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# RAINBOW CONNECTION NUMBER OF GRAPHS WITH DIAMETER 3

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#### Abstract

A path in an edge-colored graph G is rainbow if no two edges of the path are colored the same. The rainbow connection number rc(G) of G is the smallest integer k for which there exists a k-edge-coloring of G such that every pair of distinct vertices of G is connected by a rainbow path. Let f(d)denote the minimum number such that  $rc(G) \leq f(d)$  for each bridgeless graph G with diameter d. In this paper, we shall show that  $7 \leq f(3) \leq 9$ .

**Keywords:** edge-coloring, rainbow path, rainbow connection number, diameter.

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## 1. INTRODUCTION

All graphs in this paper are undirected, finite, and simple. We refer to book [2] for notation and terminology not described here. A path  $u_0u_1\cdots u_k$  is called a  $P_{uv}$ path, where  $u = u_0$  and  $u_k = v$ . The distance between two vertices x and y in G, denoted by d(x, y), is the number of edges of a shortest path between them. The eccentricity of a vertex x, denoted by ecc(x), is  $\max_{y \in V(G)} d(x, y)$ . The radius and diameter of G, denoted by rad(G) and diam(G), are  $\min_{x \in V(G)} ecc(x)$  and  $\max_{x \in V(G)} ecc(x)$ , respectively. A vertex u is a center if ecc(u) = rad(G). A path in an edge-colored graph G, where adjacent edges may have the same color, is *rainbow* if no two edges of the path are colored the same. An edgecoloring of a graph G is a *rainbow-connected edge-coloring* if every pair of distinct vertices of G is connected by a rainbow path. The *rainbow connection number* rc(G) of G is the minimum integer k for which there exists a *rainbow-connected* k-edge-coloring of G. It is easy to see that  $diam(G) \leq rc(G)$  for any connected graph G.

The rainbow connection number was introduced by Chartrand, Johns, McKeon, and Zhang in [4]. It has application in transferring information of high security in multicomputer networks. We refer the readers to [3, 8] for details.

Chakraborty, Fischer, Matsliah, and Yuster [3] investigated the hardness and algorithms for the rainbow connection number, and showed that given a graph G, deciding if rc(G) = 2 is NP-complete. Bounds for the rainbow connection number of a graph have also been studied in terms of other graph parameters, for example, radius and diameter, etc. [1, 5, 6, 7].

Let f(d) denote the minimum number such that each bridgeless graph G with diameter d has a rainbow-connected f(d)-edge-coloring. It is easy to check that f(1) = 1. In [7], we showed that f(2) = 5. In this paper, we shall show that  $7 \le f(3) \le 9$ .

The following theorem will be used in this paper.

**Theorem 1** [5]. For every bridgeless graph G,

$$rc(G) \leq \sum_{i=1}^{\operatorname{rad}(G)} \min\{2i+1, \eta(G)\} \leq \operatorname{rad}(G)\eta(G),$$

where  $\eta(G)$  is the smallest integer such that every edge of G is contained in a cycle of length at most  $\eta(G)$ .

In this paper, we investigate the upper bound on the rainbow connection number of bridgeless graphs with diameter 3, and obtain the following result.

**Theorem 2.** For every bridgeless graph G with diameter 3,  $rc(G) \leq 9$ .

If each edge of a bridgeless graph G with diameter 3 belongs to a triangle, then  $rc(G) \leq 9$  by Theorem 1. Thus, we suppose that there exists an edge e such that e does not belong to any triangle in G.

This paper is organized as follows. In Section 2, we partition V(G), and present a partial edge-coloring of G under this partition. In Section 3, we further partition V(G) and give a complete edge-coloring of G under this partition. In Section 4, we prove that the edge-coloring in Section 3 is a rainbow-connected 9-edge-coloring of G, and give a class of bridgeless graphs with diameter 3 and rainbow connection number at least 7.

# 2. A PARTIAL EDGE-COLORING

Let G be a graph. For any integer  $k \ge 1$ , the k-step open neighborhood  $N^k(X)$ is  $\{y \in V(G) : d(X, y) = k\}$ . We simply write N(X) for  $N^1(X)$  and  $N^k(x)$  for  $N^k(\{x\})$ . Similarly, the k-step closed neighborhood  $N^k[X]$  is  $\{y \in V(G) : d(X, y) \le k\}$ . We simply write N[X] for  $N^1[X]$  and  $N^k[x]$  for  $N^k[\{x\}]$ .

Let c be an edge-coloring of G, and let P be a rainbow path in G. We use c(P) to denote the set of colors used on P, that is,  $c(P) = \{c(e): \text{ the edge } e \text{ belongs}$  to P}. If  $c(P) \subseteq \{k_1, k_2, \ldots, k_r\}$ , then P is a  $\{k_1, k_2, \ldots, k_r\}$ -rainbow path. In particular, an edge e is a k-color edge if c(e) = k. We use  $x_0 \stackrel{c_1}{\sim} x_1 \stackrel{c_2}{\sim} \cdots \stackrel{c_k}{\sim} x_k$  to denote a rainbow path  $x_0 x_1 \cdots x_k$  with  $c(x_{i-1}x_i) = c_i$  for each  $1 \leq i \leq k$ . Let  $X_1, X_2, \ldots, X_{k-1}$  be pairwise disjoint vertex subsets of G. The notation  $x_0 \stackrel{c_1}{\sim} X_1 \stackrel{c_2}{\sim} \cdots \stackrel{c_k}{\sim} x_k$  means that there exists a rainbow path  $x_0 \stackrel{c_1}{\sim} x_1 \stackrel{c_2}{\sim} \cdots \stackrel{c_k}{\sim} x_k$ , where  $x_i \in X_i$  for  $1 \leq i \leq k - 1$ .

Recall that e is an edge not belonging to any triangle in G. Let u and v be the ends of e.

Since e does not belong to any triangle, for the open neighborhood,  $N(\{u, v\})$ , of  $\{u, v\}$  in G, we can divide it as follows:

$$A = N(u) \setminus \{v\},\$$
  
$$B = N(v) \setminus \{u\}.$$

See Figure 1 for details.

For the 2-step open neighborhood,  $N^2(\{u, v\})$ , of  $\{u, v\}$  in G, we can divide it as follows:

$$\begin{aligned} X &= \{ x \in N(A) \setminus N(B) : x \notin A \cup B \cup \{u, v\} \}, \\ Y &= \{ x \in N(B) \setminus N(A) : x \notin A \cup B \cup \{u, v\} \}, \\ Z &= \{ x \in N(A) \cap N(B) : x \notin A \cup B \cup \{u, v\} \}. \end{aligned}$$

See Figure 1 for details. It is easy to see that  $x \in X$  if and only if  $x \notin N[\{u, v\}]$ , d(x, u) = 2 and d(x, v) = 3;  $y \in Y$  if and only if  $y \notin N[\{u, v\}]$ , d(y, u) = 3 and d(y, v) = 2;  $z \in Z$  if and only if  $z \notin N[\{u, v\}]$ , d(x, u) = 2 and d(x, v) = 2.

Note that for  $x \in N^3(\{u, v\})$ , we have d(x, u) = d(x, v) = 3, since diam(G) = 3, that is,  $N(x) \cap N(A) \neq \emptyset$  and  $N(x) \cap N(B) \neq \emptyset$ .

For the 3-step open neighborhood,  $N^3(\{u, v\})$ , of  $\{u, v\}$  in G, we can partition  $N^3\{u, v\}$  based on the distribution of the neighbors of x as follows:

$$W = \{x \in N^3(\{u, v\}) : N(x) \cap X \neq \emptyset \text{ and } N(x) \cap Y \neq \emptyset\},\$$
  
$$I = \{x \in N^3(\{u, v\}) \setminus W : N(x) \cap X \neq \emptyset \text{ and } N(x) \cap Z \neq \emptyset\},\$$

$$K = \{x \in N^3(\{u, v\}) \setminus (W \cup I) : N(x) \cap Y \neq \emptyset \text{ and } N(x) \cap Z \neq \emptyset\},\$$
  
$$J = \{x \in N^3(\{u, v\}) \setminus (W \cup I \cup K) : N(x) \cap Z \neq \emptyset\}.$$

See Figure 1 for details. It is easy to see that  $N^3(\{u, v\}) = I \cup J \cup K \cup W$ . At this point, we further partition A and B as follows:

$$A_{1} = \{x \in A : N(x) \cap (B \cup X \cup Z) \neq \emptyset\},\$$

$$A_{2} = \{x \in A \setminus A_{1} : N(x) \cap (A \setminus A_{1}) \neq \emptyset\},\$$

$$A_{3} = A \setminus (A_{1} \cup A_{2}),\$$

$$B_{1} = \{x \in B : N(x) \cap (A \cup Y \cup Z) \neq \emptyset\},\$$

$$B_{2} = \{x \in B \setminus B_{1} : N(x) \cap (B \setminus B_{1}) \neq \emptyset\},\$$

$$B_{3} = B \setminus (B_{1} \cup B_{2}).\$$

That is,  $A_1$  consists of vertices which have neighbors outside  $A \cup \{u\}$ ,  $A_2$  consists of vertices which do not have neighbors outside A (apart from u) but have neighbors in  $A \setminus A_1$ , and  $A_3$  consists of vertices which have neighbors only in  $A_1$  (apart from u). It is clear that for each  $x \in A_2$ , there exists a vertex  $x' \in A_2$  such that xx'u is a triangle. Similar results also hold for  $B_1, B_2$  and  $B_3$ .

Note that there may exist edges between between  $A_1$  and  $A_2$ , but it does not matter for our proof.

Meanwhile, we partition X and Y as follows:

$$\begin{split} X_1 &= \{x \in X : N(x) \cap (Y \cup Z \cup I \cup W) \neq \emptyset\}, \\ X_2 &= \{x \in X \setminus X_1 : N(x) \cap (X \setminus X_1) \neq \emptyset\}, \\ X_3 &= \{x \in X \setminus (X_1 \cup X_2) : N(x) \subseteq A\}, \\ X_4 &= X \setminus (X_1 \cup X_2 \cup X_3), \\ Y_1 &= \{y \in Y : N(y) \cap (X \cup Z \cup K \cup W) \neq \emptyset\}, \\ Y_2 &= \{y \in Y \setminus Y_1 : N(y) \cap (Y \setminus Y_1) \neq \emptyset\}, \\ Y_3 &= \{y \in Y \setminus (Y_1 \cup Y_2) : N(y) \subseteq B\}, \\ Y_4 &= Y \setminus (Y_1 \cup Y_2 \cup Y_3). \end{split}$$

That is,  $X_1$  consists of vertices which have neighbors outside X (apart from  $A_1$ ),  $X_2$  consists of vertices which do not have neighbors outside X (apart from  $A_1$ ) but have neighbors in  $X \setminus X_1$ ,  $X_3$  consists of vertices which have neighbors only in  $A_1$ , and  $X_4$  consists of vertices which have neighbors only in  $X_1$  (apart from  $A_1$ ). Similar results also hold for  $Y_1, Y_2, Y_3$  and  $Y_4$ .

By the definitions of sets  $A_1, A_2$  and  $A_3$ , we know that  $N(X_3) \subseteq A_1$  and  $N(Y_3) \subseteq B_1$ . Thus  $X_3 = \{x \in X \setminus (X_1 \cup X_2) : N(x) \subseteq A_1\}$  and  $Y_3 = \{y \in Y \setminus (Y_1 \cup Y_2) : N(y) \subseteq B_1\}$ .



Figure 1. A partial edge-coloring of G.

We denote the above set partition by  $\mathcal{P}$ . The following observation holds for  $\mathcal{P}$  since G is bridgeless.

**Lemma 3.** (1) For  $x \in A_3$ ,  $N(x) \cap A_1 \neq \emptyset$ .

- (2) For  $x \in B_3$ ,  $N(x) \cap B_1 \neq \emptyset$ .
- (3) For  $x \in X_4$ ,  $N(x) \cap X_1 \neq \emptyset$ .
- (4) For  $x \in Y_4$ ,  $N(x) \cap Y_1 \neq \emptyset$ .

We give a partial 9-edge-coloring of G as follows:

$$c(e) = \begin{cases} 1, & \text{if } e = uv; \\ 2, & \text{if } e \in E[u, A_3] \cup E[v, B_1]; \\ 3, & \text{if } e \in E[u, A_1] \cup E[v, B_3]; \\ 4, & \text{if } e \in E[A_1, X_1 \cup Z] \cup E(G[A_1]); \\ 5, & \text{if } e \in E[B_1, Y_1 \cup Z] \cup E(G[B_1]); \\ 6, & \text{if } e \in E[A_1, B_1] \cup E[Z, K] \cup E[X_1, Z \cup I \cup W \cup Y_1]; \\ 7, & \text{if } e \in E[Z, I] \cup E[Y_1, K \cup W \cup Z]; \\ 8, & \text{if } e \in E[A_1, A_3] \cup E[B_1, B_3] \cup E[X_1, X_4] \\ & \cup E[Y_1, Y_4] \cup E[J, I \cup K \cup W]; \\ 9, & \text{if } e \in E[A_1, X_4] \cup E[B_1, Y_4]. \end{cases}$$

See Figure 1 for details.

For each  $x \in X_3$ ,  $N(x) \subseteq A_1$  by the above set partition. Since G is a bridgeless graph,  $|N(x)| \geq 2$ . Thus, we can color one edge incident to x by 8, and color the others incident to x by 9. Similarly, for each vertex  $y \in Y_3$ , we can color edges incident to y by colors 8 and 9.

**Lemma 4.** (1) For  $x \in X_1$ , there exists an  $x \stackrel{6}{\sim} Y_1 \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path, or  $x \stackrel{6}{\sim} Z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path, or  $x \stackrel{6}{\sim} I \stackrel{7}{\sim} Z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path, or  $x \stackrel{6}{\sim} W \stackrel{7}{\sim} Y_1 \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path under the above partial edge-coloring.

(2) For  $y \in Y_1$ , there exists a  $y \stackrel{6}{\sim} X_1 \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path, or  $y \stackrel{7}{\sim} Z \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path, or  $y \stackrel{7}{\sim} W \stackrel{6}{\sim} X_1 \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path, or  $y \stackrel{7}{\sim} K \stackrel{6}{\sim} Z \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path under the above partial edge-coloring.

**Proof.** We only show (1) since the proofs are similar. For any  $x \in X_1$ , by the definition of set  $X_1$ , we know that x has a neighbor, say x', in  $Y \cup Z \cup I \cup W$ .

If  $x' \in Y$ , then  $x' \in Y_1$  by the definition of set  $Y_1$ . Thus xx'x''v is an  $x \stackrel{6}{\sim} Y_1 \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path under the above partial edge-coloring, where x'' is a neighbor of x' in  $B_1$ .

If  $x' \in Z$ , then xx'x''v is an  $x \stackrel{6}{\sim} Z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path under the above partial edge-coloring, where x'' is a neighbor of x' in  $B_1$ .

If  $x' \in I$ , then xx'x''x'''v is an  $x \stackrel{6}{\sim} I \stackrel{7}{\sim} Z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path under the above partial edge-coloring, where x'' is a neighbor of x' in Z and x''' is a neighbor of x'' in  $B_1$ .

Otherwise,  $x' \in W$ , and then xx'x''x'''v is an  $x \stackrel{6}{\sim} W \stackrel{7}{\sim} Y_1 \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path under the above partial edge-coloring, where x'' is a neighbor of x' in  $Y_1$  and x''' is a neighbor of x'' in  $B_1$ .

**Lemma 5.** (1) For  $x \in A_1$ , there exists an  $x \stackrel{6}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path, or  $x \stackrel{4}{\sim} Z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path, or  $x \stackrel{4}{\sim} X_1 \stackrel{6}{\sim} Y_1 \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path, or  $x \stackrel{4}{\sim} X_1 \stackrel{6}{\sim} Z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path, or  $x \stackrel{4}{\sim} X_1 \stackrel{6}{\sim} I \stackrel{7}{\sim} Z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path, or  $x \stackrel{4}{\sim} X_1 \stackrel{6}{\sim} I \stackrel{7}{\sim} Z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path, or  $x \stackrel{4}{\sim} X_1 \stackrel{6}{\sim} I \stackrel{7}{\sim} Z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path, or  $x \stackrel{4}{\sim} X_1 \stackrel{6}{\sim} V \stackrel{7}{\sim} V_1 \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path under the above partial edge-coloring.

(2) For  $y \in B_1$ , there exists a  $y \stackrel{6}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path, or  $y \stackrel{5}{\sim} Z \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path, or  $y \stackrel{5}{\sim} Y_1 \stackrel{6}{\sim} X_1 \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path, or  $y \stackrel{5}{\sim} Y_1 \stackrel{7}{\sim} Z \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path, or  $y \stackrel{5}{\sim} Y_1 \stackrel{7}{\sim} Z \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path, or  $y \stackrel{5}{\sim} Y_1 \stackrel{7}{\sim} K \stackrel{6}{\sim} Z \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path, or  $y \stackrel{5}{\sim} Y_1 \stackrel{7}{\sim} K \stackrel{6}{\sim} Z \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path or  $y \stackrel{5}{\sim} Y_1 \stackrel{7}{\sim} K \stackrel{6}{\sim} Z \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path under the above partial edge-coloring.

**Proof.** We only show (1) since the proofs are similar. For any  $x \in A_1$ , by the definition of set  $A_1$ , we know that x has a neighbor, say, x', in  $B_1 \cup Z \cup X_1$ .

If  $x' \in B_1$ , then xx'v is an  $x \stackrel{6}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path under the above partial edge-coloring.

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If  $x' \in Z$ , then xx'x''v is an  $x \stackrel{4}{\sim} Z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v$ -rainbow path, where x'' is a neighbor of x' in  $B_1$ .

Otherwise,  $x' \in X_1$ . By Lemma 4, there exists a desired rainbow path.

**Lemma 6.** (1) For  $x \in Z$ , there exists an  $x \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v \stackrel{1}{\sim} u \stackrel{3}{\sim} A_1 \stackrel{4}{\sim} x$ -rainbow cycle under the above partial edge-coloring.

(2) For  $x \in I$ , there exists an  $x \stackrel{7}{\sim} Z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v \stackrel{1}{\sim} u \stackrel{3}{\sim} A_1 \stackrel{4}{\sim} X_1 \stackrel{6}{\sim} x$ -rainbow cycle under the above partial edge-coloring.

(3) For  $x \in K$ , there exists an  $x \stackrel{7}{\sim} Y_1 \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v \stackrel{1}{\sim} u \stackrel{3}{\sim} A_1 \stackrel{4}{\sim} Z \stackrel{6}{\sim} x$ -rainbow cycle under the above partial edge-coloring.

(4) For  $x \in W$ , there exists an  $x \stackrel{7}{\sim} Y_1 \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v \stackrel{1}{\sim} u \stackrel{3}{\sim} A_1 \stackrel{4}{\sim} X_1 \stackrel{6}{\sim} x$ -rainbow cycle under the above partial edge-coloring.

**Proof.** We only show (4) since (1), (2) and (3) can be proved similarly. For any  $x \in W$ , by the definition of set W, the vertex x has a neighbor  $v_1 \in X_1$  and a neighbor  $v_2 \in Y_1$ . Moreover, by the definitions of sets  $X_1$  and  $Y_1$ , the vertex  $v_1$  has a neighbor  $v_3 \in A_1$ , and the vertex  $v_2$  has a neighbor  $v_4 \in B_1$ . Thus  $x \stackrel{7}{\sim} v_2 \stackrel{5}{\sim} v_4 \stackrel{2}{\sim} v \stackrel{1}{\sim} u \stackrel{3}{\sim} v_3 \stackrel{4}{\sim} v_1 \stackrel{6}{\sim} x$  is a rainbow cycle, that is, there exists an  $x \stackrel{7}{\sim} Y_1 \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v \stackrel{1}{\sim} u \stackrel{3}{\sim} A_1 \stackrel{4}{\sim} X_1 \stackrel{6}{\sim} x$ -rainbow cycle under the above partial edge-coloring.

**Lemma 7.** For any two vertices  $x, y \in V(G) \setminus (A_2 \cup B_2 \cup X_2 \cup Y_2 \cup J)$ , there exists a rainbow path joining x and y under the above partial edge-coloring.

**Proof.** Let x and y be any two vertices in  $V(G) \setminus (A_2 \cup B_2 \cup X_2 \cup Y_2 \cup J)$ . It is easy to see that there exists a rainbow path between u (respectively v) and another vertex  $w \in V(G) \setminus (A_2 \cup B_2 \cup X_2 \cup Y_2 \cup J)$  in the partial edge-color graph G. Thus suppose that  $\{u, v\} \cap \{x, y\} = \emptyset$ .

Case 1.  $x, y \in A_1 \cup B_1 \cup X_1 \cup Y_1 \cup Z \cup I \cup K \cup W$ . By Lemmas 4, 5 and 6, we can pick a special rainbow path  $P_1$  between x and v and a special rainbow path  $P_2$  between y and v such that  $c(P_1) \cap c(P_2) = \emptyset$ . Thus we can obtain a rainbow path joining x and y by combining the paths  $P_1$  and  $P_2$ .

Case 2. Exactly one of x and y belongs to  $A_1 \cup B_1 \cup X_1 \cup Y_1 \cup Z \cup I \cup K \cup W$ . Without loss of generality, say  $x \in A_1 \cup B_1 \cup X_1 \cup Y_1 \cup Z \cup I \cup K \cup W$  and  $y \in A_3 \cup B_3 \cup X_3 \cup X_4 \cup Y_3 \cup Y_4$ . We only check the case  $y \in A_3 \cup X_3 \cup X_4$  since the case  $y \in B_3 \cup Y_3 \cup Y_4$  can be checked similarly.

For  $y \in A_3 \cup X_3 \cup X_4$ , there exists a  $y \stackrel{9 \text{ (or 8)}}{\sim} A_1 \stackrel{3}{\sim} u \stackrel{1}{\sim} v$ -rainbow path  $P_1$  joining y and v. Moreover, there exists a  $\{2, 4, 5, 6, 7\}$ -rainbow path  $P_2$  joining x and v. Thus a rainbow path joining x and y can be obtained from  $P_1$  and  $P_2$ .

Case 3.  $x \in A_3 \cup X_3 \cup X_4$  and  $y \in B_3 \cup Y_3 \cup Y_4$ .

Subcase 3.1.  $x \in A_3$ . There exist an  $x \stackrel{2}{\sim} u$ -rainbow path  $P_1$  and an  $x \stackrel{8}{\sim} A_1$  $\stackrel{3}{\sim} u$ -rainbow path  $P_2$  by Figure 1 and Lemma 3.

If  $y \in Y_3 \cup Y_4$ , then there exists a  $y \stackrel{9}{\sim} B_1 \stackrel{2}{\sim} v \stackrel{1}{\sim} u$ -rainbow path  $P_3$ . Thus a rainbow path joining x and y can be obtained from  $P_2$  and  $P_3$ .

If  $y \in B_3$ , then there exists a  $y \stackrel{3}{\sim} v \stackrel{1}{\sim} u$ -rainbow path  $P_4$ . Thus a rainbow path joining x and y can be obtained from  $P_1$  and  $P_4$ .

Subcase 3.2.  $x \in X_3 \cup X_4$ . There exists an  $x \stackrel{9}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path  $P_1$  by Figure 1. Moreover, there exists a  $y \stackrel{8}{\sim} B_1 \stackrel{2}{\sim} v \stackrel{1}{\sim} u$ -rainbow path  $P_2$  if  $y \in B_3 \cup Y_3$ , or there exists a  $y \stackrel{8}{\sim} Y_1 \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v \stackrel{1}{\sim} u$ -rainbow path  $P_2$  if  $y \in Y_4$ . Thus a rainbow path joining x and y can be obtained from  $P_1$  and  $P_2$ .

Case 4.  $x, y \in A_3 \cup X_3 \cup X_4$  or  $x, y \in B_3 \cup Y_3 \cup Y_4$ . We only check the case  $x, y \in A_3 \cup X_3 \cup X_4$  since the case  $x, y \in B_3 \cup Y_3 \cup Y_4$  can be checked similarly.

Subcase 4.1.  $x \in A_3$  or  $y \in A_3$ . Without loss of generality, say  $x \in A_3$ . Then there exists a  $x \stackrel{2}{\sim} u \stackrel{3}{\sim} A_1 \stackrel{8(\text{or } 9)}{\sim} y$ -rainbow path connecting x and y.

Subcase 4.2. At least one of x and y belongs to  $X_3$ . Without loss of generality, assume that  $x \in X_3$ . Let x' and y' be neighbors of x and y in  $A_1$  such that c(xx') = 8 and c(yy') = 9. By Lemma 5, there exists a  $\{2, 4, 5, 6, 7\}$ -rainbow path P joining y' and v. Thus yy'Pvux'x is a rainbow path connecting x and y.

Subcase 4.3. Both x and y belong to  $X_4$ . Let x' be a neighbor of x in  $A_1$ , and let y' be a neighbor of y in  $X_1$ . By Lemma 4, there exists a  $\{2, 5, 6, 7\}$ -rainbow path P joining y' and v. Thus yy'Pvux'x is a rainbow path connecting x and y.

# 3. A Complete Edge-Coloring

To complete our edge-coloring, we further partition J as follows:

 $J_0 = \{x \in J : x \text{ is not an isolated vertex in } G[J]\},$   $J_1 = \{x \in J \setminus J_0 : x \text{ has at least a neighbor in } K\},$   $J_2 = \{x \in J \setminus (J_0 \cup J_1) : x \text{ has at least a neighbor in } W\},$   $J_3 = \{x \in J \setminus (J_0 \cup J_1 \cup J_2) : x \text{ has at least a neighbor in } I\},$  $J_4 = J \setminus (J_0 \cup J_1 \cup J_2 \cup J_3).$ 

Now we further color the edges of G as follows: color the edges in  $E[Z, J_1 \cup J_2 \cup J_3]$  by color 7; for any  $x \in J_4$ , color one in E[x, Z] by 8, color the others in E[x, Z] by 9 (there exists at least one such edge since G is bridgeless).

To color the remaining edges, we need the following lemma.

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Figure 2. A complete edge-coloring of G (we omit the line between Z and  $J_1$ , the line between Z and  $J_2$ , and the line between Z and  $J_3$ ).

**Lemma 8.** Let S and T be two disjoint vertex sets of a graph G such that  $S \subseteq N(T)$ . If the induced subgraph G[S] has no trivial components, then there is an  $\{\alpha, \beta, \gamma\}$ -edge-coloring of  $G[S] \cup E[S, T]$  such that there exist two rainbow paths  $P_1$  and  $P_2$  joining s and T for every  $s \in S$ . Furthermore, if  $P_1$  has color  $\{\alpha\}$ , then  $P_2$  has colors  $\{\beta, \gamma\}$ ; if  $P_1$  has color  $\{\beta\}$ , then  $P_2$  has colors  $\{\alpha, \gamma\}$ .

**Proof.** Let F be a maximal spanning forest of G[S], and let (X, Y) be any of bipartitions defined by this forest F. We give a 3-edge-coloring  $c : E(G[S]) \cup E[S,T] \to \{\alpha, \beta, \gamma\}$  of G by defining

$$c(e) = \begin{cases} \alpha, & \text{if } e \in E[T, X]; \\ \beta, & \text{if } e \in E[T, Y]; \\ \gamma, & \text{otherwise.} \end{cases}$$

Clearly, for the edge-coloring above, there exist two rainbow paths  $P_1$  and  $P_2$  joining s and T for every  $s \in S$ . Furthermore, if  $P_1$  has color  $\{\alpha\}$ , then  $P_2$  has colors  $\{\beta, \gamma\}$ ; if  $P_2$  has color  $\{\beta\}$ , then  $P_2$  has colors  $\{\alpha, \gamma\}$ .

**Remark.** The edge-coloring in Lemma 8 is called an  $\langle \alpha, \beta, \gamma \rangle$ -edge-coloring for T and  $X \cup Y$ . Let  $T_{A_2}, T_{B_2}, T_{X_2}, T_{Y_2}$  and  $T_{J_0}$  be maximal spanning forests of  $G[A_2], G[B_2], G[X_2], G[Y_2]$  and  $G[J_0]$ , respectively. Clearly, the forests have no isolated vertex. Let  $A_2^0$  and  $A_2^1$ ,  $B_2^0$  and  $B_2^1$ ,  $X_2^0$  and  $X_2^1$ ,  $Y_2^0$  and  $Y_2^1$ , and  $J_0^0$  and  $J_0^1$  be bipartitions of  $T_{A_2}, T_{B_2}, T_{X_2}, T_{Y_2}$  and  $T_{J_0}$ . Now we give a  $\langle 2, 3, 8 \rangle$ -edge-coloring for u and  $A_2^0 \cup A_2^1$ , a  $\langle 2, 3, 8 \rangle$ -edge-coloring for v and  $B_2^0 \cup B_2^1$ , an  $\langle 8, 9, 7 \rangle$ -edge-coloring for  $A_1$  and  $X_2^0 \cup X_2^1$ , an  $\langle 8, 9, 7 \rangle$ -edge-coloring for Z and  $J_0^0 \cup J_0^1$  as shown in Figure 2.

Furthermore, we color the edges in subgraphs  $G[A_1]$ ,  $G[X_2^0]$  and  $G[X_2^1]$  by 4, the edges in subgraphs  $G[B_1]$ ,  $G[Y_2^0]$  and  $G[Y_2^1]$  by 5, the edges in  $E[X_1, X_2^1]$  and  $E[Y_1, Y_2^1]$  by 8, and the edges in  $E[X_1, X_2^0]$  and  $E[Y_1, Y_2^0]$  by 9.

For the remaining edges, we can color them arbitrarily. Up to now, we give the graph G a complete edge-coloring. Let  $\mathcal{P}$  be our final vertex set partition and let c be our final edge-coloring.

**Lemma 9.** For any two vertices  $x \in A_2 \cup B_2 \cup X_2 \cup Y_2 \cup J$  and  $y \in V(G) \setminus (A_2 \cup B_2 \cup X_2 \cup Y_2 \cup J)$ , there exists a rainbow path under the above partial edge-coloring.

**Proof.** We consider the following three cases.

Case 1.  $x \in A_2 \cup B_2$ . We only consider the case  $x \in A_2$  since the case  $x \in B_2$  can be checked similarly.

Subcase 1.1.  $x \in A_2^0$ . By observing Figure 2, there exist an  $x \stackrel{2}{\sim} u$ -rainbow path  $P_1$  joining x and u, or an  $x \stackrel{8}{\sim} A_2^1 \stackrel{3}{\sim} u$ -rainbow path  $P_2$  joining x and u.

If  $y \in A_3$ , then  $P_2 y$  is a rainbow path joining x and y.

If  $y \in B_3$ , then  $P_1vy$  is a rainbow path joining x and y.

If  $y \in B_1 \cup Y_1 \cup Y_3 \cup Y_4 \cup Z \cup I \cup K \cup W$ , then there exists a  $\{1, 2, 5, 6, 7, 9\}$ -rainbow path  $Q_1$  joining u and y. Thus a rainbow path joining x and y can be obtained by combining  $P_2$  and  $Q_1$ .

If  $y \in A_1 \cup X_1 \cup X_3 \cup X_4$ , then there exists a  $\{3, 4, 9\}$ -rainbow path  $Q_2$  joining u and y. Thus a rainbow path joining x and y can be obtained by combining  $P_1$  and  $Q_2$ .

Subcase 1.2.  $x \in A_2^1$ . By observing Figure 2, there exist an  $x \stackrel{3}{\sim} u$ -rainbow path  $P_1$  joining x and u, or an  $x \stackrel{8}{\sim} A_2^0 \stackrel{2}{\sim} u$ -rainbow path  $P_2$  joining x and u.

If  $y \in A_3$ , then  $P_1 y$  is a rainbow path joining x and y.

If  $y \in B_3$ , then  $P_2vy$  is a rainbow path joining x and y.

If  $y \in B_1 \cup Y_1 \cup Y_3 \cup Y_4 \cup Z \cup I \cup K \cup W$ , then there exists a  $\{1, 2, 5, 6, 7, 9\}$ -rainbow path  $Q_1$  joining u and y. Thus a rainbow path joining x and y can be obtained by combining  $P_1$  and  $Q_1$ .

If  $y \in A_1 \cup X_1 \cup X_3 \cup X_4$ , then there exists a  $\{3, 4, 9\}$ -rainbow path  $Q_2$  joining u and y. Thus a rainbow path joining x and y can be obtained by combining  $P_2$  and  $Q_2$ .

Case 2.  $x \in X_2 \cup Y_2$ . We only consider the case  $x \in X_2$  since the case  $x \in Y_2$  can be checked similarly.

Subcase 2.1.  $x \in X_2^0$ . By observing Figure 2, there exists an  $x \stackrel{8}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path  $P_1$  joining x and u.

If  $y \in A_3$ , then  $P_1 y$  is a rainbow path joining x and y.

If  $y \in B_3$ , then  $x \stackrel{7}{\sim} X_2^1 \stackrel{9}{\sim} A_1 \stackrel{3}{\sim} u \stackrel{1}{\sim} v \stackrel{2}{\sim} B_1 \stackrel{8}{\sim} y$  is a rainbow path joining x and y.

If  $y \in B_1 \cup Y_1 \cup Y_3 \cup Y_4 \cup Z \cup I \cup K \cup W$ , then there exists a  $\{1, 2, 5, 6, 7, 9\}$ -rainbow path  $Q_1$  joining v and y. Thus a rainbow path joining x and y can be obtained by combining  $P_1$  and  $Q_1$ .

If  $y \in A_1 \cup X_1$ , then there exists a  $\{2, 4, 5, 6, 7\}$ -rainbow path  $Q_1$  joining v and y by Lemmas 4 and 5. Thus a  $\{1, 2, 3, 4, 5, 6, 7, 8\}$ -rainbow path joining x and y can be obtained by combining  $P_1$ ,  $Q_1$  and edge uv.

If  $y \in X_3 \cup X_4$ , then y has a neighbor y' in  $A_1$  such that c(yy') = 9. Note that there exists a  $\{1, 2, 3, 4, 5, 6, 7, 8\}$ -rainbow path P joining x and y' by the arguments of the above paragraph. Thus Py is a rainbow path joining x and y.

Subcase 2.2.  $x \in X_2^1$ . By observing Figure 2, there exist an  $x \stackrel{9}{\sim} A_1 \stackrel{3}{\sim} u$ -rainbow path  $P_1$  joining x and u.

If  $y \in A_3$ , then  $P_1 y$  is a rainbow path joining x and y.

If  $y \in B_3$ , then  $P_1vy'y$  is a rainbow path joining x and y, where y' is a neighbor of y in  $B_1$ .

If  $y \in B_1 \cup Y_1 \cup Y_3 \cup Y_4 \cup Z \cup I \cup K \cup W$ , then there exists a  $\{1, 2, 5, 6, 7, 8\}$ -rainbow path  $Q_1$  joining u and y. Thus a rainbow path joining x and y can be obtained by combining  $P_1$  and  $Q_1$ .

If  $y \in A_1 \cup X_1$ , then there exists a  $\{2, 4, 5, 6, 7\}$ -rainbow path  $Q_1$  joining v and y by Lemmas 4 and 5. Thus a  $\{1, 2, 3, 4, 5, 6, 7, 9\}$ -rainbow path joining x and y can be obtained by combining  $P_1$ ,  $Q_1$  and edge uv.

If  $y \in X_3 \cup X_4$ , then y has a neighbor y' in  $A_1$  or  $X_1$  such that c(yy') = 8. Note that there exists a  $\{1, 2, 3, 4, 5, 6, 7, 9\}$ -rainbow path P joining x and y' by the arguments of the above paragraph. Thus Py is a rainbow path joining x and y.

Case 3.  $x \in J$ . By observing Figure 2, there exists a  $\{7,9\}$ -rainbow path P joining x and some vertex  $z \in Z$ . Furthermore, there exist a  $z \stackrel{4}{\sim} A_1 \stackrel{3}{\sim} u \stackrel{1}{\sim} v$ -rainbow path  $Q_1$  joining z and v, and a  $z \stackrel{5}{\sim} B_1 \stackrel{2}{\sim} v \stackrel{1}{\sim} u$ -rainbow path  $Q_2$  joining z and u. Thus a  $\{1,3,4,7,9\}$ -rainbow path  $Q'_1$  joining x and v can be obtained from P and  $Q_1$ , and a  $\{1,2,5,7,9\}$ -rainbow path  $Q'_2$  joining x and u can be obtained from P and  $Q_2$ .

If  $y \in B_1 \cup B_3 \cup Y_1 \cup Y_3 \cup Y_4$ , then there exists a  $\{2, 5, 8\}$ -rainbow path  $R_1$  between v and y. Thus a rainbow path joining x and y can be obtained from  $Q'_1$  and  $R_1$ .

If  $y \in A_1 \cup A_3 \cup X_1 \cup X_3 \cup X_4 \cup Z \cup I \cup K \cup W$ , then there exists a  $\{3, 4, 6, 8\}$ -rainbow path  $R_2$  between u and y. Thus a rainbow path joining x and y can be obtained from  $Q'_2$  and  $R_2$ .

#### 4. 9-RAINBOW-CONNECTED EDGE-COLORING

In this section, we check that the above 9-edge-coloring is rainbow-connected 9edge-coloring. It suffices to check that for any two vertices  $x, y \in A_2 \cup B_2 \cup X_2 \cup Y_2 \cup J$ , there exists a rainbow path under the above partial edge-coloring.

**Lemma 10.** There exists a rainbow path joining any two vertices of  $X_2$  under the edge-coloring c.

**Proof.** Let x and y be any two vertices in  $X_2$ . We consider the following two cases.

Case 1.  $x \in X_2^0$  and  $y \in X_2^1$ , or  $x \in X_2^0$  and  $y \in X_2^1$ . Without loss of generality, assume that  $x \in X_2^0$  and  $y \in X_2^1$ . Let x' and y' be neighbors of x and y in  $A_1$ , respectively. By Figure 2, we know that c(xx') = 8 and c(yy') = 9. By Lemma 5, there exists a  $\{2, 4, 5, 6, 7\}$ -rainbow path  $P_{y',v}$ . Thus, a  $\{1, 2, 3, 4, 5, 6, 7, 8, 9\}$ -rainbow path joining x and y is obtained from the edge yy', rainbow paths  $P_{y',v}$  and vux'x.

Case 2.  $x, y \in X_2^0$  or  $x, y \in X_2^1$ . We only check the case  $x, y \in X_2^0$  since the case  $x, y \in X_2^1$  can be checked similarly.

Subcase 2.1.  $d(x, B_1) = 2$  or  $d(y, B_1) = 2$ . Without loss of generality, assume  $d(x, B_1) = 2$ . Let  $x' \in N(x) \cap N(B_1)$ . By the definition of the above set partition, we know  $x' \in A_1$ . So, xx'x''vu is a  $\{1, 2, 6, 8\}$ -rainbow path, where x'' is a neighbor of x' in  $B_1$ . By Figure 2 and Lemma 8, u and y are connected by a  $\{3, 7, 9\}$ -rainbow path P. Thus a  $\{1, 2, 3, 6, 7, 8, 9\}$ -rainbow path joining x and y can be obtained from rainbow paths xx'x''vu and P.

Subcase 2.2.  $d(x, B_1) = d(y, B_1) = 3$ . Let  $xx_1x_2x_3$  be a path joining x and some vertex  $x_3 \in B_1$ . By the set partition above,  $x_1 \in A_1 \cup X_1 \cup X_2$ .

Subsubcase 2.2.1.  $x_1 \in A_1$ . By the definition of  $\mathcal{P}$ ,  $x_2 \in A_1 \cup B_1 \cup Z$ . So  $xx_1x_2x_3$  is a  $\{4, 5, 6\}$ -rainbow path. Furthermore,  $xx_1x_2x_3vu$  is  $\{1, 2, 4, 5, 6, 8\}$ -rainbow. By Figure 2, there exists a  $\{3, 7, 9\}$ -rainbow path P joining u and y. Hence a rainbow path joining x and y can be obtained from rainbow paths  $xx_1x_2x_3vu$  and P.

Subsubcase 2.2.2.  $x_1 \in X_1$ . By the definition of the above set partition,  $x_2 \in A_1 \cup Z$ . Thus  $xx_1x_2x_3$  is a  $\{4, 5, 6, 9\}$ -rainbow path. Thus  $xx_1x_2x_3vuy'y$  is a  $\{1, 2, 3, 4, 5, 6, 8, 9\}$ -rainbow path joining x and y, where  $y_1$  is a neighbor of yin  $A_1$ . Subsubcase 2.2.3.  $x_1 \in X_2$ . If  $x_1 \in X_2^0$ , then  $c(xx_1) = 4$ . Furthermore,  $x_2 \in A_1$ . Thus  $xx_1x_2x_3vu$  is a  $\{1, 2, 4, 6, 8\}$ -rainbow path. By Figure 2, there exists a  $\{3, 7, 9\}$ -rainbow path P joining u and y. Hence a rainbow path joining x and y can be obtained from  $xx_1x_2x_3vu$  and P.

If  $x_1 \in X_2^1$ , then  $c(xx_1) = 7$ . Furthermore,  $x_2 \in A_1$ . Thus  $xx_1x_2x_3vu$  is a  $\{1, 2, 4, 6, 7, 9\}$ -rainbow path. By Figure 2, there exists a  $\{3, 8\}$ -rainbow path P joining u and y. Hence a rainbow path joining x and y can be obtained from  $xx_1x_2x_3vu$  and P.

Similarly to Lemma 8, the following lemma holds.

**Lemma 11.** There exists a rainbow path joining any two vertices of  $Y_2$  under the edge-coloring above.

**Lemma 12.** For any two vertices  $x, y \in A_2 \cup B_2 \cup X_2 \cup Y_2 \cup J$ , there exists a rainbow path under the above partial edge-coloring.

**Proof.** For  $x, y \in X_2$  or  $x, y \in Y_2$ , there exists a rainbow path joining x and y by Lemmas 10 or 11. For the others, we can easily check them by Lemmas 4, 5, 6 and 8 in a similar way.

Combining Lemmas 7, 9 and 12, we have the following result.

**Theorem 13.** Let G be a bridgeless graph with diameter 3. If there exists an edge e such that e does not belongs to any triangle in G, then  $rc(G) \leq 9$ .

For a bridgeless graph G with diameter 3, if each edge belongs to a triangle in G, then  $rc(G) \leq 9$  by Theorem 1. Combining this result with Theorem 13, we know that Theorem 2 holds.

We can give the following example of graphs with diameter 3 for which the rainbow connection number reaches 7.

**Example 2.** Let  $K_n$  be a complete graph with vertex set  $\{v_1, \ldots, v_n\}$ , where  $n \ge 217$ . For every  $v_i$ , we add a pendant path  $\langle v_i, v_{i,1}, v_{i,2}, v_{i,3} \rangle$ , denoted by  $P_i$ , and then we identify the vertex  $v_{i,3}$  with a vertex v. The resulting graph is denoted by G. Clearly, diam(G) = 3. Let c be any 6-edge-coloring of G with colors  $\{1, \ldots, 6\}$ . Since  $6^3 = 216$ , at least two of them are colored the same. Without loss generality, say  $P_1$  and  $P_2$ , that is,  $c(v_1v_{1,1}) = c(v_2v_{2,1}), c(v_{1,1}v_{1,2}) = c(v_{2,1}v_{2,2})$  and  $c(v_{1,2}v) = c(v_{2,2}v)$ . By the structure of G, it is easy to see that there exists no rainbow path joining  $v_{1,1}$  and  $v_{2,1}$  in G under c. Thus  $rc(G) \ge 7$ .

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