# CONNECTED GLOBAL OFFENSIVE $\boldsymbol{k}$-ALLIANCES IN GRAPHS 

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#### Abstract

We consider finite graphs $G$ with vertex set $V(G)$. For a subset $S \subseteq V(G)$, we define by $G[S]$ the subgraph induced by $S$. By $n(G)=$ $|V(G)|$ and $\delta(G)$ we denote the order and the minimum degree of $G$, respectively. Let $k$ be a positive integer. A subset $S \subseteq V(G)$ is a connected global offensive $k$-alliance of the connected graph $G$, if $G[S]$ is connected and $|N(v) \cap S| \geq|N(v)-S|+k$ for every vertex $v \in V(G)-S$, where $N(v)$ is the neighborhood of $v$. The connected global offensive $k$-alliance number $\gamma_{o}^{k, c}(G)$ is the minimum cardinality of a connected global offensive $k$-alliance in $G$.

In this paper we characterize connected graphs $G$ with $\gamma_{o}^{k, c}(G)=$ $n(G)$. In the case that $\delta(G) \geq k \geq 2$, we also characterize the family of connected graphs $G$ with $\gamma_{o}^{k, c}(G)=n(G)-1$. Furthermore, we present different tight bounds of $\gamma_{o}^{k, c}(G)$. Keywords: alliances in graphs, connected global offensive $k$-alliance, global offensive $k$-alliance, domination. 2010 Mathematics Subject Classification: 05C69.


## 1. TERMINOLOGY AND Introduction

We consider finite, undirected and simple graphs $G$ with vertex set $V(G)$. The number of vertices $|V(G)|$ of a graph $G$ is called the order and is denoted by $n=n(G)$. The neighborhood $N(v)=N_{G}(v)$ of a vertex $v$ consists of the
vertices adjacent to $v$ and $d(v)=d_{G}(v)=|N(v)|$ is the degree of $v$. By $\delta=\delta(G)$ and $\Delta=\Delta(G)$, we denote the minimum degree and the maximum degree of the graph $G$, respectively. For a subset $S \subseteq V(G)$, we define by $G[S]$ the subgraph induced by $S$.

The complete graph of order $n$ is denoted by $K_{n}$, and $K_{s, t}$ is the complete bipartite graph with the two parts of cardinality $s$ and $t$.

Two vertices that are not adjacent in a graph $G$ are said to be independent. A set $I$ of vertices is independent if every two vertices of $I$ are independent. The independence number $\alpha(G)$ of a graph $G$ is the maximum cardinality among the independent sets of vertices of $G$.

A vertex-cut in a connected graph $G$ is a set $S$ of vertices of $G$ such that $G-S$ is disconnected. The connectivity $\kappa(G)$ of a graph $G$ is the minimum cardinality of a vertex-cut of $G$ if $G$ is not complete, and $\kappa(G)=n-1$ if $G$ is isomorphic to the complete graph $K_{n}$.

Kristiansen, Hedetniemi and Hedetniemi [9] introduced several types of alliances in graphs, including defensive and offensive alliances. As a generalization of the offensive alliance, Shafique and Dutton [11, 12] defined the global offensive $k$-alliance for a positive integer $k$ as follows. A subset $S \subseteq V(G)$ is a global offensive $k$-alliance of the graph $G$ if $|N(v) \cap S| \geq$ $|N(v)-S|+k$ for every vertex $v \in V(G)-S$. The global offensive $k$-alliance number $\gamma_{o}^{k}(G)$ is the minimum cardinality of a global offensive $k$-alliance in $G$. A global offensive $k$-alliance set of the minimum cardinality of a graph $G$ is called a $\gamma_{o}^{k}(G)$-set. Results on global offensive $k$-alliances were given, for example, by Bermudo, Rodríguez-Velázquez, Sigarreta and Yero [1], Chellali [2], Chellali, Haynes, Randerath and Volkmann [3] and Fernau, Rodríguez and Sigarreta [4].

In this paper, we are interested in connected global offensive $k$-alliances. Analogously to the definition above, a subset $S \subseteq V(G)$ is a connected global offensive $k$-alliance of the connected graph $G$, if $G[S]$ is connected and $|N(v) \cap S| \geq|N(v)-S|+k$ for every vertex $v \in V(G)-S$. The connected global offensive $k$-alliance number $\gamma_{o}^{k, c}(G)$ is the minimum cardinality of a connected global offensive $k$-alliance in $G$. A connected global offensive $k$ alliance set of the minimum cardinality of a connected graph $G$ is called a $\gamma_{o}^{k, c}(G)$-set.

A subset $D \subseteq V(G)$ is a $k$-dominating set of the graph $G$ if $\left|N_{G}(v) \cap D\right| \geq$ $k$ for every $v \in V(G)-D$. The $k$-domination number $\gamma^{k}(G)$ is the minimum cardinality among the $k$-dominating sets of $G$. Note that the 1-domination number $\gamma^{1}(G)$ is the usual domination number $\gamma(G)$. A subset $D \subseteq V(G)$ is
a connected $k$-dominating set of a connected graph $G$, if $D$ is a $k$-dominating set of $G$ and the induced subgraph $G[D]$ is connected. The connected $k$ domination number $\gamma^{k, c}(G)$ is the minimum cardinality among the connected $k$-dominating sets of $G$.

In $[5,6]$, Fink and Jacobson introduced the concept of $k$-domination. For a comprehensive treatment of domination in graphs, see the monographs by Haynes, Hedetniemi and Slater [7, 8].

In this paper we characterize the connected graphs $G$ with $\gamma_{o}^{k, c}(G)=$ $n(G)$. If $G$ is a connected graph with $\delta(G) \geq k \geq 3$, then we show that $\gamma_{o}^{k, c}(G)=n(G)-1$ if and only if $G$ is isomorphic to the complete graph $K_{k+1}$ or $K_{k+2}$. In addition, we derive different sharp bounds on $\gamma_{o}^{k, c}(G)$, as for example, $\gamma_{o}^{k, c}(G) \leq 2 \gamma_{o}^{k}(G)-k+1$.

## 2. Main Results

Observation 1. If $k \geq 1$ is an integer, then $\gamma_{o}^{k, c}(G) \geq \gamma^{k, c}(G)$ for any connected graph $G$.

Proof. If $S$ is a $\gamma_{o}^{k, c}(G)$-set, then $G[S]$ is connected and every vertex of $V(G)-S$ has at least $k$ neighbors in $S$. Thus $S$ is a connected $k$-dominating set of $G$ and so $\gamma^{k, c}(G) \leq|S|=\gamma_{o}^{k, c}(G)$.

In view of Observation 1, each lower bound of $\gamma^{k, c}(G)$ is also a lower bound of $\gamma_{o}^{k, c}(G)$. Now we characterize all connected graphs $G$ with the property that $\gamma_{o}^{k, c}(G)=n(G)$.

Observation 2. Let $k \geq 2$ be an integer, and let $G$ be a connected graph of order $n \geq 2$. Then $\gamma_{o}^{k, c}(G)=n$ if and only if all vertices of $G$ are either cut-vertices or vertices of degree less than $k$.

Proof. If each vertex of $G$ is either a cut-vertex or has degree less than $k$, then the definition of the connected global offensive $k$-alliance number leads to $\gamma_{o}^{k, c}(G)=n$ immediately.

Conversely, assume that $\gamma_{o}^{k, c}(G)=n$. Suppose to the contrary that $G$ contains a non-cut-vertex $u$ with $d_{G}(u) \geq k$. This implies that $G-u$ is a connected graph. Since $d_{G}(u) \geq k$, we deduce that $V(G-u)$ is a connected global offensive $k$-alliance of $G$. Therefore we obtain the contradiction $\gamma_{o}^{k, c}(G) \leq n-1$, and the proof is complete.

Corollary 3. Let $k \geq 2$ be an integer. If $T$ is a tree, then $\gamma_{o}^{k, c}(T)=n(T)$.
Corollary 4. If $k \geq 2$ is an integer, and $G$ is a connected graph with $\delta(G) \geq k$, then $\gamma_{o}^{k, c}(G) \leq n(G)-1$.
Next we derive a characterization of all connected graphs $G$ with $\gamma_{o}^{k, c}(G)=$ $n(G)-1$ when $\delta(G) \geq k \geq 2$.
Theorem 5. Let $k \geq 2$ be an integer, and let $G$ be a connected graph of order $n$ and minimum degree $\delta$.
(i) If $\delta \geq 2$, then $\gamma_{o}^{2, c}(G)=n-1$ if and only if $G$ is a cycle or $G$ is isomorphic to the complete graph $K_{4}$.
(ii) If $\delta \geq k \geq 3$, then $\gamma_{o}^{k, c}(G)=n-1$ if and only if $G$ is isomorphic to the complete graph $K_{k+1}$ or $K_{k+2}$.
Proof. Obviously, if $G$ is a cycle or $G$ is isomorphic to $K_{4}$, then $\gamma_{o}^{2, c}(G)=$ $n-1$, and if $G$ is isomorphic to the complete graphs $K_{k+1}$ or $K_{k+2}$, then $\gamma_{o}^{k, c}(G)=n-1$.

Conversely, assume that $\gamma_{o}^{k, c}(G)=n-1$, and let $P=u_{1} u_{2} \ldots u_{t}$ be the longest path in $G$. The condition $\delta \geq k \geq 2$ implies that $u_{1} \neq u_{t}$ and $G-\left\{u_{1}, u_{t}\right\}$ is a connected subgraph of $G$. If $u_{1}$ and $u_{t}$ are not adjacent in $G$, then we arrive at the contradiction that $V(G)-\left\{u_{1}, u_{t}\right\}$ is a connected global offensive $k$-alliance of $G$. In the remaining case that $u_{1}$ and $u_{t}$ are adjacent in $G$, we observe that $C=u_{1} u_{2} \ldots u_{t} u_{1}$ is a Hamiltonian cycle of $G$, because $P$ is the longest path in $G$. This yields $t=n$.
(i) Assume that $k=2$. Suppose that the Hamiltonian cycle $C=$ $u_{1} u_{2} \ldots u_{n} u_{1}$ has a chord. If, without loss of generality, $u_{1} u_{s}$ with $3 \leq s \leq$ $n-1$ is a chord of $C$, then we obtain the contradiction that $V(G)-\left\{u_{2}, u_{n}\right\}$ is a connected global offensive 2 -alliance of $G$ or $u_{2}$ and $u_{n}$ are adjacent. Therefore assume in the following that $u_{2}$ and $u_{n}$ are adjacent. If $n=4$, then $G=K_{4}$. If $n \geq 5$, then we distinguish the cases $s=3$ and $s \geq 4$.

Assume first that $s=3$. Then we obtain the contradiction that $V(G)-$ $\left\{u_{2}, u_{4}\right\}$ is a connected global offensive 2 -alliance of $G$ or $u_{2}$ and $u_{4}$ are adjacent. If $u_{2}$ and $u_{4}$ are adjacent, then we have the contradiction that $V(G)-\left\{u_{3}, u_{n}\right\}$ is a connected global offensive 2-alliance of $G$ or $u_{3}$ and $u_{n}$ are adjacent. However, if $u_{3}$ and $u_{n}$ are adjacent, then $d_{G}\left(u_{2}\right), d_{G}\left(u_{n}\right) \geq 4$, and thus we arrive at the contradiction that $V(G)-\left\{u_{2}, u_{n}\right\}$ is a connected global offensive 2-alliance of $G$.

Assume now that $s \geq 4$. Then we obtain the contradiction that $V(G)-$ $\left\{u_{1}, u_{3}\right\}$ is a connected global offensive 2-alliance of $G$ or $u_{1}$ and $u_{3}$ are
adjacent. If $u_{1}$ and $u_{3}$ are adjacent, then we have the contradiction that $V(G)-\left\{u_{3}, u_{n}\right\}$ is a connected global offensive 2-alliance of $G$ or $u_{3}$ and $u_{n}$ are adjacent. However, if $u_{3}$ and $u_{n}$ are adjacent, then $d_{G}\left(u_{1}\right), d_{G}\left(u_{n}\right) \geq 4$, and thus we arrive at the contradiction that $V(G)-\left\{u_{1}, u_{n}\right\}$ is a connected global offensive 2 -alliance of $G$.
(ii) Assume that $k \geq 3$. In the following all indices are taken modulo $n$. If the vertices $u_{i}$ and $u_{i+2}$ are not adjacent for any index $i$ with $1 \leq$ $i \leq n$, then the hypothesis $\delta \geq k \geq 3$ leads to the contradiction that $V(G)-\left\{u_{i}, u_{i+2}\right\}$ is a connected global offensive $k$-alliance of $G$. Hence assume that $u_{i}$ and $u_{i+2}$ are adjacent for each index $i \in\{1,2, \ldots, n\}$. Now let $s$ be an arbitrary integer with $3 \leq s \leq n-3$. If $u_{i}$ and $u_{i+s}$ are not adjacent, then $V(G)-\left\{u_{i}, u_{i+s}\right\}$ is a connected global offensive $k$-alliance of $G$, since there exists the edge $u_{i-1} u_{i+1}$ in $G$. Therefore it remains the case that $G$ is a complete graph. If $G$ is isomorphic to $K_{k+1}$ or $K_{k+2}$, then $\gamma_{o}^{k, c}(G)=n-1$. However, if $G$ is isomorphic to $K_{q}$ for any integer $q \geq k+3$, then $V(G)-\left\{u_{1}, u_{2}\right\}$ is a connected global offensive $k$-alliance of $G$. This contradiction completes the proof of Theorem 5.

Proposition 6. Let $G$ be a graph of order $n$, and let $k, p$ be two integers such that $k \geq 1$ and $-1 \leq p \leq \alpha(G)-2$. If $\delta(G) \geq k$ and $\kappa(G) \geq \alpha(G)-p$, then

$$
\gamma_{o}^{k, c}(G) \leq n(G)-\alpha(G)+p+1 .
$$

Proof. Let $I \subset V(G)$ be an independent set of cardinality $\alpha(G)-p-1$. The hypothesis $\kappa(G) \geq \alpha(G)-p$ implies that $G[V(G)-I]$ is connected. Since $I$ is an independent set, the condition $\delta(G) \geq k$ shows that each vertex in $I$ has at least $k$ neighbors in $V(G)-I$. Thus $V(G)-I$ is a connected global offensive $k$-alliance of $G$ such that $|V(G)-I| \leq n-(\alpha(G)-p-1)$, and the proof is complete.

If $H$ is the complete bipartite graph $K_{k, k}$, then $\delta(H)=\alpha(H)=\kappa(H)=k$ and $\gamma_{o}^{k, c}(H)=k+1=n(H)-\alpha(H)+1$. This example demonstrates that Proposition 6 is the best possible, at least for $p=0$.

Theorem 7. Let $G$ be a connected graph and $k$ an integer with $1 \leq k \leq$ $\Delta(G)$. Then

$$
\gamma_{o}^{k, c}(G) \leq 2 \gamma_{o}^{k}(G)-k+1 .
$$

Proof. Let $S$ be a $\gamma_{o}^{k}(G)$-set. Since $k \leq \Delta(G)$, we observe that $|S|=$ $\gamma_{o}^{k}(G) \leq n(G)-1$. Now let $x \in V(G)-S$ be an arbitrary vertex.

If $G[S \cup\{x\}]$ is connected, then the inequality $k \leq \gamma_{o}^{k}(G)$ implies that $\gamma_{o}^{k, c}(G) \leq \gamma_{o}^{k}(G)+1 \leq 2 \gamma_{o}^{k}(G)-k+1$, and we are done.

Thus assume next that $G[S \cup\{x\}]$ is not connected. We will add successively vertices from $V(G)-(S \cup\{x\})$ to $S \cup\{x\}$ in order to decrease the number of components, at least one in each step, until we obtain a set of vertices whose induced subgraph is connected. Note that if we partition $S \cup\{x\}$ into two parts $A$ and $B$ such that there is no edge between $A$ and $B$, and we take vertices $a \in A$ and $b \in B$ such that the distance between $a$ and $b$ is minimum in $G$, then the property of $S$ of being dominating implies that $d_{G}(a, b) \leq 3$. It follows that in each step of increasing $S \cup\{x\}$ we need to add at most 2 vertices from $V(G)-(S \cup\{x\})$. Let $r_{1}$ and $r_{2}$ be the number of steps where we include one vertex and two vertices from $V(G)-(S \cup\{x\})$, respectively, and define $r=r_{1}+r_{2}$. Let $S_{0} \subset S \cup\{x\}$ be the set of vertices of the component of $G[S \cup\{x\}]$ to which $x$ belongs, and let $S_{i} \subset S$ be the set of vertices connected to $\bigcup_{j=0}^{i-1} S_{j}$ in step $i \geq 1$. Clearly, $\left|S_{0}\right| \geq k+1$ and $\left|S_{i}\right| \geq 1$ for $1 \leq i \leq r$. Furthermore, since $S$ is a global offensive $k$-alliance, in the steps where two vertices from $V(G)-(S \cup\{x\})$ are added, we observe that $\left|S_{i}\right| \geq k+1$. This leads to

$$
\gamma_{o}^{k}(G)=|S|=\left|S_{0}-\{x\}\right|+\sum_{i=1}^{r}\left|S_{i}\right| \geq k+r_{2}(k+1)+r_{1}
$$

and therefore $r_{1} \leq \gamma_{o}^{k}(G)-k-r_{2}(k+1)$. As a further consequence, we see that $S \cup\{x\}$ together with all vertices from $V(G)-(S \cup\{x\})$ added in steps 1 to $r$ form a connected global offensive $k$-alliance of $G$. Altogether, we deduce that

$$
\begin{aligned}
\gamma_{o}^{k, c}(G) & \leq|S|+1+r_{1}+2 r_{2} \\
& \leq \gamma_{o}^{k}(G)+1+\gamma_{o}^{k}(G)-k-r_{2}(k+1)+2 r_{2} \\
& =2 \gamma_{o}^{k}(G)-k+1-r_{2}(k+1)+2 r_{2} \\
& \leq 2 \gamma_{o}^{k}(G)-k+1
\end{aligned}
$$

and the proof is complete.
If $H$ is the complete bipartite graph $K_{k, p}$, then $\gamma_{o}^{k}(H)=k$ and $\gamma_{o}^{k, c}(H)=$ $k+1$. This example shows that the bound given in Theorem 7 is tight.

Theorem 8. Let $G$ be a connected graph and $k \geq 1$ an integer. If $\delta(G) \geq$ $k+1$, then

$$
\gamma_{o}^{k+1, c}(G) \leq \frac{\gamma_{o}^{k, c}(G)+n(G)}{2}
$$

Proof. Let $S$ be a $\gamma_{o}^{k, c}(G)$-set, and let $A$ be the set of isolated vertices in the subgraph $G-S$. Then the subgraph $G-(S \cup A)$ contains no isolated vertices. If $D$ is a minimum dominating set of $G-(S \cup A)$, then the wellknown inequality of Ore [10] implies

$$
|D| \leq \frac{|V(G)-(S \cup A)|}{2} \leq \frac{|V(G)-S|}{2}=\frac{n(G)-\gamma_{o}^{k, c}(G)}{2} .
$$

If $S^{\prime}=S \cup D$, then $G\left[S^{\prime}\right]$ is connected. In addition, for each vertex $v \in$ $V(G)-\left(S^{\prime} \cup A\right)$, we have

$$
\begin{aligned}
\left|N(v) \cap S^{\prime}\right| & =|N(v) \cap S|+|N(v) \cap D| \\
& \geq|N(v)-S|+k+1 \\
& =\left|N(v)-S^{\prime}\right|+|N(v) \cap D|+k+1 \\
& \geq\left|N(v)-S^{\prime}\right|+k+2 .
\end{aligned}
$$

Since $\delta(G) \geq k+1$, every vertex of $A$ has at least $k+1$ neighbors in $S$, and therefore $S^{\prime}$ is a connected global offensive ( $k+1$ )-alliance of $G$ and thus

$$
\begin{aligned}
\gamma_{o}^{k+1, c}(G) & \leq\left|S^{\prime}\right|=|S|+|D|=\gamma_{o}^{k, c}(G)+|D| \\
& \leq \gamma_{o}^{k, c}(G)+\frac{n(G)-\gamma_{o}^{k, c}(G)}{2}=\frac{n(G)+\gamma_{o}^{k, c}(G)}{2} .
\end{aligned}
$$

The inequality $\left|N(v) \cap S^{\prime}\right| \geq\left|N(v)-S^{\prime}\right|+k+2$ for each vertex $v \in V(G)-$ ( $S^{\prime} \cup A$ ) in the proof of Theorem 8 leads immediately to the next result.

Theorem 9. Let $G$ be a connected graph and $k \geq 1$ an integer. If $\delta(G) \geq$ $k+2$, then

$$
\gamma_{o}^{k+2, c}(G) \leq \frac{\gamma_{o}^{k, c}(G)+n(G)}{2}
$$

If $H=K_{k+3}$, then $\gamma_{o}^{k+2, c}(H)=\gamma_{o}^{k+1, c}(H)=k+2$ and $\gamma_{o}^{k, c}(H)=k+1$ and thus

$$
\gamma_{o}^{k+2, c}(H)=\gamma_{o}^{k+1, c}(H)=k+2=\frac{\gamma_{o}^{k, c}(H)+n(H)}{2} .
$$

Let $k \geq 2$ be an even integer, and let $F=K_{k+6}-M$, where $M$ is a perfect matching of the complete graph $K_{k+6}$. Then $\gamma_{o}^{k+2, c}(F)=\gamma_{o}^{k+1, c}(F)=k+4$ and $\gamma_{o}^{k, c}(F)=k+2$, and so

$$
\gamma_{o}^{k+2, c}(F)=\gamma_{o}^{k+1, c}(F)=k+4=\frac{\gamma_{o}^{k, c}(F)+n(F)}{2} .
$$

These two graphs $H$ and $F$ demonstrate that Theorem 8 as well as Theorem 9 are the best possible.

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