

NOTE

## SMALLEST REGULAR GRAPHS OF GIVEN DEGREE AND DIAMETER

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### Abstract

In this note we present a sharp lower bound on the number of vertices in a regular graph of given degree and diameter.

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### 1. INTRODUCTION

The *degree/diameter* problem consists in determination of the largest order  $N(d, k)$  of a graph with (maximum) degree  $d$  and diameter  $k$ . An upper bound for  $N(d, k)$  is the *Moore bound*  $M(d, k) = 1 + d + d(d - 1) + \cdots + d(d - 1)^{k-1}$  and graphs achieving this bound are called *Moore graphs*. As shown in [1, 3, 5], Moore graphs exist only when  $d = 2$  or  $k = 1$  or when  $k = 2$  and the degree is either 3 or 7 or possibly 57. For all other pairs  $(d, k)$  we have  $N(d, k) \leq M(d, k) - 2$ , see [2, 4]. Recently, there are plenty of papers dealing with the degree/diameter problem, some of them constructing “large” graphs of given degree and diameter, which increases the lower bound for  $N(d, k)$  for special pairs  $(d, k)$ , other decreasing  $N(d, k)$  for special classes of graphs. For a nice survey see [7].

In this note we consider the inverse of degree/diameter problem. Since usually the degree/diameter problem is formulated for regular graphs (although some authors require only that  $d$  is the maximum degree), we ask what is the minimum

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order  $n(d, k)$  of a regular graph of degree  $d$  and diameter  $k$ . In this note we answer this question completely.

We start with some notation. Let  $G$  be a graph,  $G = (V(G), E(G))$ . For two of its vertices, say  $x$  and  $y$ , by  $\text{dist}_G(x, y)$  we denote their distance in  $G$ . By  $N_i(x)$  we denote the set of vertices that are at distance  $i$  from  $x$ . As usual,  $N_1(x)$  is often abbreviated to  $N(x)$ . The longest distance in  $G$  is the *diameter*  $\text{diam}(G)$ . The complete graph on  $n$  vertices is denoted by  $K_n$  and the discrete graph on  $n$  vertices (the complement of  $K_n$ ) is denoted by  $D_n$ . If  $G$  is a graph, then by  $G^{(-1)}$  (and  $G^{(-2)}$ ) we denote a graph obtained from  $G$  by removing all the edges of one 1-factor (one 2-factor).

If  $G$  and  $H$  are graphs, then  $G + H$  denotes the *join* of  $G$  and  $H$ , that is, a graph obtained from the disjoint union of  $G$  and  $H$  by adding all edges  $xy$ , where  $x \in V(G)$  and  $y \in V(H)$ . The *sequential join* of graphs  $G_1, G_2, \dots, G_r$  is denoted by  $G_1 + G_2 + \dots + G_r$  and is defined by

$$G_1 + G_2 + \dots + G_r = (G_1 + G_2) \cup (G_2 + G_3) \cup \dots \cup (G_{r-1} + G_r).$$

Thus, one can obtain  $G_1 + G_2 + \dots + G_r$  from the disjoint union  $G_1 \cup G_2 \cup \dots \cup G_r$  by adding all edges  $xy$  where  $x \in V(G_i)$  and  $y \in V(G_{i+1})$  for  $i = 1, 2, \dots, r-1$ . To simplify the expressions, instead of

$$\dots + \underbrace{G + G + \dots + G}_{k \text{ times}} + \dots \quad \text{we write} \quad \dots + (G)_k + \dots$$

Finally, denote by  $G \div H$  a graph obtained from the disjoint union of  $G$  and  $H$  by adding all edges of one 1-factor, every edge of which joins a vertex of  $G$  with a vertex of  $H$ . Obviously,  $G \div H$  is defined only if  $|V(G)| = |V(H)|$ . Analogously as in the case of join, by  $G_1 \div G_2 \div \dots \div G_r$  we denote the graph  $(G_1 \div G_2) \cup (G_2 \div G_3) \cup \dots \cup (G_{r-1} \div G_r)$ . We can form also more complicated expressions using both  $+$  and  $\div$ . In such a way,  $K_1 + D_2 \div D_2 \div K_2$  is a cycle of length 7; see Figure 1.

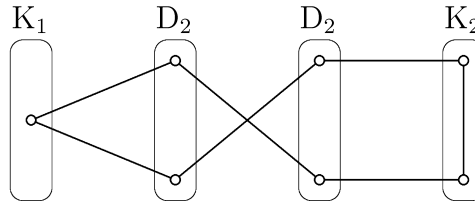


Figure 1. The graph  $K_1 + D_2 \div D_2 \div K_2$ .

## 2. RESULTS

For small diameters we have the following statement.

**Proposition 1.** *Let  $d \geq 2$ . We have*

- (i)  $n(d, 1) = d + 1$ ;
- (ii) *if  $d$  is even, then  $n(d, 2) = d + 2$ ;*
- (iii) *if  $d$  is odd, then  $n(d, 2) = d + 3$ ;*
- (iv)  $n(d, 3) = 2d + 2$ .

**Proof.** The case  $k = 1$  is obvious since  $K_{d+1}$  is the unique graph of diameter 1 and degree  $d$ .

Let  $k = 2$ . Let  $G$  be a  $d$ -regular graph of diameter 2, and let  $x, y \in V(G)$  such that  $\text{dist}_G(x, y) = 2$ . Then  $\{x\} \cup N(x) = N_0(x) \cup N_1(x)$ , which gives  $|N_0(x)| + |N_1(x)| = d + 1$ . Since  $y \in N_2(x)$ , we have  $|V(G)| = |N_0(x)| + |N_1(x)| + |N_2(x)| \geq d + 2$ , which gives  $n(d, 2) \geq d + 2$ . However, if  $d$  is odd then  $|V(G)|$  cannot be odd and so  $n(d, 2) \geq d + 3$  in this case. If  $d$  is even then  $K_{d+2}^{(-1)}$  is a  $d$ -regular graph of diameter 2 on  $d + 2$  vertices, which shows  $n(d, 2) \leq d + 2$ ; while if  $d$  is odd then  $K_{d+3}^{(-2)}$  is a  $d$ -regular graph of diameter 2 on  $d + 3$  vertices, which shows  $n(d, 2) \leq d + 3$ .

Finally, let  $k = 3$ . Analogously as above, let  $G$  be a  $d$ -regular graph of diameter 3, and let  $x, y \in V(G)$  such that  $\text{dist}_G(x, y) = 3$ . Then  $\{x\} \cup N(x) = N_0(x) \cup N_1(x)$ , which gives  $|N_0(x)| + |N_1(x)| = d + 1$ , and  $\{y\} \cup N(y) \subseteq N_2(x) \cup N_3(x)$ , which gives  $|N_2(x)| + |N_3(x)| \geq d + 1$ . Thus,  $|V(G)| = |N_0(x)| + |N_1(x)| + |N_2(x)| + |N_3(x)| \geq 2d + 2$ , and so  $n(d, 3) \geq 2d + 2$ . On the other hand, denote by  $K_{n,n}$  a complete bipartite graph on  $2n$  vertices in which the two partite sets have  $n$  vertices each. Then  $K_{d+1, d+1}^{(-1)}$  is a  $d$ -regular graph of diameter 3 on  $2d + 2$  vertices, which shows  $n(d, 3) \leq 2d + 2$ . ■

Now we turn our attention to larger diameters. Since there are only two 2-regular graphs of diameter  $k$ , namely the cycle on  $2k$  vertices and the cycle on  $2k + 1$  vertices, we have the following trivial observation.

**Proposition 2.** *If  $k \geq 4$ , then  $n(2, k) = 2k$ .*

For larger degrees we have a slightly different bound.

**Theorem 3.** *Let  $k = 3j + t$ , where  $k \geq 4$  and  $0 \leq t \leq 2$ , and let  $d \geq 3$ . Then  $n(d, k) = (d + 1)(j + 1) + t + \delta$ , where  $\delta = 1$  if either  $d$  is odd and  $t = 1$  or  $d$  is even and  $t = 2$ . Otherwise  $\delta = 0$ .*

**Proof.** First we prove a lower bound for  $n(d, k)$ . Let  $G$  be a regular graph of degree  $d$  and diameter  $k$  and let  $x, y \in V(G)$  such that  $\text{dist}_G(x, y) = k$ . Denote  $n_i = |N_i(x)|$ . Since  $x \in N_0(x)$ , we have  $\{x\} \cup N(x) \subseteq N_0(x) \cup N_1(x)$ . Thus,  $n_0 + n_1 \geq d + 1$ . Analogously  $n_{k-1} + n_k \geq d + 1$  since  $y \in N_k(x)$ . Further, for every  $i$ ,  $1 \leq i \leq j-1$ , we have  $n_{3i-1} + n_{3i} + n_{3i+1} \geq d + 1$  since for  $z_i \in N_{3i}(x)$  it holds  $\{z_i\} \cup N(z_i) \subseteq N_{3i-1}(x) \cup N_{3i}(x) \cup N_{3i+1}(x)$ . Finally, if  $t \geq 1$  then  $n_{k-1-\ell} \geq 1$  where  $1 \leq \ell \leq t$ . Summing up all these inequalities we get

$$|V(G)| = \sum_{i=0}^k n_i \geq (d+1)(j+1) + t.$$

If  $t = 2$  then we use  $n_{k-3} \geq 1$  and  $n_{k-2} \geq 1$ . But if  $d$  is even then  $G$  cannot have a bridge, and so  $n_{k-3} + n_{k-2} \geq 3$ . Thus, we get  $|V(G)| = \sum_{i=0}^k n_i \geq (d+1)(j+1) + t + 1$  in this case.

Similarly, if  $t = 1$  and  $d$  is odd then  $(d+1)(j+1) + t$  is an odd number. But a regular graph of odd degree cannot have an odd number of vertices, and so  $|V(G)| = \sum_{i=0}^k n_i \geq (d+1)(j+1) + t + 1$  also in this case.

To prove the upper bound we construct extremal graphs, that is, regular graphs of degree  $d$  and diameter  $k$  on  $n(d, k)$  vertices. First we define an extremal graph  $G$  for odd  $d$ . The case  $k = 4$  is treated separately. If  $d = 3$  then one extremal graph  $G$  is on Figure 2. For  $d \geq 5$  we set  $G = K_2 + K_{d-1}^{(-2)} + D_2 \div D_2 + K_{d-1}$ .

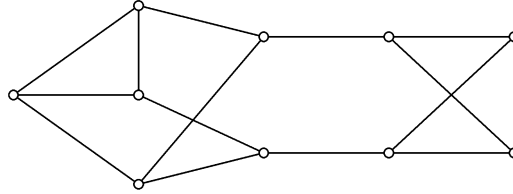


Figure 2. An extremal graph for  $d = 3$  and  $k = 4$ .

Recall that  $k = 3j + t$ . To cover the remaining diameters, that is, 5, 6, 7, ..., in the next we assume  $j \geq 1$  if  $t = 2$ , and  $j \geq 2$  if  $t = 0$  or  $t = 1$ :

$$G = K_2 + K_{d-1}^{(-1)} + (K_1 + K_1 + K_{d-1})_{j-1} + K_1 + K_1 + K_{d-1}^{(-1)} + K_2, \text{ if } t = 2;$$

$$G = K_2 + K_{d-1}^{(-1)} + (K_1 + K_1 + K_{d-1})_{j-2} + K_1 + K_1 + K_{d-1} \div K_{d-1}^{(-1)} + K_2, \text{ if } t = 0;$$

$$G = K_2 + K_{d-1}^{(-1)} + (K_1 + K_1 + K_{d-1})_{j-2} + K_1 + K_1 + K_{d-1}^{(-1)} + D_2 \div D_2 + K_{d-1}, \text{ if } t = 1.$$

Now we define an extremal graph  $G$  for even  $d$ . To cover all possible diameters, that is, 4, 5, 6, ..., in the next we assume  $j \geq 1$  if  $t = 1$  or  $t = 2$ , and  $j \geq 2$  if  $t = 0$ :

$$G = K_3 + K_{d-2}^{(-1)} + (K_1 + K_2 + K_{d-2})_{j-1} + K_1 + D_2 + K_{d-1}, \text{ if } t = 1;$$

$$G = K_3 + K_{d-2}^{(-1)} + (K_1 + K_2 + K_{d-2})_{j-1} + K_1 + K_2 + K_{d-2}^{(-2)} + K_3, \text{ if } t = 2;$$

$G = K_3 + K_{d-2}^{(-1)} + (K_1 + K_2 + K_{d-2})_{j-2} + K_1 + K_2 + K_{d-2} \div K_{d-2}^{(-1)} + K_3$ , if  $t = 0$ .

Observe that in all these graphs, whenever we removed a 1-factor out of  $K_q$ , then the number of vertices  $q$  was even. Obviously, in each case  $G$  has diameter  $k$  and it is a matter of routine to check that  $G$  is a regular graph of degree  $d$ . (For example, a vertex in the last copy of  $K_{d-2}^{(-1)}$  in the last graph is joined to 1 vertex of  $K_{d-2}$ ,  $d-4$  vertices of  $K_{d-2}^{(-1)}$  and to 3 vertices of  $K_3$ , so its degree is  $1 + d - 4 + 3 = d$ .) Also, in each of these cases the number of vertices of  $G$  attains the bound of the theorem. To verify this statement it suffices to check the number of vertices for the smallest admissible values of  $j$  since in each case in the brackets we have exactly  $d + 1$  vertices. ■

By Proposition 2, if  $d = 2$  then  $n(d, k) = dk$ . However, for higher degrees we get  $n(d, k) \sim \frac{1}{3}dk$ . Denote by  $n_{\text{VT}}(d, k)$  the minimum number of vertices in a vertex-transitive  $d$ -regular graph with diameter  $k$ . As shown in [6], for  $k \geq 4$  and “large”  $d$  we have  $n_{\text{VT}}(d, k) \sim \frac{2}{3}dk$ , and so  $n_{\text{VT}}(d, k) \doteq 2n(d, k)$  in this case. On the other hand, since the extremal graphs constructed in the proof of Proposition 1 are vertex-transitive, we have  $n_{\text{VT}}(d, k) = n(d, k)$  when  $k \leq 3$ .

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