, V_{i}) &= 0, \ \operatorname{dist}_{G}(x_{i}, V_{3}) = 1, \ \operatorname{dist}_{G}(x_{i}, V_{j}) \leq 2, 4 \leq i \leq k \ (j \neq 3, i). \end{aligned}$$

Again, for the set  $S = V(H - \{x_2, x_3\})$ , every vertex of S is good with respect to H. By Lemma 5,  $G^k$  has k completely independent spanning trees.

(2) Now suppose that k = 3.





**Claim 8.** If a connected graph G contains a subgraph which is isomorphic to one of the  $H_1, H_2, H_3$  in Figure 2 (where  $H_3 = H'_3 \cup e$  and e has exactly one end in  $V(H'_3)$ ), then  $G^3$  has 3 completely independent spanning trees.

**Proof.** If  $H_1 \subseteq G$ , then the proof follows by Case 1. If  $H_2 \subseteq G$ , then  $H_2$  is isomorphic to  $H_0 = G[V(H) \cup y_2]$  in Case 2.2(1), where H = G[W],  $W = \{x_1, x_2, x_3, y_1, y_3\}$  and  $y_2$  is adjacent to a vertex  $x_1$  of H. The proof follows by Case 2.2(1). Suppose  $H_3 \subseteq G$ , if e is adjacent to x, then there exists a subgraph of  $H_3$  which is isomorphic to  $H_2$ . Otherwise, e is adjacent to  $V(H'_3) \setminus x$ , then we obtain a subgraph of  $H_3$  which is isomorphic to  $H_3$ .

Now we begin to prove the case k = 3 and we relabel H as in left of Figure 3.

Since k = 3, we have  $h \neq y, z$ . If h = x, then  $H \cup \{wh\}$  contains a subgraph which is isomorphic to  $H_2$ , it is true by Claim 8. Thus, h = u or h = v. By symmetry, we may assume that h = u and  $H_0 = G[V(H) \cup w]$ , as in right of Figure 3.

If  $\deg_{G-H_0}(u) \ge 1$ , then G contains a subgraph which is isomorphic to  $H_1$ , it is true by Claim 8. Thus,  $N(u) \subseteq V(H_0)$ .

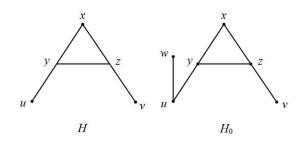


Figure 3.

If  $uv \in E(G)$ , then there exists a subgraph of G which is isomorphic to  $H_3$ , it is true by Claim 8. Otherwise,  $uv \notin E(G)$  and  $ux \in E(G)$ .

If  $xv \in E(G)$ , then we obtain a subgraph of G which is isomorphic to  $H_1$ , it is true by Claim 8. Otherwise,  $xv \notin E(G)$ . Since  $\delta \geq 3$ , we have  $\deg_{G-\{x,y,u\}}(v) \geq$ 2. Thus, we also can obtain a subgraph of G which is isomorphic to  $H_1$ , it is true by Claim 8.

Hence,  $G^3$  has 3 completely independent spanning trees.

Case 3. Every edge e of G is contained in at least two triangles. By Lemma 4, G contains a HIST T. Also, by Lemma 6, the k-th power  $T^k$  of T has k completely independent spanning trees.

Hence, the k-th power  $G^k$  has k completely independent spanning trees. The proof of Theorem 7 is completed.

An immediate consequence of Theorem 7 is the following corollary.

**Corollary 9.** If G is a k-connected graph on n vertices with  $n \ge 2k$ , then the k-th power  $G^k$  of G has k completely independent spanning trees.

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